
**Information technology — Coding of
moving pictures and associated audio for
digital storage media at up to about
1,5 Mbit/s —**

Part 2:
Video

*Technologies de l'information — Codage de l'image animée et du son
associé pour les supports de stockage numérique jusqu'à environ
1,5 Mbit/s —*

Partie 2: Vidéo



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Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work.

In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

International Standard ISO/IEC 11172-2 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Sub-Committee SC 29, *Coded representation of audio, picture, multimedia and hypermedia information*.

ISO/IEC 11172 consists of the following parts, under the general title *Information technology — Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s*:

- Part 1: *Systems*
- Part 2: *Video*
- Part 3: *Audio*
- Part 4: *Compliance testing*

Annexes A, B and C form an integral part of this part of ISO/IEC 11172. Annexes D, E and F are for information only.

Introduction

Note -- Readers interested in an overview of the MPEG Video layer should read this Introduction and then proceed to annex D, before returning to clauses 1 and 2.

0.1 Purpose

This part of ISO/IEC 11172 was developed in response to the growing need for a common format for representing compressed video on various digital storage media such as CDs, DATs, Winchester disks and optical drives. This part of ISO/IEC 11172 specifies a coded representation that can be used for compressing video sequences to bitrates around 1,5 Mbit/s. The use of this part of ISO/IEC 11172 means that motion video can be manipulated as a form of computer data and can be transmitted and received over existing and future networks. The coded representation can be used with both 625-line and 525-line television and provides flexibility for use with workstation and personal computer displays.

This part of ISO/IEC 11172 was developed to operate principally from storage media offering a continuous transfer rate of about 1,5 Mbit/s. Nevertheless it can be used more widely than this because the approach taken is generic.

0.1.1 Coding parameters

The intention in developing this part of ISO/IEC 11172 has been to define a source coding algorithm with a large degree of flexibility that can be used in many different applications. To achieve this goal, a number of the parameters defining the characteristics of coded bitstreams and decoders are contained in the bitstream itself. This allows for example, the algorithm to be used for pictures with a variety of sizes and aspect ratios and on channels or devices operating at a wide range of bitrates.

Because of the large range of the characteristics of bitstreams that can be represented by this part of ISO/IEC 11172, a sub-set of these coding parameters known as the "Constrained Parameters" has been defined. The aim in defining the constrained parameters is to offer guidance about a widely useful range of parameters. Conforming to this set of constraints is not a requirement of this part of ISO/IEC 11172. A flag in the bitstream indicates whether or not it is a Constrained Parameters bitstream.

Summary of the Constrained Parameters:

Horizontal picture size	Less than or equal to 768 pels
Vertical picture size	Less than or equal to 576 lines
Picture area	Less than or equal to 396 macroblocks
Pel rate	Less than or equal to 396x25 macroblocks/s
Picture rate	Less than or equal to 30 Hz
Motion vector range	Less than -64 to +63,5 pels (using half-pel vectors) [backward_f_code and forward_f_code ≤ 4 (see table D.7)]
Input buffer size (in VBV model)	Less than or equal to 327 680 bits
Bitrate	Less than or equal to 1 856 000 bits/s (constant bitrate)

0.2 Overview of the algorithm

The coded representation defined in this part of ISO/IEC 11172 achieves a high compression ratio while preserving good picture quality. The algorithm is not lossless as the exact pel values are not preserved during coding. The choice of the techniques is based on the need to balance a high picture quality and compression ratio with the requirement to make random access to the coded bitstream. Obtaining good picture quality at the bitrates of interest demands a very high compression ratio, which is not achievable with intraframe coding alone. The need for random access, however, is best satisfied with pure intraframe coding. This requires a careful balance between intra- and interframe coding and between recursive and non-recursive temporal redundancy reduction.

A number of techniques are used to achieve a high compression ratio. The first, which is almost independent from this part of ISO/IEC 11172, is to select an appropriate spatial resolution for the signal. The algorithm then uses block-based motion compensation to reduce the temporal redundancy. Motion compensation is used for causal prediction of the current picture from a previous picture, for non-causal prediction of the current picture from a future picture, or for interpolative prediction from past and future pictures. Motion vectors are defined for each 16-pel by 16-line region of the picture. The difference signal, the prediction error, is further compressed using the discrete cosine transform (DCT) to remove spatial correlation before it is quantized in an irreversible process that discards the less important information. Finally, the motion vectors are combined with the DCT information, and coded using variable length codes.

0.2.1 Temporal processing

Because of the conflicting requirements of random access and highly efficient compression, three main picture types are defined. Intra-coded pictures (I-Pictures) are coded without reference to other pictures. They provide access points to the coded sequence where decoding can begin, but are coded with only a moderate compression ratio. Predictive coded pictures (P-Pictures) are coded more efficiently using motion compensated prediction from a past intra or predictive coded picture and are generally used as a reference for further prediction. Bidirectionally-predictive coded pictures (B-Pictures) provide the highest degree of compression but require both past and future reference pictures for motion compensation. Bidirectionally-predictive coded pictures are never used as references for prediction. The organisation of the three picture types in a sequence is very flexible. The choice is left to the encoder and will depend on the requirements of the application. Figure 1 illustrates the relationship between the three different picture types.

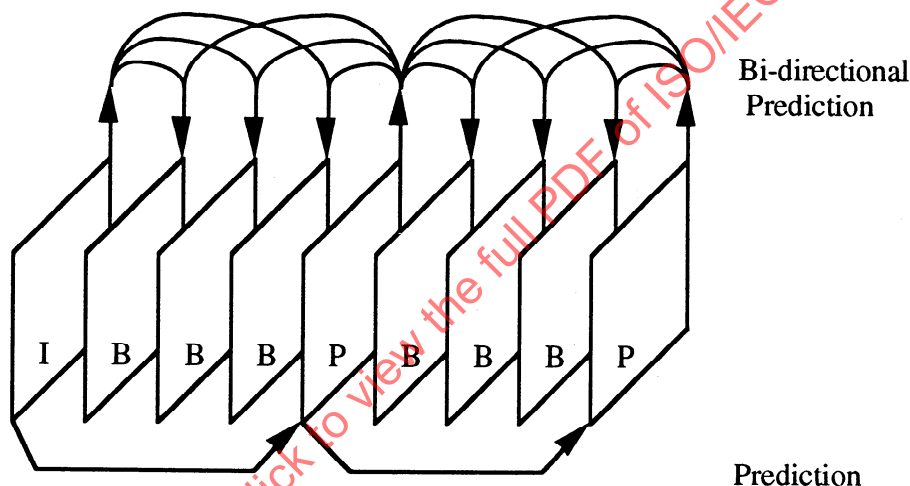


Figure 1 -- Example of temporal picture structure

The fourth picture type defined in this part of ISO/IEC 11172, the D-picture, is provided to allow a simple, but limited quality, fast-forward playback mode.

0.2.2 Motion representation - macroblocks

The choice of 16 by 16 macroblocks for the motion-compensation unit is a result of the trade-off between increasing the coding efficiency provided by using motion information and the overhead needed to store it. Each macroblock can be one of a number of different types. For example, intra-coded, forward-predictive-coded, backward-predictive coded, and bidirectionally-predictive-coded macroblocks are permitted in bidirectionally-predictive coded pictures. Depending on the type of the macroblock, motion vector information and other side information are stored with the compressed prediction error signal in each macroblock. The motion vectors are encoded differentially with respect to the last coded motion vector, using variable-length codes. The maximum length of the vectors that may be represented can be programmed, on a picture-by-picture basis, so that the most demanding applications can be met without compromising the performance of the system in more normal situations.

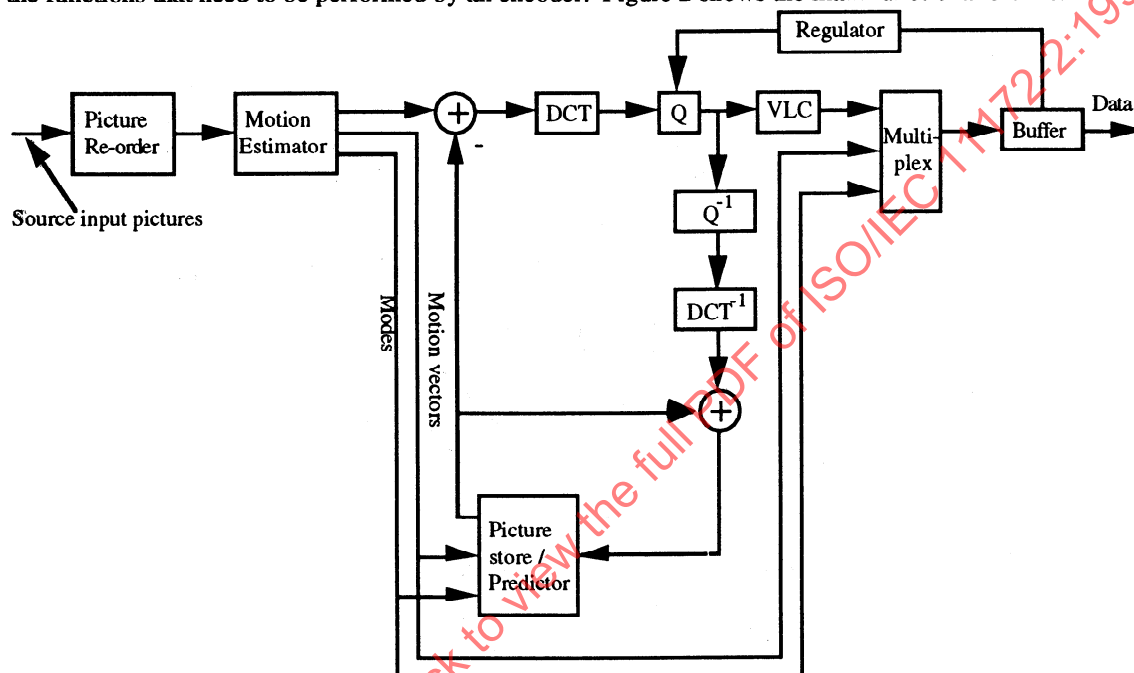
It is the responsibility of the encoder to calculate appropriate motion vectors. This part of ISO/IEC 11172 does not specify how this should be done.

0.2.3 Spatial redundancy reduction

Both original pictures and prediction error signals have high spatial redundancy. This part of ISO/IEC 11172 uses a block-based DCT method with visually weighted quantization and run-length coding. Each 8 by 8 block of the original picture for intra-coded macroblocks or of the prediction error for predictive-coded macroblocks is transformed into the DCT domain where it is scaled before being quantized. After quantization many of the coefficients are zero in value and so two-dimensional run-length and variable length coding is used to encode the remaining coefficients efficiently.

0.3 Encoding

This part of ISO/IEC 11172 does not specify an encoding process. It specifies the syntax and semantics of the bitstream and the signal processing in the decoder. As a result, many options are left open to encoders to trade-off cost and speed against picture quality and coding efficiency. This clause is a brief description of the functions that need to be performed by an encoder. Figure 2 shows the main functional blocks.



where

DCT is discrete cosine transform
 DCT^{-1} is inverse discrete cosine transform
 Q is quantization
 Q^{-1} is dequantization
 VLC is variable length coding

Figure 2 -- Simplified video encoder block diagram

The input video signal must be digitized and represented as a luminance and two colour difference signals (Y , C_b , C_r). This may be followed by preprocessing and format conversion to select an appropriate window, resolution and input format. This part of ISO/IEC 11172 requires that the colour difference signals (C_b and C_r) are subsampled with respect to the luminance by 2:1 in both vertical and horizontal directions and are reformatted, if necessary, as a non-interlaced signal.

The encoder must choose which picture type to use for each picture. Having defined the picture types, the encoder estimates motion vectors for each 16 by 16 macroblock in the picture. In P-Pictures one vector is needed for each non-intra macroblock and in B-Pictures one or two vectors are needed.

If B-Pictures are used, some reordering of the picture sequence is necessary before encoding. Because B-Pictures are coded using bidirectional motion compensated prediction, they can only be decoded after the subsequent reference picture (an I or P-Picture) has been decoded. Therefore the pictures are reordered by the

encoder so that the pictures arrive at the decoder in the order for decoding. The correct display order is recovered by the decoder.

The basic unit of coding within a picture is the macroblock. Within each picture, macroblocks are encoded in sequence, left to right, top to bottom. Each macroblock consists of six 8 by 8 blocks: four blocks of luminance, one block of Cb chrominance, and one block of Cr chrominance. See figure 3. Note that the picture area covered by the four blocks of luminance is the same as the area covered by each of the chrominance blocks. This is due to subsampling of the chrominance information to match the sensitivity of the human visual system.

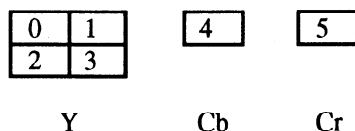


Figure 3 -- Macroblock structure

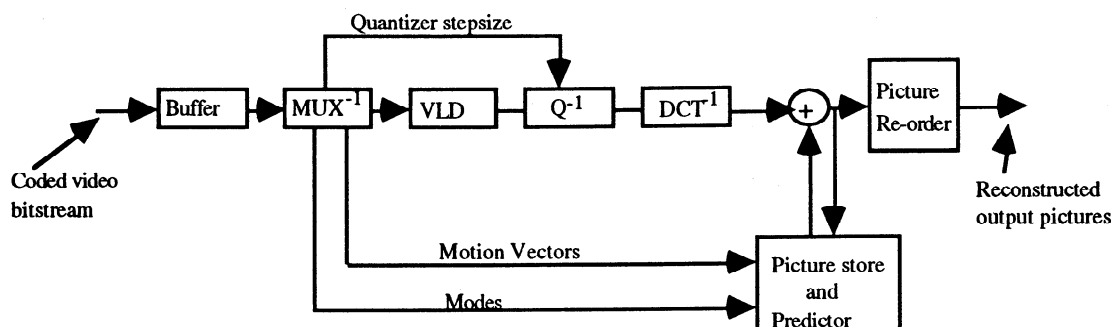
Firstly, for a given macroblock, the coding mode is chosen. It depends on the picture type, the effectiveness of motion compensated prediction in that local region, and the nature of the signal within the block. Secondly, depending on the coding mode, a motion compensated prediction of the contents of the block based on past and/or future reference pictures is formed. This prediction is subtracted from the actual data in the current macroblock to form an error signal. Thirdly, this error signal is separated into 8 by 8 blocks (4 luminance and 2 chrominance blocks in each macroblock) and a discrete cosine transform is performed on each block. Each resulting 8 by 8 block of DCT coefficients is quantized and the two-dimensional block is scanned in a zig-zag order to convert it into a one-dimensional string of quantized DCT coefficients. Fourthly, the side-information for the macroblock (mode, motion vectors etc) and the quantized coefficient data are encoded. For maximum efficiency, a number of variable length code tables are defined for the different data elements. Run-length coding is used for the quantized coefficient data.

A consequence of using different picture types and variable length coding is that the overall data rate is variable. In applications that involve a fixed-rate channel, a FIFO buffer may be used to match the encoder output to the channel. The status of this buffer may be monitored to control the number of bits generated by the encoder. Controlling the quantization process is the most direct way of controlling the bitrate. This part of ISO/IEC 11172 specifies an abstract model of the buffering system (the Video Buffering Verifier) in order to constrain the maximum variability in the number of bits that are used for a given picture. This ensures that a bitstream can be decoded with a buffer of known size.

At this stage, the coded representation of the picture has been generated. The final step in the encoder is to regenerate I-Pictures and P-Pictures by decoding the data so that they can be used as reference pictures for subsequent encoding. The quantized coefficients are dequantized and an inverse 8 by 8 DCT is performed on each block. The prediction error signal produced is then added back to the prediction signal and limited to the required range to give a decoded reference picture.

0.4 Decoding

Decoding is the inverse of the encoding operation. It is considerably simpler than encoding as there is no need to perform motion estimation and there are many fewer options. The decoding process is defined by this part of ISO/IEC 11172. The description that follows is a very brief overview of one possible way of decoding a bitstream. Other decoders with different architectures are possible. Figure 4 shows the main functional blocks.



Where

DCT^{-1} is inverse discrete cosine transform

Q^{-1} is dequantization

MUX^{-1} is demultiplexing

VLD is variable length decoding

Figure 4 -- Basic video decoder block diagram

For fixed-rate applications, the channel fills a FIFO buffer at a constant rate with the coded bitstream. The decoder reads this buffer and decodes the data elements in the bitstream according to the defined syntax.

As the decoder reads the bitstream, it identifies the start of a coded picture and then the type of the picture. It decodes each macroblock in the picture in turn. The macroblock type and the motion vectors, if present, are used to construct a prediction of the current macroblock based on past and future reference pictures that have been stored in the decoder. The coefficient data are decoded and dequantized. Each 8 by 8 block of coefficient data is transformed by an inverse DCT (specified in annex A), and the result is added to the prediction signal and limited to the defined range.

After all the macroblocks in the picture have been processed, the picture has been reconstructed. If it is an I-picture or a P-picture it is a reference picture for subsequent pictures and is stored, replacing the oldest stored reference picture. Before the pictures are displayed they may need to be re-ordered from the coded order to their natural display order. After reordering, the pictures are available, in digital form, for post-processing and display in any manner that the application chooses.

0.5 Structure of the coded video bitstream

This part of ISO/IEC 11172 specifies a syntax for a coded video bitstream. This syntax contains six layers, each of which either supports a signal processing or a system function:

Layers of the syntax	Function
Sequence layer	Random access unit: context
Group of pictures layer	Random access unit: video
Picture layer	Primary coding unit
Slice layer	Resynchronization unit
Macroblock layer	Motion compensation unit
Block layer	DCT unit

0.6 Features supported by the algorithm

Applications using compressed video on digital storage media need to be able to perform a number of operations in addition to normal forward playback of the sequence. The coded bitstream has been designed to support a number of these operations.

0.6.1 Random access

Random access is an essential feature for video on a storage medium. Random access requires that any picture can be decoded in a limited amount of time. It implies the existence of access points in the bitstream - that is segments of information that are identifiable and can be decoded without reference to other segments of data. A spacing of two random access points (Intra-Pictures) per second can be achieved without significant loss of picture quality.

0.6.2 Fast search

Depending on the storage medium, it is possible to scan the access points in a coded bitstream (with the help of an application-specific directory or other knowledge beyond the scope of this part of ISO/IEC 11172) to obtain a fast-forward and fast-reverse playback effect.

0.6.3 Reverse playback

Some applications may require the video signal to be played in reverse order. This can be achieved in a decoder by using memory to store entire groups of pictures after they have been decoded before being displayed in reverse order. An encoder can make this feature easier by reducing the length of groups of pictures.

0.6.4 Error robustness

Most digital storage media and communication channels are not error-free. Appropriate channel coding schemes should be used and are beyond the scope of this part of ISO/IEC 11172. Nevertheless the compression scheme defined in this part of ISO/IEC 11172 is robust to residual errors. The slice structure allows a decoder to recover after a data error and to resynchronize its decoding. Therefore, bit errors in the compressed data will cause errors in the decoded pictures to be limited in area. Decoders may be able to use concealment strategies to disguise these errors.

0.6.5 Editing

There is a conflict between the requirement for high coding efficiency and easy editing. The coding structure and syntax have not been designed with the primary aim of simplifying editing at any picture. Nevertheless a number of features have been included that enable editing of coded data.

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Information technology — Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s —

Part 2: Video

Section 1: General

1.1 Scope

This part of ISO/IEC 11172 specifies the coded representation of video for digital storage media and specifies the decoding process. The representation supports normal speed forward playback, as well as special functions such as random access, fast forward playback, fast reverse playback, normal speed reverse playback, pause and still pictures. This part of ISO/IEC 11172 is compatible with standard 525- and 625-line television formats, and it provides flexibility for use with personal computer and workstation displays.

ISO/IEC 11172 is primarily applicable to digital storage media supporting a continuous transfer rate up to about 1,5 Mbit/s, such as Compact Disc, Digital Audio Tape, and magnetic hard disks. Nevertheless it can be used more widely than this because of the generic approach taken. The storage media may be directly connected to the decoder, or via communications means such as busses, LANs, or telecommunications links. This part of ISO/IEC 11172 is intended for non-interlaced video formats having approximately 288 lines of 352 pels and picture rates around 24 Hz to 30 Hz.

1.2 Normative references

The following International Standards contain provisions which, through reference in this text, constitute provisions of this part of ISO/IEC 11172. At the time of publication, the editions indicated were valid. All standards are subject to revision, and parties to agreements based on this part of ISO/IEC 11172 are encouraged to investigate the possibility of applying the most recent editions of the standards indicated below. Members of IEC and ISO maintain registers of currently valid International Standards.

ISO/IEC 11172-1:1993 *Information technology - Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s - Part 1: Systems.*

ISO/IEC 11172-3:1993 *Information technology - Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s - Part 3 Audio.*

CCIR Recommendation 601-2 *Encoding parameters of digital television for studios.*

CCIR Report 624-4 *Characteristics of systems for monochrome and colour television.*

CCIR Recommendation 648 *Recording of audio signals.*

CCIR Report 955-2 *Sound broadcasting by satellite for portable and mobile receivers, including Annex IV Summary description of Advanced Digital System II.*

CCITT Recommendation J.17 *Pre-emphasis used on Sound-Programme Circuits.*

IEEE Draft Standard P1180/D2 1990 *Specification for the implementation of 8x 8 inverse discrete cosine transform*".

IEC publication 908:1987 *CD Digital Audio System*.

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Section 2: Technical elements

2.1 Definitions

For the purposes of ISO/IEC 11172, the following definitions apply. If specific to a part, this is noted in square brackets.

2.1.1 ac coefficient [video]: Any DCT coefficient for which the frequency in one or both dimensions is non-zero.

2.1.2 access unit [system]: In the case of compressed audio an access unit is an audio access unit. In the case of compressed video an access unit is the coded representation of a picture.

2.1.3 adaptive segmentation [audio]: A subdivision of the digital representation of an audio signal in variable segments of time.

2.1.4 adaptive bit allocation [audio]: The assignment of bits to subbands in a time and frequency varying fashion according to a psychoacoustic model.

2.1.5 adaptive noise allocation [audio]: The assignment of coding noise to frequency bands in a time and frequency varying fashion according to a psychoacoustic model.

2.1.6 alias [audio]: Mirrored signal component resulting from sub-Nyquist sampling.

2.1.7 analysis filterbank [audio]: Filterbank in the encoder that transforms a broadband PCM audio signal into a set of subsampled subband samples.

2.1.8 audio access unit [audio]: For Layers I and II an audio access unit is defined as the smallest part of the encoded bitstream which can be decoded by itself, where decoded means "fully reconstructed sound". For Layer III an audio access unit is part of the bitstream that is decodable with the use of previously acquired main information.

2.1.9 audio buffer [audio]: A buffer in the system target decoder for storage of compressed audio data.

2.1.10 audio sequence [audio]: A non-interrupted series of audio frames in which the following parameters are not changed:

- ID
- Layer
- Sampling Frequency
- For Layer I and II: Bitrate index

2.1.11 backward motion vector [video]: A motion vector that is used for motion compensation from a reference picture at a later time in display order.

2.1.12 Bark [audio]: Unit of critical band rate. The Bark scale is a non-linear mapping of the frequency scale over the audio range closely corresponding with the frequency selectivity of the human ear across the band.

2.1.13 bidirectionally predictive-coded picture; B-picture [video]: A picture that is coded using motion compensated prediction from a past and/or future reference picture.

2.1.14 bitrate: The rate at which the compressed bitstream is delivered from the storage medium to the input of a decoder.

2.1.15 block companding [audio]: Normalizing of the digital representation of an audio signal within a certain time period.

2.1.16 block [video]: An 8-row by 8-column orthogonal block of pels.

2.1.17 bound [audio]: The lowest subband in which intensity stereo coding is used.

2.1.18 byte aligned: A bit in a coded bitstream is byte-aligned if its position is a multiple of 8-bits from the first bit in the stream.

2.1.19 byte: Sequence of 8-bits.

2.1.20 channel: A digital medium that stores or transports an ISO/IEC 11172 stream.

2.1.21 channel [audio]: The left and right channels of a stereo signal

2.1.22 chrominance (component) [video]: A matrix, block or single pel representing one of the two colour difference signals related to the primary colours in the manner defined in CCIR Rec 601. The symbols used for the colour difference signals are Cr and Cb.

2.1.23 coded audio bitstream [audio]: A coded representation of an audio signal as specified in ISO/IEC 11172-3.

2.1.24 coded video bitstream [video]: A coded representation of a series of one or more pictures as specified in this part of ISO/IEC 11172.

2.1.25 coded order [video]: The order in which the pictures are stored and decoded. This order is not necessarily the same as the display order.

2.1.26 coded representation: A data element as represented in its encoded form.

2.1.27 coding parameters [video]: The set of user-definable parameters that characterize a coded video bitstream. Bitstreams are characterised by coding parameters. Decoders are characterised by the bitstreams that they are capable of decoding.

2.1.28 component [video]: A matrix, block or single pel from one of the three matrices (luminance and two chrominance) that make up a picture.

2.1.29 compression: Reduction in the number of bits used to represent an item of data.

2.1.30 constant bitrate coded video [video]: A compressed video bitstream with a constant average bitrate.

2.1.31 constant bitrate: Operation where the bitrate is constant from start to finish of the compressed bitstream.

2.1.32 constrained parameters [video]: The values of the set of coding parameters defined in 2.4.3.2.

2.1.33 constrained system parameter stream (CSPS) [system]: An ISO/IEC 11172 multiplexed stream for which the constraints defined in 2.4.6 of ISO/IEC 11172-1 apply.

2.1.34 CRC: Cyclic redundancy code.

2.1.35 critical band rate [audio]: Psychoacoustic function of frequency. At a given audible frequency it is proportional to the number of critical bands below that frequency. The units of the critical band rate scale are Barks.

2.1.36 critical band [audio]: Psychoacoustic measure in the spectral domain which corresponds to the frequency selectivity of the human ear. This selectivity is expressed in Bark.

2.1.37 data element: An item of data as represented before encoding and after decoding.

2.1.38 dc-coefficient [video]: The DCT coefficient for which the frequency is zero in both dimensions.

2.1.39 dc-coded picture; D-picture [video]: A picture that is coded using only information from itself. Of the DCT coefficients in the coded representation, only the dc-coefficients are present.

2.1.40 DCT coefficient: The amplitude of a specific cosine basis function.

2.1.41 decoded stream: The decoded reconstruction of a compressed bitstream.

2.1.42 decoder input buffer [video]: The first-in first-out (FIFO) buffer specified in the video buffering verifier.

2.1.43 decoder input rate [video]: The data rate specified in the video buffering verifier and encoded in the coded video bitstream.

2.1.44 decoder: An embodiment of a decoding process.

2.1.45 decoding (process): The process defined in ISO/IEC 11172 that reads an input coded bitstream and produces decoded pictures or audio samples.

2.1.46 decoding time-stamp; DTS [system]: A field that may be present in a packet header that indicates the time that an access unit is decoded in the system target decoder.

2.1.47 de-emphasis [audio]: Filtering applied to an audio signal after storage or transmission to undo a linear distortion due to emphasis.

2.1.48 dequantization [video]: The process of rescaling the quantized DCT coefficients after their representation in the bitstream has been decoded and before they are presented to the inverse DCT.

2.1.49 digital storage media; DSM: A digital storage or transmission device or system.

2.1.50 discrete cosine transform; DCT [video]: Either the forward discrete cosine transform or the inverse discrete cosine transform. The DCT is an invertible, discrete orthogonal transformation. The inverse DCT is defined in annex A.

2.1.51 display order [video]: The order in which the decoded pictures should be displayed. Normally this is the same order in which they were presented at the input of the encoder.

2.1.52 dual channel mode [audio]: A mode, where two audio channels with independent programme contents (e.g. bilingual) are encoded within one bitstream. The coding process is the same as for the stereo mode.

2.1.53 editing: The process by which one or more compressed bitstreams are manipulated to produce a new compressed bitstream. Conforming edited bitstreams must meet the requirements defined in this part of ISO/IEC 11172.

2.1.54 elementary stream [system]: A generic term for one of the coded video, coded audio or other coded bitstreams.

2.1.55 emphasis [audio]: Filtering applied to an audio signal before storage or transmission to improve the signal-to-noise ratio at high frequencies.

2.1.56 encoder: An embodiment of an encoding process.

2.1.57 encoding (process): A process, not specified in ISO/IEC 11172, that reads a stream of input pictures or audio samples and produces a valid coded bitstream as defined in ISO/IEC 11172.

2.1.58 entropy coding: Variable length lossless coding of the digital representation of a signal to reduce redundancy.

2.1.59 fast forward playback [video]: The process of displaying a sequence, or parts of a sequence, of pictures in display-order faster than real-time.

2.1.60 FFT: Fast Fourier Transformation. A fast algorithm for performing a discrete Fourier transform (an orthogonal transform).

2.1.61 filterbank [audio]: A set of band-pass filters covering the entire audio frequency range.

2.1.62 fixed segmentation [audio]: A subdivision of the digital representation of an audio signal into fixed segments of time.

2.1.63 forbidden: The term "forbidden" when used in the clauses defining the coded bitstream indicates that the value shall never be used. This is usually to avoid emulation of start codes.

2.1.64 forced updating [video]: The process by which macroblocks are intra-coded from time-to-time to ensure that mismatch errors between the inverse DCT processes in encoders and decoders cannot build up excessively.

2.1.65 forward motion vector [video]: A motion vector that is used for motion compensation from a reference picture at an earlier time in display order.

2.1.66 frame [audio]: A part of the audio signal that corresponds to audio PCM samples from an Audio Access Unit.

2.1.67 free format [audio]: Any bitrate other than the defined bitrates that is less than the maximum valid bitrate for each layer.

2.1.68 future reference picture [video]: The future reference picture is the reference picture that occurs at a later time than the current picture in display order.

2.1.69 granules [Layer II] [audio]: The set of 3 consecutive subband samples from all 32 subbands that are considered together before quantization. They correspond to 96 PCM samples.

2.1.70 granules [Layer III] [audio]: 576 frequency lines that carry their own side information.

2.1.71 group of pictures [video]: A series of one or more coded pictures intended to assist random access. The group of pictures is one of the layers in the coding syntax defined in this part of ISO/IEC 11172.

2.1.72 Hann window [audio]: A time function applied sample-by-sample to a block of audio samples before Fourier transformation.

2.1.73 Huffman coding: A specific method for entropy coding.

2.1.74 hybrid filterbank [audio]: A serial combination of subband filterbank and MDCT.

2.1.75 IMDCT [audio]: Inverse Modified Discrete Cosine Transform.

2.1.76 intensity stereo [audio]: A method of exploiting stereo irrelevance or redundancy in stereophonic audio programmes based on retaining at high frequencies only the energy envelope of the right and left channels.

2.1.77 interlace [video]: The property of conventional television pictures where alternating lines of the picture represent different instances in time.

2.1.78 intra coding [video]: Coding of a macroblock or picture that uses information only from that macroblock or picture.

2.1.79 intra-coded picture; I-picture [video]: A picture coded using information only from itself.

2.1.80 ISO/IEC 11172 (multiplexed) stream [system]: A bitstream composed of zero or more elementary streams combined in the manner defined in ISO/IEC 11172-1.

2.1.81 joint stereo coding [audio]: Any method that exploits stereophonic irrelevance or stereophonic redundancy.

2.1.82 joint stereo mode [audio]: A mode of the audio coding algorithm using joint stereo coding.

2.1.83 layer [audio]: One of the levels in the coding hierarchy of the audio system defined in ISO/IEC 11172-3.

2.1.84 layer [video and systems]: One of the levels in the data hierarchy of the video and system specifications defined in ISO/IEC 11172-1 and this part of ISO/IEC 11172.

2.1.85 luminance (component) [video]: A matrix, block or single pel representing a monochrome representation of the signal and related to the primary colours in the manner defined in CCIR Rec 601. The symbol used for luminance is Y.

2.1.86 macroblock [video]: The four 8 by 8 blocks of luminance data and the two corresponding 8 by 8 blocks of chrominance data coming from a 16 by 16 section of the luminance component of the picture. Macroblock is sometimes used to refer to the pel data and sometimes to the coded representation of the pel values and other data elements defined in the macroblock layer of the syntax defined in this part of ISO/IEC 11172. The usage is clear from the context.

2.1.87 mapping [audio]: Conversion of an audio signal from time to frequency domain by subband filtering and/or by MDCT.

2.1.88 masking [audio]: A property of the human auditory system by which an audio signal cannot be perceived in the presence of another audio signal.

2.1.89 masking threshold [audio]: A function in frequency and time below which an audio signal cannot be perceived by the human auditory system.

2.1.90 MDCT [audio]: Modified Discrete Cosine Transform.

2.1.91 motion compensation [video]: The use of motion vectors to improve the efficiency of the prediction of pel values. The prediction uses motion vectors to provide offsets into the past and/or future reference pictures containing previously decoded pel values that are used to form the prediction error signal.

2.1.92 motion estimation [video]: The process of estimating motion vectors during the encoding process.

2.1.93 motion vector [video]: A two-dimensional vector used for motion compensation that provides an offset from the coordinate position in the current picture to the coordinates in a reference picture.

2.1.94 MS stereo [audio]: A method of exploiting stereo irrelevance or redundancy in stereophonic audio programmes based on coding the sum and difference signal instead of the left and right channels.

2.1.95 non-intra coding [video]: Coding of a macroblock or picture that uses information both from itself and from macroblocks and pictures occurring at other times.

2.1.96 non-tonal component [audio]: A noise-like component of an audio signal.

2.1.97 Nyquist sampling: Sampling at or above twice the maximum bandwidth of a signal.

2.1.98 pack [system]: A pack consists of a pack header followed by one or more packets. It is a layer in the system coding syntax described in ISO/IEC 11172-1.

2.1.99 packet data [system]: Contiguous bytes of data from an elementary stream present in a packet.

2.1.100 packet header [system]: The data structure used to convey information about the elementary stream data contained in the packet data.

2.1.101 packet [system]: A packet consists of a header followed by a number of contiguous bytes from an elementary data stream. It is a layer in the system coding syntax described in ISO/IEC 11172-1.

2.1.102 padding [audio]: A method to adjust the average length in time of an audio frame to the duration of the corresponding PCM samples, by conditionally adding a slot to the audio frame.

2.1.103 past reference picture [video]: The past reference picture is the reference picture that occurs at an earlier time than the current picture in display order.

2.1.104 pel aspect ratio [video]: The ratio of the nominal vertical height of pel on the display to its nominal horizontal width.

2.1.105 pel [video]: Picture element.

2.1.106 picture period [video]: The reciprocal of the picture rate.

2.1.107 picture rate [video]: The nominal rate at which pictures should be output from the decoding process.

2.1.108 picture [video]: Source, coded or reconstructed image data. A source or reconstructed picture consists of three rectangular matrices of 8-bit numbers representing the luminance and two chrominance signals. The Picture layer is one of the layers in the coding syntax defined in this part of ISO/IEC 11172. Note that the term "picture" is always used in ISO/IEC 11172 in preference to the terms field or frame.

2.1.109 polyphase filterbank [audio]: A set of equal bandwidth filters with special phase interrelationships, allowing for an efficient implementation of the filterbank.

2.1.110 prediction [video]: The use of a predictor to provide an estimate of the pel value or data element currently being decoded.

2.1.111 predictive-coded picture; P-picture [video]: A picture that is coded using motion compensated prediction from the past reference picture.

2.1.112 prediction error [video]: The difference between the actual value of a pel or data element and its predictor.

2.1.113 predictor [video]: A linear combination of previously decoded pel values or data elements.

2.1.114 presentation time-stamp; PTS [system]: A field that may be present in a packet header that indicates the time that a presentation unit is presented in the system target decoder.

2.1.115 presentation unit; PU [system]: A decoded audio access unit or a decoded picture.

2.1.116 psychoacoustic model [audio]: A mathematical model of the masking behaviour of the human auditory system.

2.1.117 quantization matrix [video]: A set of sixty-four 8-bit values used by the dequantizer.

2.1.118 quantized DCT coefficients [video]: DCT coefficients before dequantization. A variable length coded representation of quantized DCT coefficients is stored as part of the compressed video bitstream.

2.1.119 quantizer scalefactor [video]: A data element represented in the bitstream and used by the decoding process to scale the dequantization.

2.1.120 random access: The process of beginning to read and decode the coded bitstream at an arbitrary point.

2.1.121 reference picture [video]: Reference pictures are the nearest adjacent I- or P-pictures to the current picture in display order.

2.1.122 reorder buffer [video]: A buffer in the system target decoder for storage of a reconstructed I-picture or a reconstructed P-picture.

2.1.123 requantization [audio]: Decoding of coded subband samples in order to recover the original quantized values.

2.1.124 reserved: The term "reserved" when used in the clauses defining the coded bitstream indicates that the value may be used in the future for ISO/IEC defined extensions.

2.1.125 reverse playback [video]: The process of displaying the picture sequence in the reverse of display order.

2.1.126 scalefactor band [audio]: A set of frequency lines in Layer III which are scaled by one scalefactor.

2.1.127 scalefactor index [audio]: A numerical code for a scalefactor.

2.1.128 scalefactor [audio]: Factor by which a set of values is scaled before quantization.

2.1.129 sequence header [video]: A block of data in the coded bitstream containing the coded representation of a number of data elements.

2.1.130 side information: Information in the bitstream necessary for controlling the decoder.

2.1.131 skipped macroblock [video]: A macroblock for which no data are stored.

2.1.132 slice [video]: A series of macroblocks. It is one of the layers of the coding syntax defined in this part of ISO/IEC 11172.

2.1.133 slot [audio]: A slot is an elementary part in the bitstream. In Layer I a slot equals four bytes, in Layers II and III one byte.

2.1.134 source stream: A single non-multiplexed stream of samples before compression coding.

2.1.135 spreading function [audio]: A function that describes the frequency spread of masking.

2.1.136 start codes [system and video]: 32-bit codes embedded in that coded bitstream that are unique. They are used for several purposes including identifying some of the layers in the coding syntax.

2.1.137 STD input buffer [system]: A first-in first-out buffer at the input of the system target decoder for storage of compressed data from elementary streams before decoding.

2.1.138 stereo mode [audio]: Mode, where two audio channels which form a stereo pair (left and right) are encoded within one bitstream. The coding process is the same as for the dual channel mode.

2.1.139 stuffing (bits); stuffing (bytes) : Code-words that may be inserted into the compressed bitstream that are discarded in the decoding process. Their purpose is to increase the bitrate of the stream.

2.1.140 subband [audio]: Subdivision of the audio frequency band.

2.1.141 subband filterbank [audio]: A set of band filters covering the entire audio frequency range. In ISO/IEC 11172-3 the subband filterbank is a polyphase filterbank.

2.1.142 subband samples [audio]: The subband filterbank within the audio encoder creates a filtered and subsampled representation of the input audio stream. The filtered samples are called subband samples. From 384 time-consecutive input audio samples, 12 time-consecutive subband samples are generated within each of the 32 subbands.

2.1.143 syncword [audio]: A 12-bit code embedded in the audio bitstream that identifies the start of a frame.

2.1.144 synthesis filterbank [audio]: Filterbank in the decoder that reconstructs a PCM audio signal from subband samples.

2.1.145 system header [system]: The system header is a data structure defined in ISO/IEC 11172-1 that carries information summarising the system characteristics of the ISO/IEC 11172 multiplexed stream.

2.1.146 system target decoder; STD [system]: A hypothetical reference model of a decoding process used to describe the semantics of an ISO/IEC 11172 multiplexed bitstream.

2.1.147 time-stamp [system]: A term that indicates the time of an event.

2.1.148 triplet [audio]: A set of 3 consecutive subband samples from one subband. A triplet from each of the 32 subbands forms a granule.

2.1.149 tonal component [audio]: A sinusoid-like component of an audio signal.

2.1.150 variable bitrate: Operation where the bitrate varies with time during the decoding of a compressed bitstream.

2.1.151 variable length coding; VLC: A reversible procedure for coding that assigns shorter code-words to frequent events and longer code-words to less frequent events.

2.1.152 video buffering verifier; VBV [video]: A hypothetical decoder that is conceptually connected to the output of the encoder. Its purpose is to provide a constraint on the variability of the data rate that an encoder or editing process may produce.

2.1.153 video sequence [video]: A series of one or more groups of pictures. It is one of the layers of the coding syntax defined in this part of ISO/IEC 11172.

2.1.154 zig-zag scanning order [video]: A specific sequential ordering of the DCT coefficients from (approximately) the lowest spatial frequency to the highest.

2.2 Symbols and abbreviations

The mathematical operators used to describe this International Standard are similar to those used in the C programming language. However, integer division with truncation and rounding are specifically defined. The bitwise operators are defined assuming two's-complement representation of integers. Numbering and counting loops generally begin from zero.

2.2.1 Arithmetic operators

+	Addition.
-	Subtraction (as a binary operator) or negation (as a unary operator).
++	Increment.
--	Decrement.
*	Multiplication.
^	Power.
/	Integer division with truncation of the result toward zero. For example, $7/4$ and $-7/-4$ are truncated to 1 and $-7/4$ and $7/-4$ are truncated to -1.
//	Integer division with rounding to the nearest integer. Half-integer values are rounded away from zero unless otherwise specified. For example $3//2$ is rounded to 2, and $-3//2$ is rounded to -2.
DIV	Integer division with truncation of the result towards $-\infty$.
	Absolute value. $\begin{aligned} x &= x && \text{when } x > 0 \\ x &= 0 && \text{when } x = 0 \\ x &= -x && \text{when } x < 0 \end{aligned}$
%	Modulus operator. Defined only for positive numbers.
Sign()	$\text{Sign}(x) = \begin{cases} 1 & x > 0 \\ 0 & x = 0 \\ -1 & x < 0 \end{cases}$
NINT ()	Nearest integer operator. Returns the nearest integer value to the real-valued argument. Half-integer values are rounded away from zero.
sin	Sine.
cos	Cosine.
exp	Exponential.
$\sqrt{}$	Square root.
\log_{10}	Logarithm to base ten.
\log_e	Logarithm to base e.
\log_2	Logarithm to base 2.

2.2.2 Logical operators

	Logical OR.
&&	Logical AND.

! Logical NOT.

2.2.3 Relational operators

> Greater than.

>= Greater than or equal to.

< Less than.

<= Less than or equal to.

= Equal to.

!= Not equal to.

max [...], the maximum value in the argument list.

min [...], the minimum value in the argument list.

2.2.4 Bitwise operators

A twos complement number representation is assumed where the bitwise operators are used.

& AND.

| OR.

>> Shift right with sign extension.

<< Shift left with zero fill.

2.2.5 Assignment

= Assignment operator.

2.2.6 Mnemonics

The following mnemonics are defined to describe the different data types used in the coded bit-stream.

bslbf	Bit string, left bit first, where "left" is the order in which bit strings are written in ISO/IEC 11172. Bit strings are written as a string of 1s and 0s within single quote marks, e.g. '1000 0001'. Blanks within a bit string are for ease of reading and have no significance.
ch	Channel. If ch has the value 0, the left channel of a stereo signal or the first of two independent signals is indicated. (Audio)
nch	Number of channels; equal to 1 for single_channel mode, 2 in other modes. (Audio)
gr	Granule of 3 * 32 subband samples in audio Layer II, 18 * 32 sub-band samples in audio Layer III. (Audio)
main_data	The main_data portion of the bitstream contains the scalefactors, Huffman encoded data, and ancillary information. (Audio)
main_data_beg	The location in the bitstream of the beginning of the main_data for the frame. The location is equal to the ending location of the previous frame's main_data plus one bit. It is calculated from the main_data_end value of the previous frame. (Audio)
part2_length	The number of main_data bits used for scalefactors. (Audio)

rpchof	Remainder polynomial coefficients, highest order first. (Audio)
sb	Subband. (Audio)
sblimit	The number of the lowest sub-band for which no bits are allocated. (Audio)
scfsi	Scalefactor selection information. (Audio)
switch_point_l	Number of scalefactor band (long block scalefactor band) from which point on window switching is used. (Audio)
switch_point_s	Number of scalefactor band (short block scalefactor band) from which point on window switching is used. (Audio)
uimsbf	Unsigned integer, most significant bit first.
vclcbf	Variable length code, left bit first, where "left" refers to the order in which the VLC codes are written.
window	Number of the actual time slot in case of block_type==2, $0 \leq \text{window} \leq 2$. (Audio)

The byte order of multi-byte words is most significant byte first.

2.2.7 Constants

π	3,14159265358...
e	2,71828182845...

2.3 Method of describing bitstream syntax

The bitstream retrieved by the decoder is described in 2.4.2. Each data item in the bitstream is in bold type. It is described by its name, its length in bits, and a mnemonic for its type and order of transmission.

The action caused by a decoded data element in a bitstream depends on the value of that data element and on data elements previously decoded. The decoding of the data elements and definition of the state variables used in their decoding are described in 2.4.3. The following constructs are used to express the conditions when data elements are present, and are in normal type:

Note this syntax uses the 'C'-code convention that a variable or expression evaluating to a non-zero value is equivalent to a condition that is true.

while (condition) { data_element ... }	If the condition is true, then the group of data elements occurs next in the data stream. This repeats until the condition is not true.
do { data_element ... } while (condition)	The data element always occurs at least once. The data element is repeated until the condition is not true.
if (condition) { data_element ... }	If the condition is true, then the first group of data elements occurs next in the data stream.
else { data_element ... }	If the condition is not true, then the second group of data elements occurs next in the data stream.

for (expr1; expr2; expr3) { expr1 is an expression specifying the initialization of the loop. Normally it specifies the initial state of the counter. expr2 is a condition specifying a test made before each iteration of the loop. The loop terminates when the condition is not true. expr3 is an expression that is performed at the end of each iteration of the loop, normally it increments a counter.

Note that the most common usage of this construct is as follows:

```
for ( i = 0; i < n; i++) { The group of data elements occurs n times. Conditional constructs
  data_element          within the group of data elements may depend on the value of the
  . . .                 loop control variable i, which is set to zero for the first occurrence,
}                       incremented to one for the second occurrence, and so forth.
```

As noted, the group of data elements may contain nested conditional constructs. For compactness, the { } may be omitted when only one data element follows.

data_element [] data_element [] is an array of data. The number of data elements is indicated by the context.

data_element [n] data_element [n] is the n+1th element of an array of data.

data_element [m][n] data_element [m][n] is the m+1,n+1 th element of a two-dimensional array of data.

data_element [l][m][n] data_element [l][m][n] is the l+1,m+1,n+1 th element of a three-dimensional array of data.

data_element [m..n] is the inclusive range of bits between bit m and bit n in the data_element.

While the syntax is expressed in procedural terms, it should not be assumed that 2.4.3 implements a satisfactory decoding procedure. In particular, it defines a correct and error-free input bitstream. Actual decoders must include a means to look for start codes in order to begin decoding correctly, and to identify errors, erasures or insertions while decoding. The methods to identify these situations, and the actions to be taken, are not standardized.

Definition of bytealigned function

The function bytealigned () returns 1 if the current position is on a byte boundary, that is the next bit in the bitstream is the first bit in a byte. Otherwise it returns 0.

Definition of nextbits function

The function nextbits () permits comparison of a bit string with the next bits to be decoded in the bitstream.

Definition of next_start_code function

The next_start_code function removes any zero bit and zero byte stuffing and locates the next start code.

Syntax	No. of bits	Mnemonic
next_start_code() {		
while (!bytealigned())		
zero_bit	1	"0"
while (nextbits() != '0000 0000 0000 0000 0000 0001')		
zero_byte	8	"00000000"
}		

This function checks whether the current position is bytealigned. If it is not, zero stuffing bits are present. After that any number of zero bytes may be present before the start-code. Therefore start-codes are always bytealigned and may be preceded by any number of zero stuffing bits.

2.4 Requirements

2.4.1 Coding structure and parameters

Video sequence

A coded video sequence commences with a sequence header and is followed by one or more groups of pictures and is ended by a sequence_end_code. Immediately before each of the groups of pictures there may be a sequence header. Within each sequence, pictures shall be decodable continuously.

In each of these repeated sequence headers all of the data elements with the permitted exception of those defining the quantization matrices (load_intra_quantizer_matrix, load_non_intra_quantizer_matrix and optionally intra_quantizer_matrix and non_intra_quantizer_matrix) shall have the same values as in the first sequence header. The quantization matrices may be redefined each time that a sequence header occurs in the bitstream. Thus the data elements load_intra_quantizer_matrix, load_non_intra_quantizer_matrix and optionally intra_quantizer_matrix and non_intra_quantizer_matrix may have any (non-forbidden) values.

Repeating the sequence header allows the data elements of the initial sequence header to be repeated in order that random access into the video sequence is possible. In addition the quantization matrices may be changed inside the video sequence as required.

Sequence header

A video sequence header commences with a sequence_header_code and is followed by a series of data elements.

Group of pictures

A **group of pictures** is a series of one or more coded pictures intended to assist random access into the sequence. In the stored bitstream, the first coded picture in a group of pictures is an I-Picture. The order of the pictures in the coded stream is the order in which the decoder processes them in normal playback. In particular, adjacent B-Pictures in the coded stream are in display order. The last coded picture, in display order, of a group of pictures is either an I-Picture or a P-Picture.

The following is an example of groups of pictures taken from the beginning of a video sequence. In this example the first group of pictures contains seven pictures and subsequent groups of pictures contain nine pictures. There are two B-pictures between successive P-pictures and also two B-pictures between successive I- and P-pictures. Picture '1I' is used to form a prediction for picture '4P'. Pictures '4P' and '1I' are both used to form predictions for pictures '2B' and '3B'. Therefore the order of pictures in the coded sequence shall be '1I', '4P', '2B', '3B'. However, the decoder should display them in the order '1I', '2B', '3B', '4P'.

At the encoder input,

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
I	B	B	P	B	B	P	B	B	I	B	B	P	B	B	P	B	B	I	B	B	P	B	B	P

At the encoder output, in the stored bitstream, and at the decoder input,

1	4	2	3	7	5	6	10	8	9	13	11	12	16	14	15	19	17	18	22	20	21	25	23	24
I	P	B	B	P	B	B	I	B	B	P	B	B	P	B	B	I	B	B	P	B	B	P	B	B

where the double vertical bars mark the group of pictures boundaries. Note that in this example, the first group of pictures is two pictures shorter than subsequent groups of pictures, since at the beginning of video coding there are no B-pictures preceding the first I-Picture. However, in general, in display order, there may be B-Pictures preceding the first I-Picture in the group of pictures, even for the first group of pictures to be decoded.

At the decoder output,

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

A group of pictures may be of any length. A group of pictures shall contain one or more I-Pictures. Applications requiring random access, fast-forward playback, or fast and normal reverse playback may use relatively short groups of pictures. Groups of pictures may also be started at scene cuts or other cases where motion compensation is ineffective.

The number of consecutive B-Pictures is variable. Neither B- nor P-Pictures need be present.

A video sequence of groups of pictures that is read by the decoder may be different from the one at the encoder output due to editing.

Picture

A source or reconstructed picture consists of three rectangular matrices of eight-bit numbers; a luminance matrix (Y), and two chrominance matrices (Cb and Cr). The Y-matrix shall have an even number of rows and columns, and the Cb and Cr matrices shall be one half the size of the Y-matrix in both horizontal and vertical dimensions.

The Y, Cb and Cr components are related to the primary (analogue) Red, Green and Blue Signals (E'_R , E'_G and E'_B) as described in CCIR Recommendation 601. These primary signals are gamma pre-corrected. The assumed value of gamma is not defined in this part of ISO/IEC 11172 but may typically be in the region approximately 2,2 to approximately 2,8. Applications which require accurate colour reproduction may choose to specify the value of gamma more accurately, but this is outside the scope of this part of ISO/IEC 11172.

The luminance and chrominance samples are positioned as shown in figure 5, where "x" marks the position of the luminance (Y) samples and "0" marks the position of the chrominance (Cb and Cr) samples:

x		x		x		x		x		x
	0				0				0	
x		x		x		x		x		x
---	---	---		---	---	---		---	---	---
x		x		x		x		x		x
	0				0				0	
x		x		x		x		x		x
---	---	---		---	---	---		---	---	---
x		x		x		x		x		x
	0				0				0	
x		x		x		x		x		x

Figure 5 -- The position of luminance and chrominance samples.

There are four types of coded picture that use different coding methods.

An **Intra-coded picture (I-picture)** is coded using information only from itself.

A **Predictive-coded picture (P-picture)** is a picture which is coded using motion compensated prediction from a past I-Picture or P-Picture.

A **Bidirectionally predictive-coded picture (B-picture)** is a picture which is coded using motion compensated prediction from a past and/or future I-Picture or P-Picture.

A **dc coded (D) picture** is coded using information only from itself. Of the DCT coefficients only the dc ones are present. The D-Pictures shall not be in a sequence containing any other picture types.

Slice

A **slice** is a series of an arbitrary number of macroblocks with the order of macroblocks starting from the upper-left of the picture and proceeding by raster-scan order from left to right and top to bottom. The first and last macroblocks of a slice shall not be skipped macroblocks (see 2.4.4.4). Every slice shall contain at least one macroblock. Slices shall not overlap and there shall be no gaps between slices. The position of slices may change from picture to picture. The first slice shall start with the first macroblock in the picture and the last slice shall end with the last macroblock in the picture.

Macroblock

A **macroblock** contains a 16-pel by 16-line section of luminance component and the spatially corresponding 8-pel by 8-line section of each chrominance component. A macroblock has 4 luminance blocks and 2 chrominance blocks. The term "macroblock" can refer to source or reconstructed data or to scaled, quantized coefficients. The order of blocks in a macroblock is top-left, top-right, bottom-left, bottom-right blocks for Y, followed by Cb and Cr. Figure 6 shows the arrangement of these blocks. A skipped macroblock is one for which no information is stored (see 2.4.4.4).

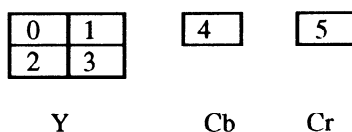


Figure 6 -- The arrangement of blocks in a macroblock.

Block

A **block** is an orthogonal 8-pel by 8-line section of a luminance or chrominance component.

The term "block" can refer either to source and reconstructed data or to the corresponding coded data elements.

Reserved, Forbidden and Marker bit

The terms "reserved" and "forbidden" are used in the description of some values of several fields in the coded bitstream.

The term "reserved" indicates that the value may be used in the future for ISO/IEC-defined extensions.

The term "forbidden" indicates a value that shall never be used (usually in order to avoid emulation of start codes).

The term "marker_bit" indicates a one bit field in which the value zero is forbidden. These marker bits are introduced at several points in the syntax to avoid start-code emulation.

2.4.2 Specification of the coded video bitstream syntax

2.4.2.1 Start codes

Start codes are reserved bit patterns that do not otherwise occur in the video stream. All start codes are bytealigned.

Name	Hexadecimal value
picture_start_code	00000100
slice_start_codes (including slice_vertical_positions)	00000101
	through
	000001AF
reserved	000001B0
reserved	000001B1
user_data_start_code	000001B2
sequence_header_code	000001B3
sequence_error_code	000001B4
extension_start_code	000001B5
reserved	000001B6
sequence_end_code	000001B7
group_start_code	000001B8
system start codes (see note)	000001B9
	through
	000001FF
NOTE - System start codes are defined in ISO/IEC 11172-1.	

The use of the start codes is defined in the following syntax description with the exception of the sequence_error_code. The sequence_error_code has been allocated for use by the digital storage media interface to indicate where uncorrectable errors have been detected.

2.4.2.2 Video sequence layer

Syntax	No. of bits	Mnemonic
<pre> video_sequence() { next_start_code() do { sequence_header() do { group_of_pictures() } while (nextbits() == group_start_code) } while (nextbits() == sequence_header_code) sequence_end_code } </pre>	32	bslbf

2.4.2.3 Sequence header

Syntax	No. of bits	Mnemonic
sequence_header() {		
sequence_header_code	32	bslbf
horizontal_size	12	uimsbf
vertical_size	12	uimsbf
pel_aspect_ratio	4	uimsbf
picture_rate	4	uimsbf
bit_rate	18	uimsbf
marker_bit	1	"1"
vbv_buffer_size	10	uimsbf
constrained_parameters_flag	1	
load_intra_quantizer_matrix	1	
if (load_intra_quantizer_matrix)		
intra_quantizer_matrix []	8*64	uimsbf
load_non_intra_quantizer_matrix	1	
if (load_non_intra_quantizer_matrix)		
non_intra_quantizer_matrix []	8*64	uimsbf
next_start_code()		
if (nextbits() == extension_start_code) {		
extension_start_code	32	bslbf
while (nextbits() != '0000 0000 0000 0000 0000 0001') {		
sequence_extension_data	8	
}		
next_start_code()		
}		
if (nextbits() == user_data_start_code) {		
user_data_start_code	32	bslbf
while (nextbits() != '0000 0000 0000 0000 0000 0001') {		
user_data	8	
}		
next_start_code()		
}		
}		

2.4.2.4 Group of pictures layer

Syntax	No. of bits	Mnemonic
group_of_pictures() {		
group_start_code	32	bslbf
time_code	25	
closed_gop	1	
broken_link	1	
next_start_code()		
if (nextbits() == extension_start_code) {		
extension_start_code	32	bslbf
while (nextbits() != '0000 0000 0000 0000 0000 0001') {		
group_extension_data	8	
}		
next_start_code()		
}		
if (nextbits() == user_data_start_code) {		
user_data_start_code	32	bslbf
while (nextbits() != '0000 0000 0000 0000 0000 0001') {		
user_data	8	
}		
next_start_code()		
}		
do {		
picture()		
} while (nextbits() == picture_start_code)		
}		

2.4.2.5 Picture layer

Syntax	No. of bits	Mnemonic
picture() {		
picture_start_code	32	bslbf
temporal_reference	10	uimsbf
picture_coding_type	3	uimsbf
vbv_delay	16	uimsbf
if ((picture_coding_type == 2) (picture_coding_type == 3)) {		
full_pel_forward_vector	1	
forward_f_code	3	uimsbf
}		
if (picture_coding_type == 3) {		
full_pel_backward_vector	1	
backward_f_code	3	uimsbf
}		
while (nextbits() == '1') {		
extra_bit_picture	1	"1"
extra_information_picture	8	
}		
extra_bit_picture	1	"0"
next_start_code()		
if (nextbits() == extension_start_code) {		
extension_start_code	32	bslbf
while (nextbits() != '0000 0000 0000 0000 0000 0001') {		
picture_extension_data	8	
}		
next_start_code()		
}		
if (nextbits() == user_data_start_code) {		
user_data_start_code	32	bslbf
while (nextbits() != '0000 0000 0000 0000 0000 0001') {		
user_data	8	
}		
next_start_code()		
}		
do {		
slice()		
} while (nextbits() == slice_start_code)		
}		

2.4.2.6 Slice layer

Syntax	No. of bits	Mnemonic
slice() {		
slice_start_code	32	bslbf
quantizer_scale	5	uimsbf
while (nextbits() == '1') {		
extra_bit_slice	1	"1"
extra_information_slice	8	
}		
extra_bit_slice	1	"0"
do {		
macroblock()		
} while (nextbits() != '000 0000 0000 0000 0000 0000')		
next_start_code()		
}		

2.4.2.7 Macroblock layer

Syntax	No. of bits	Mnemonic
macroblock () {		
while (nextbits() == '0000 0001 111')		
macroblock_stuffing	11	vlc1bf
while (nextbits() == '0000 0001 000')		
macroblock_escape	11	vlc1bf
macroblock_address_increment	1-11	vlc1bf
macroblock_type	1-6	vlc1bf
if (macroblock_quant)		
quantizer_scale	5	uimsbf
if (macroblock_motion_forward) {		
motion_horizontal_forward_code	1-11	vlc1bf
if ((forward_f != 1) &&		
(motion_horizontal_forward_code != 0))		
motion_horizontal_forward_r	1-6	uimsbf
motion_vertical_forward_code	1-11	vlc1bf
if ((forward_f != 1) &&		
(motion_vertical_forward_code != 0))		
motion_vertical_forward_r	1-6	uimsbf
}		
if (macroblock_motion_backward) {		
motion_horizontal_backward_code	1-11	vlc1bf
if ((backward_f != 1) &&		
(motion_horizontal_backward_code != 0))		
motion_horizontal_backward_r	1-6	uimsbf
motion_vertical_backward_code	1-11	vlc1bf
if ((backward_f != 1) &&		
(motion_vertical_backward_code != 0))		
motion_vertical_backward_r	1-6	uimsbf
}		
if (macroblock_pattern)		
coded_block_pattern	3-9	vlc1bf
for (i=0; i<6; i++)		
block(i)		
if (picture_coding_type == 4)		
end_of_macroblock	1	"1"
}		

2.4.2.8 Block layer

Syntax	No. of bits	Mnemonic
block(i) {		
if (pattern_code[i]) {		
if (macroblock_intra) {		
if (i<4) {		
dct_dc_size_luminance	2-7	vlclbf
if(dct_dc_size_luminance != 0)		
dct_dc_differential	1-8	uimsbf
}		
else {		
dct_dc_size_chrominance	2-8	vlclbf
if(dct_dc_size_chrominance !=0)		
dct_dc_differential	1-8	uimsbf
}		
}		
else {		
dct_coeff_first	2-28	vlclbf
}		
if (picture_coding_type != 4) {		
while (nextbits() != '10')		
dct_coeff_next	3-28	vlclbf
end_of_block	2	vlclbf
}		
}		
}		

2.4.3 Semantics for the video bitstream syntax

2.4.3.1 Video sequence layer

sequence_end_code -- The sequence_end_code is the bit string 000001B7 in hexadecimal. It terminates a video sequence.

2.4.3.2 Sequence header

sequence_header_code -- The sequence_header_code is the bit string 000001B3 in hexadecimal. It identifies the beginning of a sequence header.

horizontal_size -- The horizontal_size is the width of the displayable part of the luminance component in pels. The width of the encoded luminance component in macroblocks, mb_width, is $(horizontal_size+15)/16$. The displayable part of the picture is left-aligned in the encoded picture.

vertical_size -- The vertical_size is the height of the displayable part of the luminance component in pels. The height of the encoded luminance component in macroblocks, mb_height, is $(vertical_size+15)/16$. The displayable part of the picture is top-aligned in the encoded picture.

pel_aspect_ratio -- This is a four-bit integer defined in the following table.

pel_aspect_ratio	height/width	example
0000	forbidden	
0001	1,0000	VGA etc.
0010	0,6735	
0011	0,7031	16:9, 625line
0100	0,7615	
0101	0,8055	
0110	0,8437	16:9, 525line
0111	0,8935	
1000	0,9157	CCIR601, 625line
1001	0,9815	
1010	1,0255	
1011	1,0695	
1100	1,0950	CCIR601, 525line
1101	1,1575	
1110	1,2015	
1111	reserved	

picture_rate -- This is a four-bit integer defined in the following table.

picture_rate	pictures per second
0000	forbidden
0001	23,976
0010	24
0011	25
0100	29,97
0101	30
0110	50
0111	59,94
1000	60
...	reserved
1111	reserved

Applications and encoders should take into account the fact that 23,976, 29,97 and 59,94 are not exact representations of the nominal picture rate. The exact values are found from 24 000/1 001, 30 000/1 001, and 60 000/1 001 and can be derived from CCIR Report 624-4.

bit_rate -- This is an integer specifying the bitrate of the bitstream measured in units of 400 bits/s, rounded upwards. The value zero is forbidden. The value 3FFFF (11 1111 1111 1111 1111) identifies variable bit rate operation.

marker_bit -- This is one bit that shall be set to "1".

vbv_buffer_size -- This is a 10-bit integer defining the size of the VBV (Video Buffering Verifier, see annex C) buffer needed to decode the sequence. It is defined as:

$$B = 16 * 1024 * vbv_buffer_size$$

where B is the minimum VBV buffer size in bits required to decode the sequence (see annex C).

constrained_parameters_flag -- This is a one-bit flag which may be set to "1" if the following data elements meet the following constraints:

horizontal_size <= 768 pels,
vertical_size <= 576 pels,
((horizontal_size+15)/16)*((vertical_size+15)/16) <= 396,
((horizontal_size+15)/16)*((vertical_size+15)/16)*picture_rate <= 396*25,
picture_rate <= 30 pictures/s.
forward_f_code <= 4 (see 2.4.3.4)
backward_f_code <= 4 (see 2.4.3.4)

If the constrained_parameters_flag is set, then the vbv_buffer_size field shall indicate a VBV buffer size less than or equal to 327 680 bits (20*1024*16; i.e. 40 kbytes).

If the constrained_parameters_flag is set, then the bit_rate field shall indicate a coded data rate less than or equal to 1 856 000 bits/s.

load_intra_quantizer_matrix -- This is a one-bit flag which is set to "1" if an intra_quantizer_matrix follows. If it is set to "0" then the default values defined below in raster-scan order, are used until the next occurrence of the sequence header.

8	16	19	22	26	27	29	34
16	16	22	24	27	29	34	37
19	22	26	27	29	34	34	38
22	22	26	27	29	34	37	40
22	26	27	29	32	35	40	48
26	27	29	32	35	40	48	58
26	27	29	34	38	46	56	69
27	29	35	38	46	56	69	83

intra_quantizer_matrix -- This is a list of sixty-four 8-bit unsigned integers. The new values, stored in the zigzag scanning order shown in 2.4.4.1, replace the default values shown above. The value zero is forbidden. The value for intra_quant[0][0] shall always be 8. The new values shall be in effect until the next occurrence of a sequence header.

load_non_intra_quantizer_matrix -- This is a one-bit flag which is set to "1" if a non_intra_quantizer_matrix follows. If it is set to "0" then the default values defined below are used until the next occurrence of the sequence header.

16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16

non_intra_quantizer_matrix -- This is a list of sixty-four 8-bit unsigned integers. The new values, stored in the zigzag scanning order shown in 2.4.4.1, replace the default values shown above. The value zero is forbidden. The new values shall be in effect until the next occurrence of a sequence header.

extension_start_code -- The extension_start_code is the bit string 000001B5 in hexadecimal. It identifies the beginning of extension data. The extension data continue until receipt of another start code. It is a requirement to parse extension data correctly.

sequence_extension_data -- Reserved.

user_data_start_code -- The user_data_start_code is the bit string 000001B2 in hexadecimal. It identifies the beginning of user data. The user data continues until receipt of another start code.

user_data -- The user_data is defined by the users for their specific applications. The user data shall not contain a string of 23 or more zero bits.

2.4.3.3 Group of pictures layer

group_start_code -- The group_start_code is the bit string 000001B8 in hexadecimal. It identifies the beginning of a group of pictures.

time_code -- This is a 25-bit field containing the following: drop_frame_flag, time_code_hours, time_code_minutes, marker_bit, time_code_seconds and time_code_pictures. The fields correspond to the fields defined in the IEC standard (Publication 461) for "time and control codes for video tape recorders" (see annex E). The code refers to the first picture in the group of pictures that has a temporal reference of zero. The drop_frame_flag can be set to either "0" or "1". It may be set to "1" only if the picture rate is 29,97Hz. If it is "0" then pictures are counted assuming rounding to the nearest integral number of pictures per second, for example 29,97 Hz would be rounded to and counted as 30 Hz. If it is "1" then picture numbers 0 and 1 at the start of each minute, except minutes 0, 10, 20, 30, 40, 50 are omitted from the count.

time_code	range of value	bits	
drop_frame_flag		1	
time_code_hours	0 - 23	5	uimbsf
time_code_minutes	0 - 59	6	uimbsf
marker_bit	1	1	"1"
time_code_seconds	0 - 59	6	uimbsf
time_code_pictures	0 - 59	6	uimbsf

closed_gop -- This is a one-bit flag which may be set to "1" if the group of pictures has been encoded without motion vectors pointing to the previous group of pictures.

This bit is provided for use during any editing which occurs after encoding. If the previous group of pictures is removed by editing, **broken_link** may be set to "1" so that a decoder may avoid displaying the B-Pictures immediately following the first I-Picture of the group of pictures. However if the **closed_gop** bit indicates that there are no prediction references to the previous group of pictures then the editor may choose not to set the **broken_link** bit as these B-Pictures can be correctly decoded in this case.

broken_link -- This is a one-bit flag which shall be set to "0" during encoding. It is set to "1" to indicate that the B-Pictures immediately following the first I-Picture of a group of pictures cannot be correctly decoded because the other I-Picture or P-Picture which is used for prediction is not available (because of the action of editing).

A decoder may use this flag to avoid displaying pictures that cannot be correctly decoded.

extension_start_code -- See 2.4.3.2.

group_extension_data -- Reserved.

user_data_start_code -- See 2.4.3.2.

user_data -- See 2.4.3.2.

2.4.3.4 Picture layer

picture_start_code -- The picture_start_code is a string of 32-bits having the value 00000100 in hexadecimal.

temporal_reference -- The temporal_reference is a 10-bit unsigned integer associated with each input picture. It is incremented by one, modulo 1024, for each input picture. For the earliest picture (in display order) in each group of pictures, the temporal_reference is reset to zero.

The temporal_reference is assigned (in sequence) to the pictures in display order, no temporal_reference shall be omitted from the sequence.

picture_coding_type -- The picture_coding_type identifies whether a picture is an intra-coded picture(I), predictive-coded picture(P), bidirectionally predictive-coded picture(B), or intra-coded with only dc coefficients picture(D) according to the following table. D-pictures shall never be included in the same video sequence as the other picture coding types.

picture_coding_type	coding method
000	forbidden
001	intra-coded (I)
010	predictive-coded (P)
011	bidirectionally-predictive-coded (B)
100	dc intra-coded (D)
101	reserved
...	...
111	reserved

vbv_delay -- The vbv_delay is a 16-bit unsigned integer. For constant bitrate operation, the vbv_delay is used to set the initial occupancy of the decoder's buffer at the start of decoding the picture so that the decoder's buffer does not overflow or underflow. The vbv_delay measures the time needed to fill the VBV buffer from an initially empty state at the target bit rate, R, to the correct level immediately before the current picture is removed from the buffer.

The value of vbv_delay is the number of periods of the 90kHz system clock that the VBV should wait after receiving the final byte of the picture start code. It may be calculated from the state of the VBV as follows:

$$vbv_delay_n = 90\,000 * B_n^* / R$$

where:

$$n > 0$$

B_n^* = VBV occupancy, measured in bits, immediately before removing picture n from the buffer but after removing any group of picture layer data, sequence header data and the picture_start_code that immediately precedes the data elements of picture n.

R = bitrate measured in bits/s. The full precision of the bitrate rather than the rounded value encoded by the bit_rate field in the sequence header shall be used by the encoder in the VBV model.

For non-constant bitrate operation vbv_delay shall have the value FFFF in hexadecimal.

full_pel_forward_vector -- If set to "1", then the motion vector values decoded represent integer pel offsets (rather than half-pel units) as reflected in the equations of 2.4.4.2.

forward_f_code -- An unsigned integer taking values 1 through 7. The value zero is forbidden. The variables forward_r_size and forward_f used in the process of decoding the forward motion vectors are derived from forward_f_code as described in 2.4.4.2.

full_pel_backward_vector -- If set to "1", then the motion vector values decoded represent integer pel offsets (rather than half pel units) as reflected in the equations of 2.4.4.3.

backward_f_code -- An unsigned integer taking values 1 through 7. The value zero is forbidden. The variables **backward_r_size** and **backward_f** used in the process of decoding the backward motion vectors are derived from **backward_f_code** as described in 2.4.4.3.

extra_bit_picture -- A bit indicates the presence of the following extra information. If **extra_bit_picture** is set to "1", **extra_information_picture** will follow it. If it is set to "0", there are no data following it.

extra_information_picture -- Reserved.

extension_start_code -- See 2.4.3.2.

picture_extension_data -- Reserved.

user_data_start_code -- See 2.4.3.2.

user_data -- See 2.4.3.2.

2.4.3.5 Slice layer

slice_start_code -- The **slice_start_code** is a string of 32-bits. The first 24-bits have the value 000001 in hexadecimal and the last 8-bits are the **slice_vertical_position** having a value in the range 01 through AF hexadecimal inclusive.

slice_vertical_position -- This is given by the last eight bits of the **slice_start_code**. It is an unsigned integer giving the vertical position in macroblock units of the first macroblock in the slice. The **slice_vertical_position** of the first row of macroblocks is one. Some slices may have the same **slice_vertical_position**, since slices may start and finish anywhere. Note that the **slice_vertical_position** is constrained by 2.4.1 to define non-overlapping slices with no gaps between them. The maximum value of **slice_vertical_position** is 175.

quantizer_scale -- An unsigned integer in the range 1 to 31 used to scale the reconstruction level of the retrieved DCT coefficient levels. The decoder shall use this value until another **quantizer_scale** is encountered either at the slice layer or the macroblock layer. The value zero is forbidden.

extra_bit_slice -- A bit indicates the presence of the following extra information. If **extra_bit_slice** is set to "1", **extra_information_slice** will follow it. If it is set to "0", there are no data following it.

extra_information_slice -- Reserved.

2.4.3.6 Macroblock layer

macroblock_stuffing -- This is a fixed bit string "0000 0001 111" which can be inserted by the encoder to increase the bit rate to that required of the storage or transmission medium. It is discarded by the decoder.

macroblock_escape -- The **macroblock_escape** is a fixed bit-string "0000 0001 000" which is used when the difference between **macroblock_address** and **previous_macroblock_address** is greater than 33. It causes the value of **macroblock_address_increment** to be 33 greater than the value that will be decoded by subsequent **macroblock_escapes** and the **macroblock_address_increment** codewords.

For example, if there are two **macroblock_escape** codewords preceding the **macroblock_address_increment**, then 66 is added to the value indicated by **macroblock_address_increment**.

macroblock_address_increment -- This is a variable length coded integer coded as per table B.1 which indicates the difference between **macroblock_address** and **previous_macroblock_address**. The maximum value of **macroblock_address_increment** is 33. Values greater than this can be encoded using the **macroblock_escape** codeword.

The **macroblock_address** is a variable defining the absolute position of the current macroblock. The **macroblock_address** of the top-left macroblock is zero.

The `previous_macroblock_address` is a variable defining the absolute position of the last non-skipped macroblock (see 2.4.4.4 for the definition of skipped macroblocks) except at the start of a slice. At the start of a slice, `previous_macroblock_address` is reset as follows:

$$\text{previous_macroblock_address} = (\text{slice_vertical_position} - 1) * \text{mb_width} - 1;$$

The spatial position in macroblock units of a macroblock in the picture (`mb_row`, `mb_column`) can be computed from the `macroblock_address` as follows:

$$\begin{aligned}\text{mb_row} &= \text{macroblock_address} / \text{mb_width} \\ \text{mb_column} &= \text{macroblock_address} \% \text{mb_width}\end{aligned}$$

where `mb_width` is the number of macroblocks in one row of the picture.

NOTE - The `slice_vertical_position` differs from `mb_row` by one.

macroblock_type -- Variable length coded indicator which indicates the method of coding and content of the macroblock according to the tables B.2a through B.2d.

macroblock_quant -- Derived from `macroblock_type`.

macroblock_motion_forward -- Derived from `macroblock_type`.

macroblock_motion_backward -- Derived from `macroblock_type`.

macroblock_pattern -- Derived from `macroblock_type`.

macroblock_intra -- Derived from `macroblock_type`.

quantizer_scale -- An unsigned integer in the range 1 to 31 used to scale the reconstruction level of the retrieved DCT coefficient levels. The value zero is forbidden. The decoder shall use this value until another `quantizer_scale` is encountered either at the slice layer or the macroblock layer. The presence of `quantizer_scale` is determined from `macroblock_type`.

motion_horizontal_forward_code -- `motion_horizontal_forward_code` is decoded according to table B.4. The decoded value is required (along with `forward_f` - see 2.4.4.2) to decide whether or not `motion_horizontal_forward_r` appears in the bitstream.

motion_horizontal_forward_r -- An unsigned integer (of `forward_r_size` bits - see 2.4.4.2) used in the process of decoding forward motion vectors as described in 2.4.4.2.

motion_vertical_forward_code -- `motion_vertical_forward_code` is decoded according to table B.4. The decoded value is required (along with `forward_f` - see 2.4.4.2) to decide whether or not `motion_vertical_forward_r` appears in the bitstream.

motion_vertical_forward_r -- An unsigned integer (of `forward_r_size` bits - see 2.4.4.2) used in the process of decoding forward motion vectors as described in 2.4.4.2.

motion_horizontal_backward_code -- `motion_horizontal_backward_code` is decoded according to table B.4. The decoded value is required (along with `backward_f` - see 2.4.4.2) to decide whether or not `motion_horizontal_backward_r` appears in the bitstream.

motion_horizontal_backward_r -- An unsigned integer (of `backward_r_size` bits - see 2.4.4.2) used in the process of decoding backward motion vectors as described in 2.4.4.2.

motion_vertical_backward_code -- `motion_vertical_backward_code` is decoded according to table B.4. The decoded value is required (along with `backward_f`) to decide whether or not `motion_vertical_backward_r` appears in the bitstream.

motion_vertical_backward_r -- An unsigned integer (of `backward_r_size` bits) used in the process of decoding backward motion vectors as described in 2.4.4.3.

coded_block_pattern -- coded_block_pattern is a variable length code that is used to derive the variable cbp according to table B.3. If macroblock_intra is zero, cbp=0. Then the pattern_code[i] for i=0 to 5 is derived from cbp using the following:

```

pattern_code[i] = 0;
if ( cbp & (1<<(5-i)) ) pattern_code[i] = 1;
if ( macroblock_intra ) pattern_code[i] = 1 ;

```

pattern_code[0] -- If 1, then the upper left luminance block is to be received in this macroblock.

pattern_code[1] -- If 1, then the upper right luminance block is to be received in this macroblock.

pattern_code[2] -- If 1, then the lower left luminance block is to be received in this macroblock.

pattern_code[3] -- If 1, then the lower right luminance block is to be received in this macroblock..

pattern_code[4] -- If 1, then the chrominance block Cb is to be received in this macroblock.

pattern_code[5] -- If 1, then the chrominance block Cr is to be received in this macroblock.

end_of_macroblock -- This is a bit which is set to "1" and exists only in D-Pictures.

2.4.3.7 Block layer

dct_dc_size_luminance -- The number of bits in the following dct_dc_differential code, dc_size_luminance, is derived according to the VLC table B.5a. Note that this data element is used in intra coded blocks.

dct_dc_size_chrominance -- The number of bits in the following dct_dc_differential code, dc_size_chrominance, is derived according to the VLC table B.5b. Note that this data element is used in intra coded blocks.

dct_dc_differential -- A variable length unsigned integer. If dc_size_luminance or dc_size_chrominance (as appropriate) is zero, then dct_dc_differential is not present in the bitstream. dct_zz[] is the array of quantized DCT coefficients in zig-zag scanning order. dct_zz[i] for i=0..63 shall be set to zero initially. If dc_size_luminance or dc_size_chrominance (as appropriate) is greater than zero, then dct_zz[0] is computed as follows from dct_dc_differential:

For luminance blocks:

```

if ( dct_dc_differential & ( 1 << (dc_size_luminance-1)) ) dct_zz[0] = dct_dc_differential ;
else dct_zz[0] = ( (-1) << (dc_size_luminance) ) | (dct_dc_differential+1) ;

```

For chrominance blocks:

```

if ( dct_dc_differential & ( 1 << (dc_size_chrominance-1)) ) dct_zz[0] = dct_dc_differential ;
else dct_zz[0] = ( (-1) << (dc_size_chrominance) ) | (dct_dc_differential+1) ;

```

Note that this data element is used in intra coded blocks.

example for dc_size_luminance = 3	
dct_dc_differential	dct_zz[0]
000	-7
001	-6
010	-5
011	-4
100	4
101	5
110	6
111	7

dct_coeff_first -- A variable length code according to tables B.5c through B.5f for the first coefficient. The variables run and level are derived according to these tables. The zigzag-scanned quantized DCT coefficient list is updated as follows.

```
i = run ;  
if ( s == 0 ) dct_zz[i] = level ;  
if ( s == 1 ) dct_zz[i] = - level ;
```

The terms `dct_coeff_first` and `dct_coeff_next` are run-length encoded and `dct_zz[i]`, $i \geq 0$ shall be set to zero initially. A variable length code according to tables B.5c through B.5f is used to represent the run-length and level of the DCT coefficients. Note that this data element is used in non-intra coded blocks.

dct_coeff_next -- A variable length code according to tables B.5c through B.5f for coefficients following the first retrieved. The variables `run` and `level` are derived according to these tables. The zigzag-scanned quantized DCT coefficient list is updated as follows.

```
i = i + run + 1 ;  
if ( s == 0 ) dct_zz[i] = level ;  
if ( s == 1 ) dct_zz[i] = - level ;
```

If `macroblock_intra == 1` then the term `i` shall be set to zero before the first `dct_coeff_next` of the block. The decoding of `dct_coeff_next` shall not cause `i` to exceed 63.

end_of_block -- This symbol is always used to indicate that no additional non-zero coefficients are present. It is used even if `dct_zz[63]` is non-zero. Its value is the bit-string "10" as defined in table B.5c.

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2.4.4 The video decoding process

Compliance requirements for decoders are contained in ISO/IEC 11172-4.

2.4.4.1 Intra-coded macroblocks

In I-pictures all macroblocks are intra-coded and stored. In P-pictures and B-pictures, some macroblocks may be intra-coded as identified by `macroblock_type`. Thus, `macroblock_intra` identifies the intra-coded macroblocks.

The variables `mb_row` and `mb_column` locate the macroblock in the picture. They are defined in 2.4.3.6. The definitions of `dct_dc_differential`, and `dct_coeff_next` also have defined the zigzag-scanned quantized DCT coefficient list, `dct_zz[]`. Each `dct_zz[]` is located in the macroblock as defined by `pattern_code[]`.

Define `dct_recon[m][n]` to be the matrix of reconstructed DCT coefficients of the block, where the first index identifies the row and the second the column of the matrix. Define `dct_dc_y_past`, `dct_dc_cb_past` and `dct_dc_cr_past` to be the `dct_recon[0][0]` of the most recently decoded intra-coded Y, Cb and Cr blocks respectively. The predictors `dct_dc_y_past`, `dct_dc_cb_past` and `dct_dc_cr_past` shall all be reset at the start of a slice and at non-intra-coded macroblocks (including skipped macroblocks) to the value 1 024 (128*8).

Define `intra_quant[m][n]` to be the intra quantizer matrix that is specified in the sequence header.

Note that `intra_quant[0][0]` is used in the dequantizer calculations for simplicity of description, but the result is overwritten by the subsequent calculation for the dc coefficient.

Define `scan[m][n]` to be the matrix defining the zigzag scanning sequence as follows:

0	1	5	6	14	15	27	28
2	4	7	13	16	26	29	42
3	8	12	17	25	30	41	43
9	11	18	24	31	40	44	53
10	19	23	32	39	45	52	54
20	22	33	38	46	51	55	60
21	34	37	47	50	56	59	61
35	36	48	49	57	58	62	63

Where `n` is the horizontal index and `m` is the vertical index.

Define `past_intra_address` as the `macroblock_address` of the most recently retrieved intra-coded macroblock within the slice. It shall be reset to -2 at the beginning of each slice.

Then `dct_recon[m][n]` shall be computed by any means equivalent to the following procedure for the first luminance block:

```

for (m=0; m<8; m++) {
    for (n=0; n<8; n++) {
        i = scan[m][n];
        dct_recon[m][n] = ( 2 * dct_zz[i] * quantizer_scale * intra_quant[m][n] ) / 16 ;
        if ( ( dct_recon[m][n] & 1 ) == 0 )
            dct_recon[m][n] = dct_recon[m][n] - Sign(dct_recon[m][n]);
        if (dct_recon[m][n] > 2 047) dct_recon[m][n] = 2 047 ;
        if (dct_recon[m][n] < -2 048) dct_recon[m][n] = -2 048 ;
    }
}
dct_recon[0][0] = dct_zz[0] * 8 ;
if ( ( macroblock_address - past_intra_address > 1 ) )
    dct_recon[0][0] = (128 * 8) + dct_recon[0][0] ;
else
    dct_recon[0][0] = dct_dc_y_past + dct_recon[0][0] ;
dct_dc_y_past = dct_recon[0][0] ;

```

Note that this process disallows even valued numbers. This has been found to prevent accumulation of mismatch errors.

For the subsequent luminance blocks in the macroblock, in the order of the list defined by the array `pattern_code[]`:

```

for (m=0; m<8; m++) {
    for (n=0; n<8; n++) {
        i = scan[m][n];
        dct_recon[m][n] = ( 2 * dct_zz[i] * quantizer_scale * intra_quant[m][n] ) / 16;
        if ( ( dct_recon[m][n] & 1 ) == 0 )
            dct_recon[m][n] = dct_recon[m][n] - Sign(dct_recon[m][n]);
        if (dct_recon[m][n] > 2 047) dct_recon[m][n] = 2 047;
        if (dct_recon[m][n] < -2 048) dct_recon[m][n] = -2 048;
    }
}
dct_recon[0][0] = dct_dc_y_past + (dct_zz[0] * 8);
dct_dc_y_past = dct_recon[0][0];

```

For the chrominance Cb block,:

```

for (m=0; m<8; m++) {
    for (n=0; n<8; n++) {
        i = scan[m][n];
        dct_recon[m][n] = ( 2 * dct_zz[i] * quantizer_scale * intra_quant[m][n] ) / 16;
        if ( ( dct_recon[m][n] & 1 ) == 0 )
            dct_recon[m][n] = dct_recon[m][n] - Sign(dct_recon[m][n]);
        if (dct_recon[m][n] > 2 047) dct_recon[m][n] = 2 047;
        if (dct_recon[m][n] < -2 048) dct_recon[m][n] = -2 048;
    }
}
dct_recon[0][0] = dct_zz[0] * 8;
if ( ( macroblock_address - past_intra_address ) > 1 )
    dct_recon[0][0] = (128 * 8) + dct_recon[0][0];
else
    dct_recon[0][0] = dct_dc_cb_past + dct_recon[0][0];
dct_dc_cb_past = dct_recon[0][0];

```

For the chrominance Cr block, :

```

for (m=0; m<8; m++) {
    for (n=0; n<8; n++) {
        i = scan[m][n];
        dct_recon[m][n] = ( 2 * dct_zz[i] * quantizer_scale * intra_quant[m][n] ) / 16;
        if ( ( dct_recon[m][n] & 1 ) == 0 )
            dct_recon[m][n] = dct_recon[m][n] - Sign(dct_recon[m][n]);
        if (dct_recon[m][n] > 2 047) dct_recon[m][n] = 2 047;
        if (dct_recon[m][n] < -2 048) dct_recon[m][n] = -2 048;
    }
}
dct_recon[0][0] = dct_zz[0] * 8;
if ( ( macroblock_address - past_intra_address ) > 1 )
    dct_recon[0][0] = (128 * 8) + dct_recon[0][0];
else
    dct_recon[0][0] = dct_dc_cr_past + dct_recon[0][0];
dct_dc_cr_past = dct_recon[0][0];

```

After all the blocks in the macroblock are processed:

```
past_intra_address = macroblock_address;
```

Values in the coded data elements leading to $\text{dct_recon}[0][0] < 0$ or $\text{dct_recon}[0][0] > 2\,047$ are not permitted.

Once the DCT coefficients are reconstructed, the inverse DCT transform defined in annex A shall be applied to obtain the inverse transformed pel values in the range [-256, 255]. These pel values shall be limited to

the range [0, 255] and placed in the luminance and chrominance matrices in the positions defined by mb_row, mb_column, and the list defined by the array pattern_code[].

2.4.4.2 Predictive-coded macroblocks in P-pictures

Predictive-coded macroblocks in P-Pictures are decoded in two steps.

First, the value of the forward motion vector for the macroblock is reconstructed and a prediction macroblock is formed, as detailed below.

Second, the DCT coefficient information stored for some or all of the blocks is decoded, dequantized, inverse DCT transformed, and added to the prediction macroblock.

Let **recon_right_for** and **recon_down_for** be the reconstructed horizontal and vertical components of the motion vector for the current macroblock, and **recon_right_for_prev** and **recon_down_for_prev** be the reconstructed motion vector for the previous predictive-coded macroblock. If the current macroblock is the first macroblock in the slice, or if the last macroblock that was decoded contained no motion vector information (either because it was skipped or macroblock_motion_forward was zero), then **recon_right_for_prev** and **recon_down_for_prev** shall be set to zero.

If no forward motion vector data exists for the current macroblock (either because it was skipped or macroblock_motion_forward == 0), the motion vectors shall be set to zero.

If forward motion vector data exists for the current macroblock, then any means equivalent to the following procedure shall be used to reconstruct the motion vector horizontal and vertical components.

forward_r_size and forward_f are derived from forward_f_code as follows:

```

forward_r_size = forward_f_code - 1
forward_f = 1 << forward_r_size

if ( (forward_f == 1) || (motion_horizontal_forward_code == 0) ) {
    complement_horizontal_forward_r = 0;
} else {
    complement_horizontal_forward_r = forward_f - 1 - motion_horizontal_forward_r;
}
if ( (forward_f == 1) || (motion_vertical_forward_code == 0) ) {
    complement_vertical_forward_r = 0;
} else {
    complement_vertical_forward_r = forward_f - 1 - motion_vertical_forward_r;
}

right_little = motion_horizontal_forward_code * forward_f;
if (right_little == 0) {
    right_big = 0;
} else {
    if (right_little > 0) {
        right_little = right_little - complement_horizontal_forward_r;
        right_big = right_little - (32 * forward_f);
    } else {
        right_little = right_little + complement_horizontal_forward_r;
        right_big = right_little + (32 * forward_f);
    }
}

```

```

down_little = motion_vertical_forward_code * forward_f;
if (down_little == 0) {
    down_big = 0;
} else {
    if (down_little > 0) {
        down_little = down_little - complement_vertical_forward_r;
        down_big = down_little - (32 * forward_f);
    } else {
        down_little = down_little + complement_vertical_forward_r;
        down_big = down_little + (32 * forward_f);
    }
}

```

Values of forward_f, motion_horizontal_forward_code and if present, motion_horizontal_forward_r shall be such that right_little is not equal to forward_f * 16.

Values of forward_f, motion_vertical_forward_code and if present, motion_vertical_forward_r shall be such that down_little is not equal to forward_f * 16.

```

max = ( 16 * forward_f ) - 1 ;
min = ( -16 * forward_f ) ;

new_vector = recon_right_for_prev + right_little ;
if ( (new_vector <= max) && (new_vector >= min) )
    recon_right_for = recon_right_for_prev + right_little ;
else
    recon_right_for = recon_right_for_prev + right_big ;
recon_right_for_prev = recon_right_for ;

if ( full_pel_forward_vector ) recon_right_for = recon_right_for << 1 ;
new_vector = recon_down_for_prev + down_little ;
if ( (new_vector <= max) && (new_vector >= min) )
    recon_down_for = recon_down_for_prev + down_little ;
else
    recon_down_for = recon_down_for_prev + down_big ;
recon_down_for_prev = recon_down_for ;
if ( full_pel_forward_vector ) recon_down_for = recon_down_for << 1 ;

```

The motion vectors in whole pel units for the macroblock, right_for and down_for, and the half pel unit flags, right_half_for and down_half_for, are computed as follows:

for luminance	for chrominance
right_for = recon_right_for >> 1 ;	right_for = (recon_right_for / 2) >> 1 ;
down_for = recon_down_for >> 1 ;	down_for = (recon_down_for / 2) >> 1 ;
right_half_for = recon_right_for - (2*right_for) ;	right_half_for = recon_right_for/2 - (2*right_for) ;
down_half_for = recon_down_for - (2*down_for) ;	down_half_for = recon_down_for/2 - (2*down_for) ;

Motion vectors leading to references outside a reference picture's boundaries are not allowed.

A positive value of the reconstructed horizontal motion vector (right_for) indicates that the referenced area of the past reference picture is to the right of the macroblock in the coded picture.

A positive value of the reconstructed vertical motion vector (down_for) indicates that the referenced area of the past reference picture is below the macroblock in the coded picture.

Defining pel_past[][] as the pel values of the past picture referenced by the forward motion vector, and pel[][] as the predictors for the pel values of the block being decoded, then:

```

if ( (! right_half_for) && (! down_half_for) )
    pel[i][j] = pel_past[i+down_for][j+right_for] ;

```

```

if ( (! right_half_for) && down_half_for )
    pel[i][j] = ( pel_past[i+down_for][j+right_for] +
                  pel_past[i+down_for+1][j+right_for] ) // 2 ;

if ( right_half_for && (! down_half_for) )
    pel[i][j] = ( pel_past[i+down_for][j+right_for] +
                  pel_past[i+down_for][j+right_for+1] ) // 2 ;

if ( right_half_for && down_half_for )
    pel[i][j] = ( pel_past[i+down_for][j+right_for] + pel_past[i+down_for+1][j+right_for] +
                  pel_past[i+down_for][j+right_for+1] + pel_past[i+down_for+1][j+right_for+1] ) // 4 ;

```

Define non_intra_quant[m][n] to be the non-intra quantizer matrix that is specified in the sequence header.

The DCT coefficients for each block present in the macroblock shall be reconstructed by any means equivalent to the following procedure:

```

for ( m=0; m<8; m++ ) {
    for ( n=0; n<8; n++ ) {
        i = scan[m][n] ;
        dct_recon[m][n] = ( ( ( 2 * dct_zz[i] ) + Sign(dct_zz[i]) ) *
                             quantizer_scale * non_intra_quant[m][n] ) / 16 ;
        if ( ( dct_recon[m][n] & 1 ) == 0 )
            dct_recon[m][n] = dct_recon[m][n] - Sign(dct_recon[m][n]) ;
        if ( dct_recon[m][n] > 2047 ) dct_recon[m][n] = 2047 ;
        if ( dct_recon[m][n] < -2048 ) dct_recon[m][n] = -2048 ;
        if ( dct_zz[i] == 0 )
            dct_recon[m][n] = 0 ;
    }
}

```

dct_recon[m][n] = 0 for all m, n in skipped macroblocks and when pattern[i] == 0.

Once the DCT coefficients are reconstructed, the inverse DCT transform defined in annex A shall be applied to obtain the inverse transformed pel values in the interval [-256, 255]. The inverse DCT pel values shall be added to the pel[i][j] which were computed above using the motion vectors. The result of the addition shall be limited to the interval [0,255]. The location of the pels is determined from mb_row, mb_column and the pattern_code list.

2.4.4.3 Predictive-coded macroblocks in B-pictures

Predictive-coded macroblocks in B-Pictures are decoded in four steps.

First, the value of the forward motion vector for the macroblock is reconstructed from the retrieved forward motion vector information, and the forward motion vector reconstructed for the previous macroblock, using the same procedure as for calculating the forward motion vector in P-pictures. However, for B-pictures the previous reconstructed motion vectors shall be reset only for the first macroblock in a slice, or when the last macroblock that was decoded was an intra-coded macroblock. If no forward motion vector data exists for the current macroblock, the motion vectors shall be obtained by:

```

recon_right_for = recon_right_for_prev,
recon_down_for = recon_down_for_prev.

```

Second, the value of the backward motion vector for the macroblock shall be reconstructed from the retrieved backward motion vector information, and the backward motion vector reconstructed for the previous macroblock using the same procedure as for calculating the forward motion vector in B-pictures. In this procedure, the variables needed to find the backward motion vector are substituted for the variables needed to find the forward motion vector. The variables and coded data elements used to calculate the backward motion vector are:

```

recon_right_back_prev, recon_down_back_prev, backward_f_code, full_pel_backward_vector
motion_horizontal_backward_code, motion_horizontal_backward_r,
motion_vertical_backward_code, motion_vertical_backward_r,

```

backward_r_size and backward_f are derived from backward_f_code as follows:

```
backward_r_size = backward_f_code - 1
backward_f = 1 << backward_r_size
```

The following variables result from applying the algorithm in 2.4.4.2, modified as described in the previous paragraphs in this clause:

```
right_for    right_half_for    down_for    down_half_for
right_back   right_half_back   down_back   down_half_back
```

They define the integral and half pel value of the rightward and downward components of the forward motion vector (which references the past picture in display order) and the backward motion vector (which references the future picture in display order).

Third, the predictors of the pel values of the block being decoded, pel[], are calculated. If only forward motion vector information was retrieved for the macroblock, then pel[] of the decoded picture shall be calculated according to the formulas in 2.4.4.2. If only backward motion vector information was retrieved for the macroblock, then pel[] of the decoded picture shall be calculated according to the formulas in the predictive-coded macroblock clause, with "back" replacing "for", and pel_future[] replacing pel_past[]. If both forward and backward motion vectors information are retrieved, then let pel_for[] be the value calculated from the past picture by use of the reconstructed forward motion vector, and let pel_back[] be the value calculated from the future picture by use of the reconstructed backward motion vector. Then the value of pel[] shall be calculated by:

```
pel[] = ( pel_for[] + pel_back[] ) // 2 ;
```

Define non_intra_quant[m][n] to be the non-intra quantizer matrix that is specified in the sequence header.

Fourth, the DCT coefficients for each block present in the macroblock shall be reconstructed by any means equivalent to the following procedure:

```
for ( m=0; m<8; m++ ) {
    for ( n=0; n<8; n++ ) {
        i = scan[m][n] ;
        dct_recon[m][n] = ( ( ( 2 * dct_zz[i] ) + Sign(dct_zz[i] ) ) *
            quantizer_scale * non_intra_quant[m][n] ) / 16 ;
        if ( ( dct_recon[m][n] & 1 ) == 0 )
            dct_recon[m][n] = dct_recon[m][n] - Sign(dct_recon[m][n]) ;
        if (dct_recon[m][n] > 2 047) dct_recon[m][n] = 2 047 ;
        if (dct_recon[m][n] < -2 048) dct_recon[m][n] = -2 048 ;
        if ( dct_zz[i] == 0 )
            dct_recon[m][n] = 0 ;
    }
}
```

dct_recon[m][n] = 0 for all m, n in skipped macroblocks and when pattern[i] == 0.

Once the DCT coefficients are reconstructed, the inverse DCT transform defined in annex A shall be applied to obtain the inverse transformed pel values in the range [-256, 255]. The inverse DCT pel values shall be added to pel[], which were computed above from the motion vectors. The result of the addition shall be limited to the interval [0,255]. The location of the pels is determined from mb_row, mb_column and the pattern_code list.

2.4.4.4 Skipped macroblocks

For some macroblocks there are no coded data, that is neither motion vector information nor DCT information is available to the decoder. These macroblocks are called skipped macroblocks and are indicated when the macroblock_address_increment is greater than 1.

In I-pictures, all macroblocks shall be coded and there shall be no skipped macroblocks.

In P-pictures, the skipped macroblock is defined to be a macroblock with a reconstructed motion vector equal to zero and no DCT coefficients.

In B-pictures, the skipped macroblock is defined to have the same macroblock_type (forward, backward, or both motion vectors) as the prior macroblock, differential motion vectors equal to zero, and no DCT coefficients. In a B-picture, a skipped macroblock shall not follow an intra-coded macroblock.

2.4.4.5 Forced updating

This function is achieved by forcing the use of an intra-coded macroblock. The update pattern is not defined. For control of accumulation of IDCT mismatch error, each macroblock shall be intra-coded at least once per every 132 times it is coded in a P-picture without an intervening I-picture.

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Annex A

(normative)

8 by 8 Inverse discrete cosine transform

The 8 by 8 inverse discrete cosine transform for I-pictures and P-pictures shall conform to IEEE Draft Standard, P1180/D2, July 18, 1990. For B-pictures this specification may also be applied but may be unnecessarily stringent. Note that clause 2.3 of P1180/D2 "Considerations of Specifying IDCT Mismatch Errors" requires the specification of periodic intra-coding in order to control the accumulation of mismatch errors. The maximum refresh period requirement for this part of ISO/IEC 11172 shall be 132 intra-coded pictures or predictive-coded pictures as stated in 2.4.4.5, which is the same as indicated in P1180/D2 for visual telephony according to CCITT Recommendation H.261 [5].

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Annex B

(normative)

Variable length code tables

Introduction

This annex contains the variable length code tables for macroblock addressing, macroblock type, macroblock pattern, motion vectors, and DCT coefficients.

B.1 Macroblock addressing

Table B.1. -- Variable length codes for macroblock_address_increment.

macroblock_address_increment VLC code	increment value	macroblock_address_increment VLC code	increment value
1	1	0000 0101 10	17
011	2	0000 0101 01	18
010	3	0000 0101 00	19
0011	4	0000 0100 11	20
0010	5	0000 0100 10	21
0001 1	6	0000 0100 011	22
0001 0	7	0000 0100 010	23
0000 111	8	0000 0100 001	24
0000 110	9	0000 0100 000	25
0000 1011	10	0000 0011 111	26
0000 1010	11	0000 0011 110	27
0000 1001	12	0000 0011 101	28
0000 1000	13	0000 0011 100	29
0000 0111	14	0000 0011 011	30
0000 0110	15	0000 0011 010	31
0000 0101 11	16	0000 0011 001	32
		0000 0011 000	33
		0000 0001 111	macroblock_stuffing
		0000 0001 000	macroblock_escape

B.2 Macroblock type

The properties of the macroblock are determined by the macroblock type VLC according to these tables.

Table B.2a. -- Variable length codes for macroblock_type in intra-coded pictures (I-pictures).

macroblock_ typeVLC code	macroblock_ quant	macroblock_ motion_ forward	macroblock_ motion_ backward	macroblock_ pattern	macroblock_ intra
1	0	0	0	0	1
01	1	0	0	0	1

Table B.2b. -- Variable length codes for macroblock_type in predictive-coded pictures (P-pictures).

macroblock_ typeVLC code	macroblock_ quant	macroblock_ motion_ forward	macroblock_ motion_ backward	macroblock_ pattern	macroblock_ intra
1	0	1	0	1	0
01	0	0	0	1	0
001	0	1	0	0	0
00011	0	0	0	0	1
00010	1	1	0	1	0
00001	1	0	0	1	0
000001	1	0	0	0	1

Table B.2c. -- Variable length codes for macroblock_type in bidirectionally predictive-coded pictures (B-pictures).

macroblock_ typeVLC code	macroblock_ quant	macroblock_ motion_ forward	macroblock_ motion_ backward	macroblock_ pattern	macroblock_ intra
10	0	1	1	0	0
11	0	1	1	1	0
010	0	0	1	0	0
011	0	0	1	1	0
0010	0	1	0	0	0
0011	0	1	0	1	0
00011	0	0	0	0	1
00010	1	1	1	1	0
000011	1	1	0	1	0
000010	1	0	1	1	0
000001	1	0	0	0	1

Table B.2d. -- Variable length codes for macroblock_type in dc intra-coded pictures (D-pictures).

macroblock_ typeVLC code	macroblock_ quant	macroblock_ motion_ forward	macroblock_ motion_ backward	macroblock_ pattern	macroblock_ intra
1	0	0	0	0	1

B.3 Macroblock pattern

Table B.3. -- Variable length codes for coded_block_pattern.

coded_block_pattern VLC code	cbp	coded_block_pattern VLC code	cbp
111	60	0001 1100	35
1101	4	0001 1011	13
1100	8	0001 1010	49
1011	16	0001 1001	21
1010	32	0001 1000	41
1001 1	12	0001 0111	14
1001 0	48	0001 0110	50
1000 1	20	0001 0101	22
1000 0	40	0001 0100	42
0111 1	28	0001 0011	15
0111 0	44	0001 0010	51
0110 1	52	0001 0001	23
0110 0	56	0001 0000	43
0101 1	1	0000 1111	25
0101 0	61	0000 1110	37
0100 1	2	0000 1101	26
0100 0	62	0000 1100	38
0011 11	24	0000 1011	29
0011 10	36	0000 1010	45
0011 01	3	0000 1001	53
0011 00	63	0000 1000	57
0010 111	5	0000 0111	30
0010 110	9	0000 0110	46
0010 101	17	0000 0101	54
0010 100	33	0000 0100	58
0010 011	6	0000 0011 1	31
0010 010	10	0000 0011 0	47
0010 001	18	0000 0010 1	55
0010 000	34	0000 0010 0	59
0001 1111	7	0000 0001 1	27
0001 1110	11	0000 0001 0	39
0001 1101	19		

B.4 Motion vectors

Table B.4. -- Variable length codes for motion_horizontal_forward_code, motion_vertical_forward_code, motion_horizontal_backward_code, and motion_vertical_backward_code.

motion VLC code	code
0000 0011 001	-16
0000 0011 011	-15
0000 0011 101	-14
0000 0011 111	-13
0000 0100 001	-12
0000 0100 011	-11
0000 0100 11	-10
0000 0101 01	-9
0000 0101 11	-8
0000 0111	-7
0000 1001	-6
0000 1011	-5
0000 111	-4
0001 1	-3
0011	-2
011	-1
1	0
010	1
0010	2
0001 0	3
0000 110	4
0000 1010	5
0000 1000	6
0000 0110	7
0000 0101 10	8
0000 0101 00	9
0000 0100 10	10
0000 0100 010	11
0000 0100 000	12
0000 0011 110	13
0000 0011 100	14
0000 0011 010	15
0000 0011 000	16

B.5 DCT coefficients

Table B.5a -- Variable length codes for dct_dc_size_luminance.

VLC code	dct_dc_size_luminance
100	0
00	1
01	2
101	3
110	4
1110	5
11110	6
111110	7
1111110	8

Table B.5b. -- Variable length codes for dct_dc_size_chrominance.

VLC code	dct_dc_size_chrominance
00	0
01	1
10	2
110	3
1110	4
11110	5
111110	6
1111110	7
11111110	8

Table B.5c. -- Variable length codes for dct_coeff_first and dct_coeff_next.

dct_coeff_first and dct_coeff_next variable length code (NOTE1)	run	level
10	end_of_block	
1 s (NOTE2)	0	1
11 s (NOTE3)	0	1
011 s	1	1
0100 s	0	2
0101 s	2	1
0010 1 s	0	3
0011 1 s	3	1
0011 0 s	4	1
0001 10 s	1	2
0001 11 s	5	1
0001 01 s	6	1
0001 00 s	7	1
0000 110 s	0	4
0000 100 s	2	2
0000 111 s	8	1
0000 101 s	9	1
0000 01	escape	
0010 0110 s	0	5
0010 0001 s	0	6
0010 0101 s	1	3
0010 0100 s	3	2
0010 0111 s	10	1
0010 0011 s	11	1
0010 0010 s	12	1
0010 0000 s	13	1
0000 0010 10 s	0	7
0000 0011 00 s	1	4
0000 0010 11 s	2	3
0000 0011 11 s	4	2
0000 0010 01 s	5	2
0000 0011 10 s	14	1
0000 0011 01 s	15	1
0000 0010 00 s	16	1
NOTES 1 - The last bit 's' denotes the sign of the level, '0' for positive '1' for negative. 2 - This code shall be used for dct_coeff_first. 3 - This code shall be used for dct_coeff_next.		

Table B.5d. -- Variable length codes for dct_coeff_first and dct_coeff_next.

dct_coeff_first and dct_coeff_next variable length code (NOTE)	run	level
0000 0001 1101 s	0	8
0000 0001 1000 s	0	9
0000 0001 0011 s	0	10
0000 0001 0000 s	0	11
0000 0001 1011 s	1	5
0000 0001 0100 s	2	4
0000 0001 1100 s	3	3
0000 0001 0010 s	4	3
0000 0001 1110 s	6	2
0000 0001 0101 s	7	2
0000 0001 0001 s	8	2
0000 0001 1111 s	17	1
0000 0001 1010 s	18	1
0000 0001 1001 s	19	1
0000 0001 0111 s	20	1
0000 0001 0110 s	21	1
0000 0000 1101 0 s	0	12
0000 0000 1100 1 s	0	13
0000 0000 1100 0 s	0	14
0000 0000 1011 1 s	0	15
0000 0000 1011 0 s	1	6
0000 0000 1010 1 s	1	7
0000 0000 1010 0 s	2	5
0000 0000 1001 1 s	3	4
0000 0000 1001 0 s	5	3
0000 0000 1000 1 s	9	2
0000 0000 1000 0 s	10	2
0000 0000 1111 1 s	22	1
0000 0000 1111 0 s	23	1
0000 0000 1110 1 s	24	1
0000 0000 1110 0 s	25	1
0000 0000 1101 1 s	26	1

NOTE - The last bit 's' denotes the sign of the level, '0' for positive, '1' for negative.

Table B.5e. -- Variable length codes for dct_coeff_first and dct_coeff_next (concluded).

dct_coeff_first and dct_coeff_next variable length code (NOTE)	run	level
0000 0000 0111 11 s	0	16
0000 0000 0111 10 s	0	17
0000 0000 0111 01 s	0	18
0000 0000 0111 00 s	0	19
0000 0000 0110 11 s	0	20
0000 0000 0110 10 s	0	21
0000 0000 0110 01 s	0	22
0000 0000 0110 00 s	0	23
0000 0000 0101 11 s	0	24
0000 0000 0101 10 s	0	25
0000 0000 0101 01 s	0	26
0000 0000 0101 00 s	0	27
0000 0000 0100 11 s	0	28
0000 0000 0100 10 s	0	29
0000 0000 0100 01 s	0	30
0000 0000 0100 00 s	0	31
0000 0000 0011 000 s	0	32
0000 0000 0010 111 s	0	33
0000 0000 0010 110 s	0	34
0000 0000 0010 101 s	0	35
0000 0000 0010 100 s	0	36
0000 0000 0010 011 s	0	37
0000 0000 0010 010 s	0	38
0000 0000 0010 001 s	0	39
0000 0000 0010 000 s	0	40
0000 0000 0011 111 s	1	8
0000 0000 0011 110 s	1	9
0000 0000 0011 101 s	1	10
0000 0000 0011 100 s	1	11
0000 0000 0011 011 s	1	12
0000 0000 0011 010 s	1	13
0000 0000 0011 001 s	1	14
0000 0000 0001 0011 s	1	15
0000 0000 0001 0010 s	1	16
0000 0000 0001 0001 s	1	17
0000 0000 0001 0000 s	1	18
0000 0000 0001 0100 s	6	3
0000 0000 0001 1010 s	11	2
0000 0000 0001 1001 s	12	2
0000 0000 0001 1000 s	13	2
0000 0000 0001 0111 s	14	2
0000 0000 0001 0110 s	15	2
0000 0000 0001 0101 s	16	2
0000 0000 0001 1111 s	27	1
0000 0000 0001 1110 s	28	1
0000 0000 0001 1101 s	29	1
0000 0000 0001 1100 s	30	1
0000 0000 0001 1011 s	31	1
NOTE - The last bit 's' denotes the sign of the level, '0' for positive, '1' for negative.		

Table B.5f. -- Encoding of run and level following an escape code either as a 14-bit fixed length code (-127 ≤ level ≤ 127) or as a 22-bit fixed length code (-255 ≤ level ≤ -128, 128 ≤ level ≤ 255).

(Note - This yields total escape code lengths of 20-bits and 28-bits respectively).

fixed length code	run
0000 00	0
0000 01	1
0000 10	2
...	...
...	...
...	...
...	...
...	...
...	...
1111 11	63

fixed length code	level
forbidden	-256
1000 0000 0000 0001	-255
1000 0000 0000 0010	-254
...	...
1000 0000 0111 1111	-129
1000 0000 1000 0000	-128
1000 0001	-127
1000 0010	-126
...	...
1111 1110	-2
1111 1111	-1
forbidden	0
0000 0001	1
...	...
0111 1111	127
0000 0000 1000 0000	128
0000 0000 1000 0001	129
...	...
0000 0000 1111 1111	255

Annex C

(normative)

Video buffering verifier

Constant rate coded video bitstreams shall meet constraints imposed through a Video Buffering Verifier (VBV) defined in clause C.1.

The VBV is a hypothetical decoder which is conceptually connected to the output of an encoder. Coded data are placed in the input buffer of the model decoder at the constant bitrate that is being used. Coded data is removed from the buffer as defined in C.1.4, below. It is a requirement of the encoder (or editor) that the bitstream it produces will not cause the VBV input buffer to either overflow or underflow.

C.1 Video buffering verifier

C.1.1 The VBV and the video encoder have the same clock frequency as well as the same picture rate, and are operated synchronously.

C.1.2 The VBV has an input buffer of size B, where B is given in the `vbv_buffer_size` field in the sequence header.

C.1.3 The VBV input buffer is initially empty. After filling the input buffer with all the data that precedes the first picture start code and the picture start code itself, the input buffer is filled from the bitstream for the time specified by the `vbv_delay` field in the video bitstream.

C.1.4 All of the picture data for the picture that has been in the buffer longest is instantaneously removed. Then after each subsequent picture interval all of the picture data for the picture which at that time has been in the buffer longest is instantaneously removed.

For the purposes of this annex picture data includes any sequence header and group of picture layer data that immediately precede the picture start code as well as all the picture data elements and any trailing stuffing bits or bytes. For the first coded picture in the video sequence, any zero bit or byte stuffing immediately preceding the sequence header is also included in the picture data.

The VBV buffer is examined immediately before removing any picture data and immediately after this picture data is removed. Each time the VBV is examined its occupancy shall lie between zero bits and B bits where, B is the size of the VBV buffer indicated by `vbv_buffer_size` in the sequence header.

This is a requirement for the entire video bitstream.

To meet these requirements the number of bits for the (n+1)'th coded picture d_{n+1} shall satisfy

$$d_{n+1} > B_n + (2R/P) - B$$

$$d_{n+1} \leq B_n + (R/P)$$

Real-valued arithmetic is used in these inequalities.

where

$$n \geq 0$$

B = VBV receiving buffer size given by `vbv_buffer_size` * 16 384 bits.

B_n = the buffer occupancy (measured in bits) just after time t_n

R = bitrate measured in bits/s. The full precision of the bitrate rather than the rounded value encoded by the `bit_rate` field in the sequence header shall be used by the encoder in the VBV model.

P = nominal number of pictures per second

t_n = the time when the n'th coded picture is removed from the VBV buffer

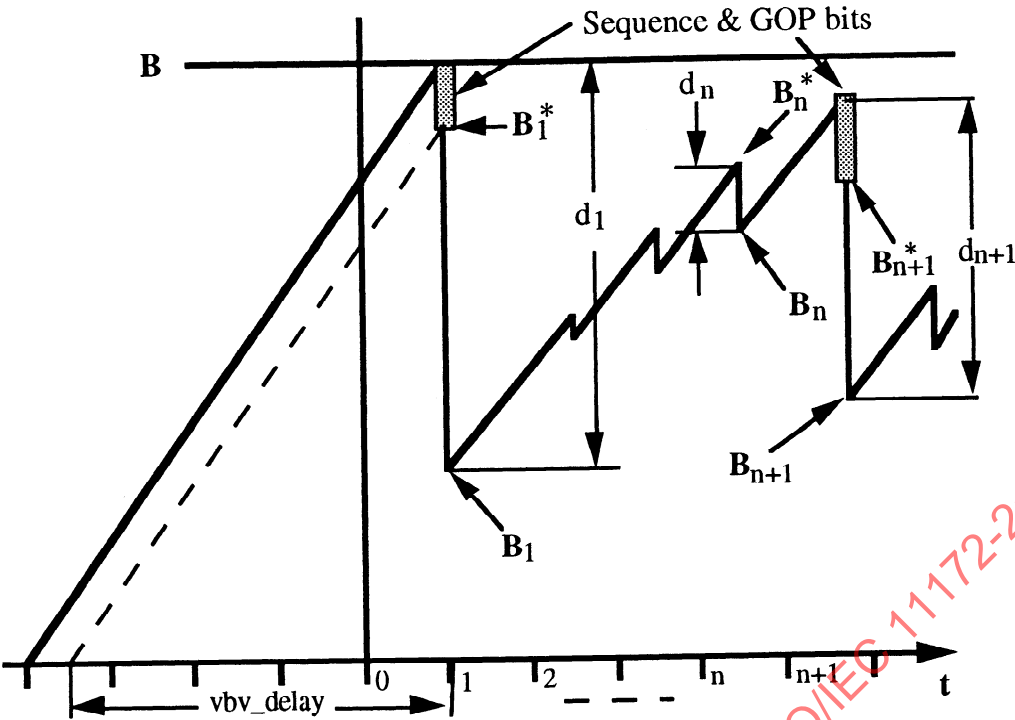


Figure C.1 -- VBV buffer occupancy

Annex D

(informative)

Guide to encoding video

D.1 Introduction

This annex provides background material to help readers understand and implement this part of ISO/IEC 11172. The normative clauses of this part of ISO/IEC 11172 do not specify the design of a decoder. They provide even less information about encoders; they do not specify what algorithms encoders should employ in order to produce a valid bitstream. The normative material is written in a concise form and contains few examples; consequently is not easy to understand. This annex attempts to address this problem by explaining coding methods, giving examples, and discussing encoding and decoding algorithms which are not directly covered by this part of ISO/IEC 11172.

The normative clauses specify the bitstream in such a way that it is fairly straightforward to design a compliant decoder. Decoders may differ considerably in architecture and implementation details, but have very few choices during the decoding process: the methods and the results of the decoding process are closely specified. Decoders do have some freedom in methods of post processing and display, but the results of such post processing cannot be used in subsequent decoding steps.

The situation is quite different for encoders. This part of ISO/IEC 11172 does not specify how to design or implement an encoder which produces good quality video. This annex devotes a major part to discussing encoder algorithms.

This part of ISO/IEC 11172 was developed by ISO/IEC JTC1/SC29/WG11 which is widely known as MPEG (Moving Pictures Expert Group). This part of ISO/IEC 11172 was developed in response to industry needs for an efficient way of storing and retrieving audio and video information on digital storage media (DSM). CD-ROM is an inexpensive medium which can deliver data at approximately 1,2 Mbits/s, and this part of ISO/IEC 11172 was aimed at approximately this data rate. The "constrained parameters bitstream", a subset of all permissible bitstreams that is expected to be widely used, is limited to data rates up to 1 856 000 bits/s. However, it should be noted that this part of ISO/IEC 11172 is not limited to this value and may be used at higher data rates.

Two other relevant International Standards were being developed during the work of the MPEG video committee: H.261 by CCITT aimed at telecommunications applications [5], and ISO/IEC 10918 by the ISO/IEC JTC1/SC29 (JPEG) committee aimed at the coding of still pictures [6]. Elements of both of these standards were incorporated into this part of ISO/IEC 11172, but subsequent development work by the committee resulted in coding elements that are new to this part of ISO/IEC 11172. Le Gall [2] gives an account of the method by which ISO/IEC JTC1/SC29/WG11 (MPEG) developed this part of ISO/IEC 11172, and a summary of this part of ISO/IEC 11172 itself.

D.2 Overview

D.2.1 Video concepts

This part of ISO/IEC 11172 defines a format for compressed digital video. This annex describes some ways in which practical encoders and decoders might be implemented.

Although this part of ISO/IEC 11172 is quite flexible, the basic algorithms have been tuned to work well at data rates of about 1 to 1,5 M bits/s, at spatial resolutions of about 350 pels horizontally by about 250 pels vertically, and picture rates of about 24 to 30 pictures/s. The use of the word "picture" as opposed to "frame" is deliberate. This part of ISO/IEC 11172 codes progressively-scanned images and does not recognize the concept of interlace. Interlaced source video must be converted to a non-interlaced format before coding. After decoding, the decoder may optionally produce an interlaced format for display.

This part of ISO/IEC 11172 is designed to permit several methods of viewing coded video which are normally associated with VCRs such as forward playback, freeze picture, fast forward, fast reverse, and slow

forward. In addition, random access may be possible. The ability of the decoder to implement these modes depends to some extent on the nature of the digital storage medium on which the coded video is stored.

The overall process of encoding and decoding is illustrated below:

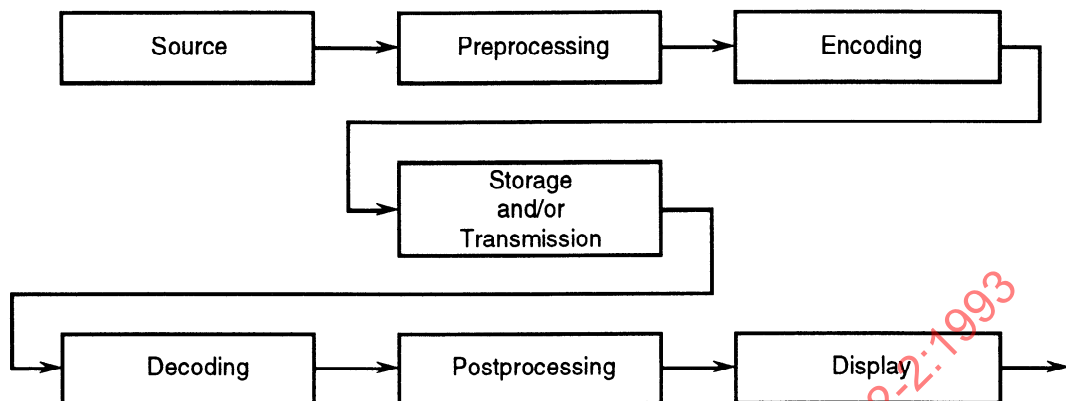


Figure D.1 -- Coding and decoding process

Figure D.1 shows a typical sequence of operations that must be performed before moving pictures can be seen by a viewer. The unencoded source may exist in many forms, such as the CCIR 601 format. Clause D.3 describes how such a source may be converted into the appropriate resolution for subsequent encoding. In the encoding step, the encoder must be aware of the decoder buffer capacity, and the need of the decoder to match the rate of the media to the rate of filling the picture buffer with each successive picture. To this end, a model of the decoder buffer and its overflow and underflow problem is introduced in D.4, and rate control is described in D.6.1 The structure of an ISO/IEC 11172-2 bitstream is covered in D.5, as are the coding operations that compress the video. Following the encoding process, the bitstream may be copied to a storage medium. To view the moving pictures, the decoder accesses the ISO/IEC 11172-2 bitstream, and decodes it as described in D.7. Postprocessing for display is described in D.8.

D.2.2 MPEG video compression techniques

Video is represented as a succession of individual pictures, and each picture is treated as a two-dimensional array of picture elements (pels). The colour representation for each pel consists of three components: Y (luminance), and two chrominance components, Cb and Cr.

Compression of digitized video comes from the use of several techniques: subsampling of the chrominance information to match the sensitivity of the human visual system (HVS), quantization, motion compensation (MC) to exploit temporal redundancy, frequency transformation by discrete cosine transform (DCT) to exploit spatial redundancy, variable length coding (VLC), and picture interpolation.

D.2.2.1 Subsampling of chrominance information

The HVS is most sensitive to the resolution of an image's luminance component, so the Y pel values are encoded at full resolution. The HVS is less sensitive to the chrominance information. Subsampling reduces the number of pel values by systematically combining them with a type of averaging process. This reduces the amount of information to be compressed by other techniques. The International Standard retains one set of chrominance pels for each 2x2 neighbourhood of luminance pels.

D.2.2.2 Quantization

Quantization represents a range of values by a single value in the range. For example, converting a real number to the nearest integer is a form of quantization. The quantized range can be concisely represented as an integer code, which can be used to recover the quantized value during decoding. The difference between the actual value and the quantized value is called the quantization noise. Under some circumstances, the HVS is less sensitive to quantization noise so such noise can be allowed to be large, thus increasing coding efficiency.

D.2.2.3 Predictive coding

Predictive coding is a technique to improve the compression through statistical redundancy. Based on values of pels previously decoded, both the encoder and decoder can estimate or predict the value of a pel yet to be encoded or decoded. The difference between the predicted and actual values is encoded. This difference value is the prediction error which the decoder can use to correct the prediction. Most error values will be small and cluster around the value 0 since pel values typically do not have large changes within a small spatial neighbourhood. The probability distribution of the prediction error is skewed and compresses better than the distribution of the pel values themselves. Additional information can be discarded by quantizing the prediction error. In this International Standard predictive coding is also used for the dc-values of successive luminance or chrominance blocks and in the encoding of motion vectors.

D.2.2.4 Motion compensation and interframe coding

Motion compensation (MC) predicts the values of a block pels in a picture by relocating a block of neighbouring pel values from a known picture. The motion is described as a two-dimensional motion vector that specifies where to retrieve a block of pel values from a previously decoded picture that is used to predict pel values of the current block. The simplest example is a scene where the camera is not moving, and no objects in the scene are moving. The pel values at each image location remain the same, and the motion vector for each block is 0. In general however, the encoder may transmit a motion vector for each macroblock. The translated block from the known picture becomes a prediction for the block in the picture to be encoded. The technique relies on the fact that within a short sequence of pictures of the same general scene, many objects remain in the same location while others move only a short distance.

D.2.2.5 Frequency transformation

The discrete cosine transform (DCT) converts an 8 by 8 block of pel values to an 8 by 8 matrix of horizontal and vertical spatial frequency coefficients. An 8 by 8 block of pel values can be reconstructed by performing the inverse discrete cosine transform (IDCT) on the spatial frequency coefficients. In general, most of the energy is concentrated in the low frequency coefficients, which are conventionally written in the upper left corner of the transformed matrix. Compression is achieved by a quantization step, where the quantization intervals are identified by an index. Since the encoder identifies the interval and not the exact value within the interval, the pel values of the block reconstructed by the IDCT have reduced accuracy.

The DCT coefficient in location (0,0) (upper left) of the block represents the zero horizontal and zero vertical frequency and is called the dc coefficient. The dc coefficient is proportional to the average pel value of the 8 by 8 block, and additional compression is provided through predictive coding since the difference in the average value of neighbouring 8 by 8 blocks tends to be relatively small. The other coefficients represent one or more nonzero horizontal or nonzero vertical spatial frequencies, and are called ac coefficients. The quantization level of the coefficients corresponding to the higher spatial frequencies favors the creation of an ac coefficient of 0 by choosing a quantization step size such that the HVS is unlikely to perceive the loss of the particular spatial frequency unless the coefficient value lies above the particular quantization level. The statistical encoding of the expected runs of consecutive zero-valued coefficients of higher-order coefficients accounts for considerable compression gain. To cluster nonzero coefficients early in the series and encode as many zero coefficients as possible following the last nonzero coefficient in the ordering, the coefficient sequence is specified to be a zig-zag ordering; see figure D.30. The ordering concentrates the highest spatial frequencies at the end of the series.

D.2.2.6 Variable-length coding

Variable-length coding (VLC) is a statistical coding technique that assigns codewords to values to be encoded. Values of high frequency of occurrence are assigned short codewords, and those of infrequent occurrence are assigned long codewords. On average, the more frequent shorter codewords dominate, such that the code string is shorter than the original data.

D.2.2.7 Picture interpolation

If the decoder reconstructs a picture from the past and a picture from the future, then the intermediate pictures can be reconstructed by the technique of interpolation, or bidirectional prediction. Blocks in the intermediate pictures can be forward and backward predicted and translated by means of motion vectors. The decoder may reconstruct pel values belonging to a given block as an average of values from the past and future pictures.

D.2.3 Bitstream hierarchy

The ISO/IEC 11172-2 coding scheme is arranged in layers corresponding to a hierarchical structure. A **sequence** is the top layer of the coding hierarchy and consists of a header and some number of **groups-of-pictures (GOPs)**. The sequence header initializes the state of the decoder. This allows decoders to decode any sequence without being affected by past decoding history.

A **GOP** is a random access point, i.e. it is the smallest coding unit that can be independently decoded within a sequence, and consists of a header and some number of **pictures**. The GOP header contains time and editing information.

A **picture** corresponds to a single frame of motion video, or to a movie frame. There are four picture types: **I-pictures**, or *intra coded pictures*, which are coded without reference to any other pictures; **P-pictures**, or *predictive coded pictures*, which are coded using motion compensation from a previous I or P-picture; **B-pictures**, or *bidirectionally predictive coded pictures*, which are coded using motion compensation from a previous and a future I or P-picture, and **D pictures**, or *D pictures*, which are intended only for a fast forward search mode. A typical coding scheme contains a mix of I, P, and B-pictures. Typically, an I-picture may occur every half a second, to give reasonably fast random access, with two B-pictures inserted between each pair of I or P-pictures.

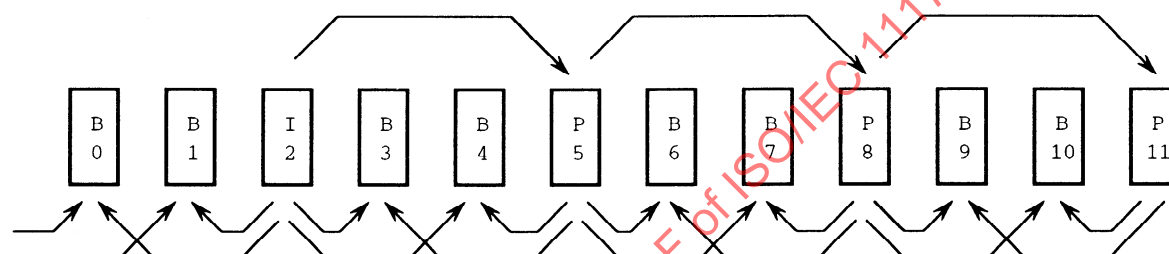


Figure D.2 -- Dependency relationship between I, B, and P-pictures

Figure D.2 illustrates a number of pictures in display order. The arrows show the dependency relationship of the predictive and bidirectionally predictive coded pictures.

Note that because of the picture dependencies, the bitstream order, i.e. the order in which pictures are transmitted, stored, or retrieved, is not the display order, but rather the order which the decoder requires them to decode the bitstream. An example of a sequence of pictures, in display order, might be:

I	B	B	P	B	B	P	B	B	P	B	B	I	B	B	P	B	B	P
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18

Figure D.3 -- Typical sequence of pictures in display order

whereas the bitstream order would be as shown below:

I	P	B	B	P	B	B	P	B	B	I	B	B	P	B	B	P	B	B
0	3	1	2	6	4	5	9	7	8	12	10	11	15	13	14	18	16	17

Figure D.4 -- Typical sequence of pictures in bitstream order

Because the B-pictures depend on the following (in display order) I or P-picture, the I or P-picture must be transmitted and decoded before the dependent B-pictures.

Pictures consist of a header and one or more **slices**. The picture header contains time, picture type, and coding information.

A **slice** provides some immunity to data corruption. Should the bitstream become unreadable within a picture, the decoder should be able to recover by waiting for the next slice, without having to drop an entire picture.

Slices consist of a header and one or more **macroblocks**. At the start of each slice all of the predictors, for dc values and motion vectors, are reset. The slice header contains position and quantizer scale information. This is sufficient for recovery from local corruption.

A **macroblock** is the basic unit for motion compensation and quantizer scale changes.

Each **macroblock** consists of a header and six component 8 by 8 **blocks**: four blocks of luminance, one block of Cb chrominance, and one block of Cr chrominance. See figure D.5. The macroblock header contains quantizer scale and motion compensation information.

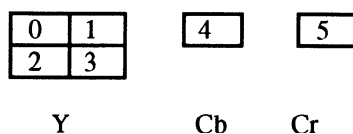


Figure D.5 -- Macroblock structure

A **macroblock** contains a 16-pel by 16-line section of luminance component and the spatially corresponding 8-pel by 8-line section of each chrominance component. A skipped macroblock is one for which no information is stored (see 2.4.4.4).

Note that the picture area covered by the four blocks of luminance is the same as the area covered by each of the chrominance blocks. This is due to subsampling of the chrominance information.

Blocks are the basic coding unit, and the DCT is applied at this block level. Each block contains 64 component **pels** arranged in an 8 by 8 array as shown in figure D.6.

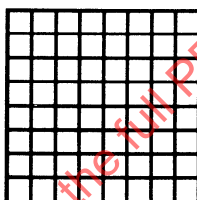


Figure D.6 -- Block structure

D.2.4 Decoder overview

A simplified block diagram of a possible decoder implementation is shown below:

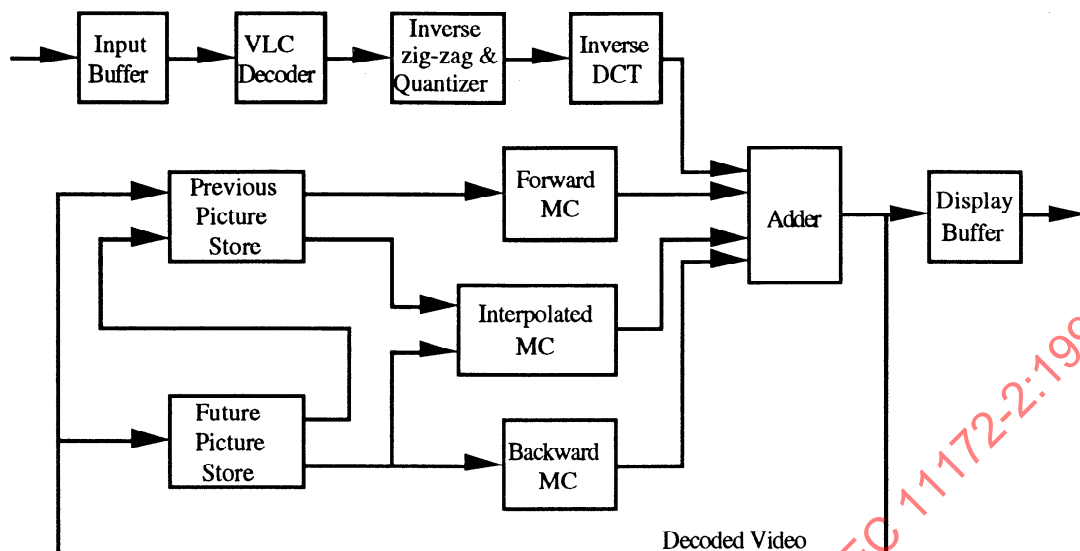


Figure D.7 -- Simplified decoder block diagram

It is instructive to follow the method which the decoder uses to decode a bitstream containing the sequence of pictures given in Fig D.4, and display them in the order given in Fig D.3. The following description is simplified for clarity.

The input bitstream is accumulated in the Input Buffer until needed. The Variable Length Code (VLC) Decoder decodes the header of the first picture, picture 0, and determines that it is an I-picture. The VLC Decoder produces quantized coefficients corresponding to the quantized DCT coefficients. These are assembled for each 8 by 8 block of pels in the image by inverse zig-zag scanning. The Inverse Quantizer produces the actual DCT coefficients using the quantization step size. The coefficients are then transformed into pel values by the Inverse DCT transformer and stored in the Previous Picture Store and the Display Buffer. The picture may be displayed at the appropriate time.

The VLC Decoder decodes the header of the next picture, picture 3, and determines that it is a P-picture. For each block, the VLC Decoder decodes motion vectors giving the displacement from the stored previous picture, and quantized coefficients corresponding to the quantized DCT coefficients of the difference block. These quantized coefficients are inverse quantized to produce the actual DCT coefficients. The coefficients are then transformed into pel difference values and added to the predicted block produced by applying the motion vectors to blocks in the stored previous picture. The resultant block is stored in the Future Picture Store and the Display Buffer. This picture cannot be displayed until B-pictures 1 and 2 have been received, decoded, and displayed.

The VLC Decoder decodes the header of the next picture, picture 1, and determines that it is a B-picture. For each block, the VLC decoder decodes motion vectors giving the displacement from the stored previous or future pictures or both, and quantized coefficients corresponding to the quantized DCT coefficients of the difference block. These quantized coefficients are inverse quantized to produce the actual DCT coefficients. The coefficients are then inverse transformed into difference pel values and added to the predicted block produced by applying the motion vectors to the stored pictures. The resultant block is then stored in the Display Buffer. It may be displayed at the appropriate time.

The VLC Decoder decodes the header of the next picture, picture 2, and determines that it is a B-picture. It is decoded using the same method as for picture 1. After decoding picture 2, picture 0, which is in the Previous Picture Store, is no longer needed and may be discarded.

The VLC Decoder decodes the header of the next picture, picture 6, and determines that it is a P-picture. The picture in the Future Picture Store is copied into the Previous Picture Store, then decoding proceeds as for picture 3. Picture 6 should not be displayed until pictures 4 and 5 have been received and displayed.

The VLC Decoder decodes the header of the next picture, picture 4, and determines that it is a B-picture. It is decoded using the same method as for picture 1.

The VLC Decoder decodes the header of the next picture, picture 5, and determines that it is a B-picture. It is decoded using the same method as for picture 1.

The VLC Decoder decodes the header of the next picture, picture 9, and determines that it is a P-picture. It then proceeds as for picture 6.

The VLC Decoder decodes the header of the next picture, picture 7, and determines that it is a B-picture. It is decoded using the same method as for picture 1.

The VLC Decoder decodes the header of the next picture, picture 8, and determines that it is a B-picture. It is decoded using the same method as for picture 1.

The VLC Decoder decodes the header of the next picture, picture 12, and determines that it is an I-picture. It is decoded using the same method as for picture 0. This process is repeated for the subsequent pictures.

D.3 Preprocessing

The source material may exist in many forms, e.g. computer files or CCIR 601 format, but in general, it must be processed before being encoded. This clause discusses some aspects of preprocessing.

For a given data rate and source material, there is an optimum picture rate and spatial resolution at which to code if the best perceived quality is desired. If the resolution is too high, then too many bits will be expended on the overhead associated with each block leaving too few to code the values of each pel accurately. If the resolution is too low, the pel values will be rendered accurately, but high frequency detail will be lost. The optimum resolution represents a tradeoff between the various coding artifacts (e.g. noise and blockiness) and the perceived resolution and sharpness of the image. This tradeoff is further complicated by the unknowns of the final viewing conditions, e.g. screen brightness and the distance of the viewer from the screen.

At data rates of 1 to 1,5 Mbits/s, reasonable choices are: picture rates of 24, 25 and 30 pictures/s, a horizontal resolution of between 250 and 400 pels, and a vertical resolution of between 200 and 300 lines. Note that these values are not normative and other picture rates and resolutions are valid.

D.3.1 Conversion from CCIR 601 video to MPEG SIF

The two widely used scanning standards for colour television are 525 and 625 lines at 29,97 and 25 frames/s respectively. The number of lines containing picture information in the transmitted signal is 484 for the 525-line system and 576 for the 625-line system. Both use interlaced scanning with two fields per picture.

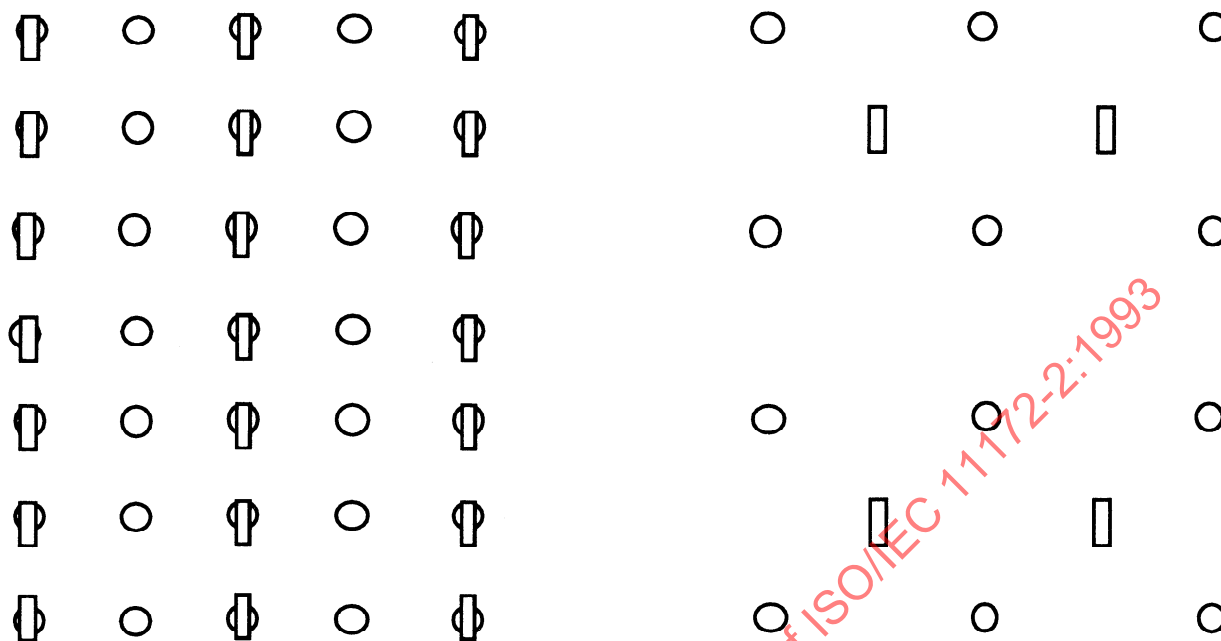
CCIR Recommendation 601 defines standards for the digital coding of colour television signals in component form. Of these the 4:2:2 standard has become widely adopted; the sampling frequency used for the luminance signal, Y, is 13,5 MHz and the two colour difference signals, Cb or B-Y and Cr or R-Y, are both sampled at 6,75 MHz. The number of luminance samples in the digital active line is 720 but only about 702 will be used in practice by the analogue active line.

The number of picture elements in the height and width of the picture, in the standards defined above, are too large for effective coding at data rates between 1 and 1,5 Mbit/s. More appropriate values are obtained by decreasing the resolution in both directions to a half. This reduces the pel rate by a factor of four. Interlace should be avoided as it increases the difficulties in achieving low data rates.

One way to reduce the vertical resolution is to use only the odd or the even fields. If the other field is simply discarded, spatial aliasing will be introduced, and this may produce visible and objectionable artifacts. More sophisticated methods of rate conversion require more computational power, but can perceptibly reduce the aliasing artifacts.

The horizontal and vertical resolutions may be halved by filtering and subsampling. Consider a picture in the 4:2:2 format. See the CCIR 601 sampling pattern of figure D.8(a). Such a sampling pattern may be converted to the SIF sampling pattern of figure D.8(b) as follows. The odd field only may be extracted, reducing the number of lines by two, and then a horizontal decimation filter used on the remaining lines to

reduce the horizontal resolution by a factor of two. In addition the chrominance values may be vertically decimated. The filters for luminance and chrominance have to be chosen carefully since particular attention has to be given to the location of the samples in the respective International Standards. The temporal relationship between luminance and chrominance must also be correct.



(a) Sampling pattern for 4:2:2 (CCIR 601)

(b) Sampling pattern for MPEG (SIF)

Circles represent luminance; Boxes represent Chrominance

Figure D.8 -- Conversion of CCIR 601 to SIF

The following 7-tap FIR filter has been found to give good results in decimating the luminance:

$$\begin{bmatrix} -29 & 0 & 88 & 138 & 88 & 0 & -29 \end{bmatrix} // 256$$

Figure D.9 -- Luminance subsampling filter tap weights

Use of a power of two for the divisor allows a simple hardware implementation.

The chrominance samples have to appear in the between the luminance samples both horizontally and vertically. The following linear filter with a phase shift of half a pel may be found useful.

$$\begin{bmatrix} 1 & 3 & 3 & 1 \end{bmatrix} // 8$$

Figure D.10 -- Chrominance subsampling filter tap weights

To recover the samples consistent with the CCIR 601 grid of figure D.8(a), the process of interpolation is used. The interpolation filter applied to a zero-padded signal can be chosen to be equal to the decimation filter employed for the luminance and the two chrominance values in the encoder.

Note that these filters are not part of the International Standard, and other filters may be used.

At the end of the lines some special technique such as renormalizing the filter or replicating the last pel, must be adopted. The following example shows a horizontal line of 16 luminance pels and the same line after filtering and subsampling. In this example the data in the line is reflected at each end.

10 12 20 30 35 15 19 11 11 19 26 45 80 90 92 90
12 32 23 9 12 49 95 92

Figure D.11 -- Example of filtering and subsampling of a line of pels

The result of this filtering and subsampling is a source input format (SIF) which has a luminance resolution of 360×240 or 360×288 , and a chrominance resolution which is half that of the luminance in each dimension.

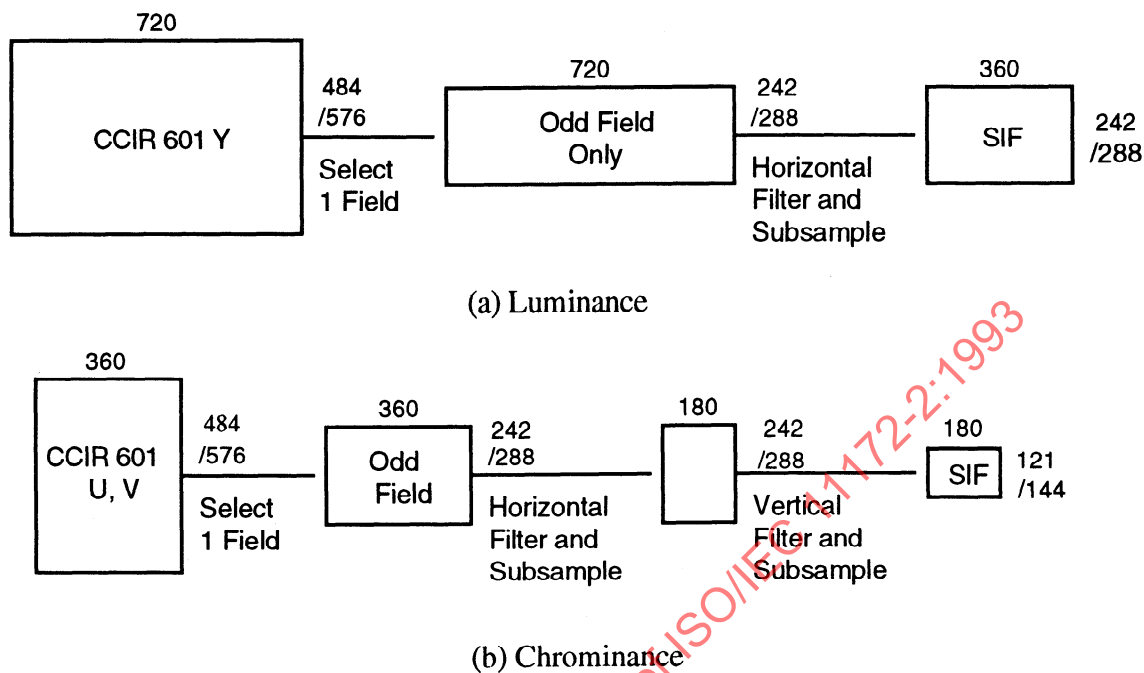


Figure D.12 -- Conversion from CCIR 601 into SIF

The SIF is not quite optimum for processing by MPEG video coders. MPEG video divides the luminance component into macroblocks of 16×16 pels. The horizontal resolution, 360, is not divisible by 16. The same is true of the vertical resolution, 242, in the case of 525-line systems. A better match is obtained in the horizontal direction by discarding the 4 pels at the end of every line of the subsampled picture. Care must be taken that this results in the correct configuration of luminance and chrominance samples in the macroblock. The remaining picture is called the significant pel area, and corresponds to the dark area in figure D.13:

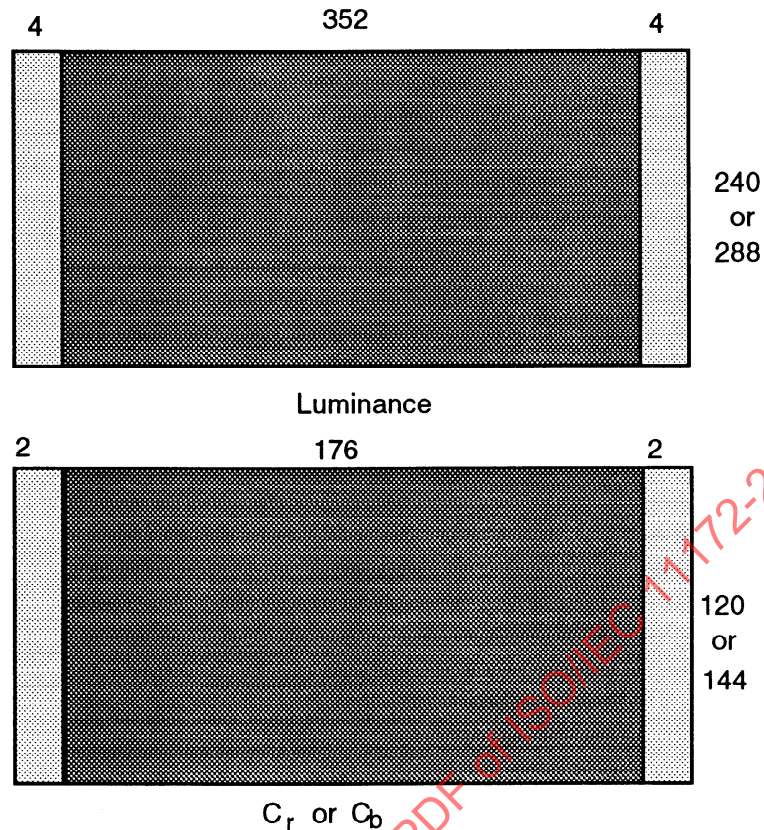


Figure D.13 -- Source input with significant pel area shaded dark

The conversion process is summarized in table D.1.

Table D.1 -- Conversion of source formats

Picture Rate (Hz)	29,97	25
Picture Aspect Ratio (width:height)	4:3	4:3
Luminance (Y)		
CCIR Sample Resolution	720 x 484	720 x 576
SIF	360 x 242	360 x 288
Significant Pel Area	352 x 240	352 x 288
Chrominance (Cb Cr)		
CCIR Sample Resolution	360 x 484	360 x 576
SIF	180 x 121	180 x 144
Significant Pel Area	176 x 120	176 x 144

The preprocessing into the SIF format is not normative, other processing steps and other resolutions may be used. The picture size need not even be a multiple of 16. In this case an MPEG video coder adds padding pels to the right or bottom edges of a picture in order to bring the transmitted resolution up to a multiple of 16, and the decoder discards these after decoding the picture. For example, a horizontal resolution of 360 pels could be coded by adding 8 padding pels to the right edge of each horizontal row bringing the total up to 368 pels. 23 macroblocks would be coded in each row. The decoder would discard the extra padding pels after decoding, giving a final decoded horizontal resolution of 360 pels.

D.3.2 Conversion from film

If film material can be digitized at 24 pictures/s, then it forms an excellent source for an ISO/IEC 11172-2 bitstream. It may be digitized at the desired spatial resolution. The picture_rate field in the video sequence header, see 2.4.2.3, allows the picture rate of 24 pictures/s to be specified exactly.

Sometimes the source material available for compression consists of film material which has been converted to video at some other rate. The encoder may detect this and recode at the original film rate. For example, 24 pictures/s film material may have been digitized and converted to a 30 frame/s system by the technique of 3:2 pulldown. In this mode digitized pictures are shown alternately for 3 and for 2 television field times. This alternation may not be exact since the actual frame rate might be 29,97 frames/s and not the 30 frames/s that the 3:2 pulldown technique gives. In addition the pulldown timing might have been changed by editing and splicing after the conversion. A sophisticated encoder might detect the duplicated fields, average them to reduce digitization noise, and code the result at the original 24 pictures/s rate. This should give a significant improvement in quality over coding at 30 pictures per second, since direct coding at 30 pictures/s destroys the 3:2 pulldown timing and gives a jerky appearance to the final decoded video.

D.4 Model decoder

D.4.1 Need for a decoder model

A coded bitstream contains different types of pictures, and each type ideally requires a different number of bits to encode. In addition, the video may vary in complexity with time, and an encoder may wish to devote more coding bits to one part of a sequence than to another. For constant bitrate coding, varying the number of bits allocated to each picture requires that the decoder have buffering to store the bits not needed to decode the immediate picture. The extent to which an encoder can vary the number of bits allocated to each picture depends on the amount of this buffering. If the amount of the buffering is large an encoder can use greater variations, increasing the picture quality, but at the cost of increasing the decoding delay. Encoders need to know the size of the amount of the decoder's buffering in order to determine to what extent they can vary the distribution of coding bits among the pictures in the sequence.

The model decoder is defined to solve two problems. It constrains the variability in the number of bits that may be allocated to different pictures and it allows a decoder to initialize its buffering when the system is started. It should be noted that Part 1 of this International Standard addresses the initialisation of buffers and the maintenance of synchronisation during playback in the case when two or more elementary streams (for example one audio and one video stream) are multiplexed together. The tools defined in ISO/IEC 11172-1 for the maintenance of synchronisation should be used by decoders when multiplexed streams are being played.

D.4.2 Decoder model

Annex C contains the definition of a parameterized model decoder for this purpose. It is known as a Video Buffer Verifier (VBV). The parameters used by a particular encoder are defined in the bitstream. This really defines a model decoder that is needed if encoders are to be assured that the coded bitstreams they produce will be decodable. The model decoder looks like this:

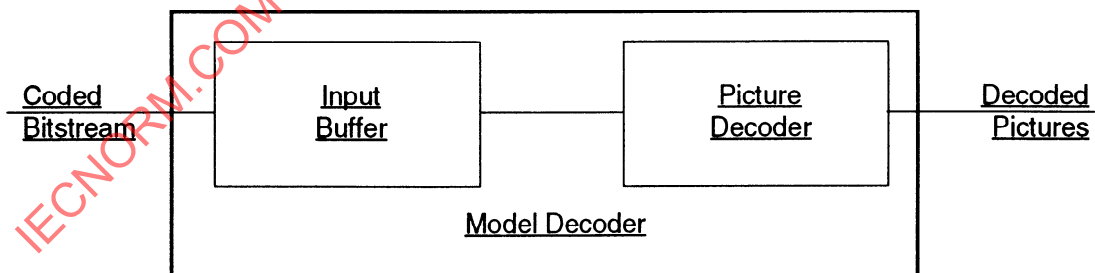


Figure D.14 -- Model decoder

A fixed-rate channel is assumed to put bits at a constant rate into the Input Buffer. At regular intervals, set by the picture rate, the Picture Decoder instantaneously removes all the bits for the next picture from the Input Buffer. If there are too few bits in the Input Buffer, i.e. all the bits for the next picture have not been received, then the Input Buffer underflows and there is an underflow error. If, during the time between picture starts, the capacity of the Input Buffer is exceeded, then there is an overflow error.

Practical decoders differ from this model in several important ways. They may implement their buffering at a different point in the decoder, or distribute it throughout the decoder. They may not remove all the bits required to decode a picture from the Input Buffer instantaneously, they may not be able to control the start

of decoding very precisely as required by the buffer fullness parameter in the picture header, and they take a finite time to decode. They may also be able to delay decoding for a short time to reduce the chances of underflow occurring. But these differences depend in degree and kind on the exact method of implementation. To satisfy requirements of different implementations, the MPEG video committee (ISO/IEC JTC1 SC29/WG11) chose a very simple model for the decoder. Practical implementations of decoders must ensure that they can decode the bitstream constrained by this model. In many cases this will be achieved by using an Input Buffer that is larger than the minimum required, and by using a decoding delay that is larger than the value derived from the `vbv_delay` parameter. The designer must compensate for differences between the actual design and the model in order to guarantee that the decoder can handle any bitstream that satisfies the model.

Encoders monitor the status of the model to control the encoder so that overflow problems do not occur. The calculated buffer fullness is transmitted at the start of each picture so that the decoder can maintain synchronization.

D.4.3 Buffer size and delay

For constant bit rate operation each picture header contains a `vbv_delay` parameter to enable decoders to start their decoding correctly. This parameter defines the time needed to fill the Input Buffer of figure D.14 from an empty state to the correct level immediately before the Picture Decoder removes all the bits for the picture. This time is thus a delay and is measured in units of 1/90 000 s. This number was chosen because it is almost an exact multiple of the picture durations: 1/24, 1/25, 1/29,97 and 1/30, and because it is comparable in duration to an audio sample.

The delay is given by

$$D = \text{vbv_delay} / 90\,000 \text{ s}$$

For example, if `vbv_delay` were 9 000, then the delay would be 0,1 sec. This means that at the start of a picture the Input Buffer of the model decoder should contain exactly 0,1 s worth of data from the input bitstream.

The bit rate, *R*, is defined in the sequence header. The number of bits in the Input Buffer at the beginning of the picture is thus given by:

$$B = D * R = \text{vbv_delay} * R / 90\,000 \text{ bits}$$

For example, if `vbv_delay` were 9 000 and *R* were 1,2 Mbits/s, then the number of bits in the Input Buffer would be 120 000.

The constrained parameter bitstream requires that the Input Buffer have a capacity of 327 680 bits, and *B* should never exceed this value.

D.5 MPEG video bitstream syntax

This clause describes the video bitstream in a top-down fashion. A sequence is the top level of video coding. It begins with a sequence header which defines important parameters needed by the decoder. The sequence header is followed by one or more groups of pictures. Groups of pictures, as the name suggests, consist of one or more individual pictures. The sequence may contain additional sequence headers. A sequence is terminated by a `sequence_end_code`. ISO/IEC 11172-1 allows considerable flexibility in specifying application parameters such as bit rate, picture rate, picture resolution, and picture aspect ratio. These parameters are specified in the sequence header.

If these parameters, and some others, fall within certain limits, then the bitstream is called a constrained parameter bitstream.

D.5.1 Sequence

A video sequence commences with a sequence header and is followed by one or more groups of pictures and is ended by a `sequence_end_code`. Additional sequence headers may appear within the sequence. In each such repeated sequence header, all of the data elements with the permitted exception of those defining quantization matrices (`load_intra_quantizer_matrix`, `load_non_intra_quantizer_matrix` and optionally

`intra_quantizer_matrix` and `non_intra_quantizer_matrix`) shall have the same values as the first sequence header. Repeating the sequence header with its data elements makes random access into the video sequence possible. The quantization matrices may be redefined as required with each repeated sequence header.

The encoder may set such parameters as the picture size and aspect ratio in the sequence header, to define the resources that a decoder requires. In addition, user data may be included.

D.5.1.1 Sequence header code

A coded sequence begins with a sequence header and the header starts with the sequence start code. Its value is:

hex: 00 00 01 B3
binary: 0000 0000 0000 0000 0000 0001 1011 0011

This is a unique string of 32 bits that cannot be emulated anywhere else in the bitstream, and is byte-aligned, as are all start codes. To achieve byte alignment the encoder may precede the sequence start code with any number of zero bits. These can have a secondary function of preventing decoder input buffer underflow. This procedure is called *bit stuffing*, and may be performed before any start code. The stuffing bits must all be zero. The decoder discards all such stuffing bits.

The sequence start code, like all video start codes, begins with a string of 23 zeros. The coding scheme ensures that such a string of consecutive zeros cannot be produced by any other combination of codes, i.e. it cannot be emulated by other codes in the video bitstream. This string of zeros can only be produced by a start code, or by stuffing bits preceding a start code.

D.5.1.2 Horizontal size

This is a 12-bit number representing the width of the picture in pels, i.e. the horizontal resolution. It is an unsigned integer with the most significant bit first. A value of zero is not allowed (to avoid start code emulation) so the legal range is from 1 to 4 095. In practice values are usually a multiple of 16. At 1,5 Mbits/s, a popular horizontal resolution is 352 pels. The value 352 is derived from half the CCIR 601 horizontal resolution of 720, rounded down to the nearest multiple of 16 pels. Otherwise the encoder must fill out the picture on the right to the next higher multiple of 16 so that the last few pels can be coded in a macroblock. The decoder should discard these extra pels before display.

For efficient coding of the extra pels, the encoder should add pel values that reduce the number of bits generated in the transformed block. Replicating the last column of pels is usually superior to filling in the remaining pels with a gray level.

D.5.1.3 Vertical size

This is a 12-bit number representing the height of the picture in pels, i.e. the vertical resolution. It is an unsigned integer with the most significant bit first. A value of zero is not allowed (to avoid start code emulation) so the legal range is from 1 to 4 095. In practice values are usually a multiple of 16. Note that the maximum value of `slice_vertical_position` is 175 (decimal), which corresponds to a picture height of 2 800 lines. At 1,5 Mbits/s, a popular vertical resolution is 240 to 288 pels. Values of 240 pels are convenient for interfacing to 525-line NTSC systems, and values of 288 pels are more appropriate for 625-line PAL and SECAM systems.

If the vertical resolution is not a multiple of 16 lines, the encoder must fill out the picture at the bottom to the next higher multiple of 16 so that the last few lines can be coded in a macroblock. The decoder should discard these extra lines before display.

For efficient coding, replicating the last line of pels is usually better than filling in the remaining pels with a gray level.

D.5.1.4 Pel aspect ratio

This is a four-bit number which defines the shape of the pel on the viewing screen. This is needed since the horizontal and vertical picture sizes by themselves do not specify the shape of the displayed picture.

The pel aspect ratio does not give the shape directly, but is an index to the following look up table:

Table D.2 -- Pel aspect ratio

CODE	HEIGHT/WIDTH	COMMENT
0000	undefined	Forbidden
0001	1,0	square pels
0010	0,6735	16:9 625-line
0011	0,7031	
0100	0,7615	
0101	0,8055	
0110	0,8437	16:9 525-line
0111	0,8935	
1000	0,9157	
1001	0,9815	
1010	1,0255	702x575 at 4:3 = 0,9157
1011	1,0695	
1100	1,0950	
1101	1,1575	
1110	1,2015	711x487 at 4:3 = 1,0950
1111	undefined	
		reserved

The code 0000 is forbidden to avoid start code emulation. The code 0001 has square pels. This is appropriate for many computer graphics systems. The code 1000 is suitable for displaying pictures on the 625-line 50Hz TV system (see CCIR Recommendation 601).

$$\text{height / width} = 0,75 * 702 / 575 = 0,9157$$

The code 1100 is suitable for displaying pictures on the 525-line 60Hz TV system (see CCIR Recommendation 601).

$$\text{height / width} = 0,75 * 711 / 487 = 1,0950$$

The code 1111 is reserved for possible future extensions to this part of ISO/IEC 11172.

The remaining points in the table were filled in by interpolating between these two points 1000 and 1100 using the formula:

$$\text{aspect ratio} = 0,5855 + 0,044N$$

where N is the value of the code in table D.2. These additional pel aspect ratios might be useful for HDTV where ratios of 16:9 and 5:3 have been proposed.

It is evident that the specification does not allow all possible pel aspect ratios to be specified. We therefore presume that a certain degree of tolerance is allowable. Encoders will convert the actual pel aspect ratio to the nearest value in the table, and decoders will display the decoded values to the nearest pel aspect ratio of which they are capable.

D.5.1.5 Picture rate

This is a four-bit integer which is an index to the following table:

Table D.3 -- Picture rate

CODE	PICTURES PER SECOND
0000	Forbidden
0001	23,976
0010	24
0011	25
0100	29,97
0101	30
0110	50
0111	59,94
1000	60
1001	Reserved
.	.
1111	Reserved

The allowed picture rates are commonly available sources of analog or digital sequences. One advantage in not allowing greater flexibility in picture rates is that standard techniques may be used to convert to the display rate of the decoder if it does not match the coded rate.

D.5.1.6 Bit rate

The bit rate is an 18-bit integer giving the bit rate of the data channel in units of 400 bits/s. The bit rate is assumed to be constant for the entire sequence. The actual bit rate is rounded up to the nearest multiple of 400 bits/s. For example, a bit rate of 830 100 bits/s would be rounded up to 830 400 bits/s giving a coded bit rate of 2 076 units.

If all 18 bits are 1 then the bitstream is intended for variable bit rate operation. The value zero is forbidden.

For constant bit rate operation, the bit rate is used by the decoder in conjunction with the `vbv_delay` parameter in the picture header to maintain synchronization of the decoder with a constant rate data channel. If the stream is multiplexed using ISO/IEC 11172-1, the time-stamps and system clock reference information defined in ISO/IEC 11172-1 provide a more appropriate tool for performing this function.

D.5.1.7 Marker bit

The bit rate is followed by a single reserved bit which is always set to 1. This bit prevents emulation of start codes.

D.5.1.8 VBV buffer size

The buffer size is a 10-bit integer giving the minimum required size of the input buffer in the model decoder in units of 16 384 bits (2 048 bytes). For example, a buffer size of 20 would require an input buffer of $20 \times 16\,384 = 327\,680$ bits (= 40 960 bytes). Decoders may provide more memory than this, but if they provide less they will probably run into buffer overflow problems while the sequence is being decoded.

D.5.1.9 Constrained Parameter flag

If certain parameters specified in the bitstream fall within predefined limits, then the bitstream is called a constrained parameter bitstream. Thus the constrained parameter bitstream is a standard of performance giving guidelines to encoders and decoders to facilitate the exchange of bitstreams.

The bitrate parameter allows values up to about 100 Mbits/s, but a constrained parameter bitstream must have a bit rate of 1,856 Mbits/s or less. Thus the bit rate parameter must be 3 712 or less.

The picture rate parameter allows picture rates up to 60 pictures/s, but a constrained parameter bitstream must have a picture rate of 30 pictures/s or less.

The resolution of the coded picture is also specified in the sequence header. Horizontal resolutions up to 4 095 pels are allowed by the syntax, but in a constrained parameter bitstream the resolution is limited to 768 pels or less. Vertical resolutions up to 4 095 pels are allowed, but that in a constrained parameter

bitstream is limited to 576 pels or less. In a constrained parameter bitstream, the total number of macroblocks per picture is limited to 396. This sets a limit on the maximum area of the picture which is only about one quarter of the area of a 720x576 pel picture. In a constrained parameter bitstream, the pel rate is limited to 2 534 400 pels/s. For a given picture rate, this sets another limit on the maximum area of the picture. If the picture has the maximum area of 396 macroblocks, then the picture rate is restricted to 25 pictures/s or less. If the picture rate has the maximum constrained value of 30 pictures/s the maximum area is limited to 330 macroblocks.

A constrained parameter bitstream can be decoded by a model decoder with a buffer size of 327 680 bits without overflowing or underflowing during the decoding process. The maximum buffer size that can be specified for a constrained parameter bitstream is 20 units.

A constrained parameter bitstream uses a forward_f_code or backward_f_code less than or equal to 4. This constrains the maximum range of motion vectors that can be represented in the bitstream (see table D.7).

If all these conditions are met, then the bitstream is constrained and the constrained_parameters_flag in the sequence header should be set to 1. If any parameter is exceeded, the flag shall be set to 0 to inform decoders that more than a minimum capability is required to decode the sequence.

D.5.1.10 Load intra quantizer matrix

This is a one-bit flag. If it is set to 1, sixty-four 8-bit integers follow. These define an 8 by 8 set of weights which are used to quantize the DCT coefficients. They are transmitted in the zigzag scan order shown in figure D.30. None of these weights can be zero. The first weight must be eight which matches the fixed quantization level of the dc coefficient.

If the flag is set to zero, the intra quantization matrix must be reset to the following default value:

8	16	19	22	26	27	29	34
16	16	22	24	27	29	34	37
19	22	26	27	29	34	34	38
22	22	26	27	29	34	37	40
22	26	27	29	32	35	40	48
26	27	29	32	35	40	48	58
26	27	29	34	38	46	56	69
27	29	35	38	46	56	69	83

Figure D.15 -- Default intra quantization matrix

The default quantization matrix is based on work performed by ISO/IEC JTC1 SC29/WG10 (JPEG) [6]. Experience has shown that it gives good results over a wide range of video material. For resolutions close to 350x250 there should normally be no need to redefine the intra quantization matrix. If the picture resolution departs significantly from this nominal resolution, then some other matrix may give perceptibly better results.

The weights increase to the right and down. This reflects the human visual system which is less sensitive to quantization noise at higher frequencies.

D.5.1.11 Load non-intra quantizer matrix

This is a one-bit flag. If it is set to 1, sixty-four 8-bit integers follow in zigzag scan order. None of these integers can be zero.

If the flag is set to zero, the non-intra quantization matrix must be reset to the following default value which consists of all 16s.

16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16

Figure D.16 -- Default non-intra quantization matrix

This flat default quantization matrix was adopted from H.261 which uses a flat matrix for the equivalent of P-pictures [5]. Little work has been performed to determine the optimum non-intra matrix for MPEG video coding, but evidence suggests that it is more dependent on video material than is the intra matrix. The optimum non-intra matrix may be somewhere between the flat default non-intra matrix and the strongly frequency-dependent values of the default intra matrix.

D.5.1.12 Extension data

This start code is byte-aligned and is 32 bits long. Its value is

hex: 00 00 01 B5
binary: 0000 0000 0000 0000 0000 0001 1011 0101

It may be preceded by any number of zeros. If it is present then it will be followed by an undetermined number of data bytes terminated by the next start code. These data bytes are reserved for future extensions to this part of ISO/IEC 11172, and should not be generated by encoders. MPEG video decoders should have the capability to discard any extension data found.

D.5.1.13 User data

A user data start code may follow the optional extension data. This start code is byte-aligned and is 32 bits long. Its value is

hex: 00 00 01 B2
binary: 0000 0000 0000 0000 0000 0001 1011 0010

It may be preceded by any number of zeros. If it is present then it will be followed by an undetermined number of data bytes terminated by the next start code. These data bytes can be used by the encoder for any purpose. The only restriction on the data is that they cannot emulate a start code, even if not byte-aligned. This means that a string of 23 consecutive zeros must not occur. One way to prevent emulation is to force the most significant bit of alternate bytes to be a 1.

In closed encoder-decoder systems the decoder may be able to use the data. In the more general case, decoders should be capable of discarding the user data.

D.5.2 Group of pictures

Two distinct picture orderings exist, the display order and the bitstream order (as they appear in the video bitstream). A group of pictures (gop) is a set of pictures which are contiguous in display order. A group of pictures must contain at least one I-picture. This required picture may be followed by any number of I and P-pictures. Any number of B-pictures may be interspersed between each pair of I or P-pictures, and may also precede the first I-picture.

Property 1. A group of pictures, in bitstream order, must start with an I-picture and may be followed by any number of I, P or B-pictures in any order.

Property 2. Another property of a group of pictures is that it must begin, in display order, with an I or a B-picture, and must end with an I or a P-picture. The smallest group of pictures consists of a single I-picture, whereas the largest size is unlimited.

The original concept of a group of pictures was a set of pictures that could be coded and displayed independently of any other group. In the final version of this part of ISO/IEC 11172 this is not always true, and any B-pictures preceding (in display order) the first I-picture in a group may require the last picture in the previous group in order to be decoded. Nevertheless encoders can still construct groups of pictures which are independent of one another. One way to do this is to omit any B-pictures preceding the first I-picture. Another way is to allow such B-pictures, but to code them using only backward motion compensation.

Property 3. From a coding point of view, a concisely stated property is that a group of pictures begins with a group of pictures header, and either ends at the next group of pictures header or at the next sequence header or at the end of sequence, whichever comes first.

Some examples of groups of pictures are given below:

```

I
I  P  P
I  B  P  B  P
B  B  I  B  P  B  P
B  B  I  B  B  P  B  B  P  B  B  P
B  I  B  B  B  B  P  B  I  B  B  I  I

```

Figure D.17 -- Examples of groups of pictures in display order

These examples illustrate what is possible, and do not constitute a suggestion for structures of groups of pictures.

Group of pictures start code

The group of pictures header starts with the Group of Pictures start code. This code is byte-aligned and is 32 bits long. Its value is

```

hex:    00 00 01 B8
binary: 0000 0000 0000 0000 0000 0001 1011 1000

```

It may be preceded by any number of zeros. The encoder may have inserted some zeros to get byte alignment, and may have inserted additional zeros to prevent buffer underflow. An editor may have inserted zeros in order to match the `vbv_delay` parameter of the first picture in the group.

Time code

A time code of 25 bits immediately follows the group of pictures start code. This encodes the same information as the SMPTE time code [4].

The time code can be broken down into six fields as shown in table D.4.

Table D.4 -- Time code fields

FIELD	BITS	VALUES
Drop frame flag	1	
Hours	5	0 to 23
Minutes	6	0 to 59
Fixed	1	1
Seconds	6	0 to 59
Picture number	6	0 to 60

The time code refers to the first picture in the group in display order, i.e. the first picture with a temporal reference of zero. The SMPTE time code is included to provide a video time identification to applications. It may be discontinuous. The presentation time-stamp in the System layer (Part 1) has a much higher precision and identifies the time of presentation of the picture.

Closed GOP

A one bit flag follows the time code. It denotes whether the group of pictures is open or closed. Closed groups can be decoded without using decoded pictures of the previous group for motion compensation, whereas open groups require such pictures to be available.

A typical example of a closed group is shown in figure D.18a.

I	B	B	P	B	B	P	B	B	P	B	B	P
0	1	2	3	4	5	6	7	8	9	10	11	12

(a) closed group

B	B	I	B	B	P	B	B	P	B	B	P	B	B	P
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

(b) open or closed group

Figure D.18 -- Example groups of pictures in display order

A less typical example of a closed group is shown in figure D.18b. In this example, the B-pictures which precede the first I-picture must use backward motion compensation only, i.e. any motion compensation must be based only on picture number 2 in the group.

If the closed_gop flag is set to 0 then the group is open. The first B-pictures that precede the first I-picture in the group may have been encoded using the last picture in the previous group for motion compensation.

Broken link

A one bit flag follows the closed_gop flag. It denotes whether the B-pictures which precede the first I-picture in the GOP can be correctly decoded. If it is set to 1, these pictures cannot be correctly decoded because the I-picture or P-picture from the previous group pictures that is required to form the predictions is not available (presumably because the preceding group of pictures has been removed by editing). The decoder will probably choose not to display these B-pictures.

If the sequence is edited so that the original group of pictures no longer precedes the current group of pictures then this flag normally will be set to 1 by the editor. However, if the closed_gop flag for the current group of pictures is set, then the editor should not set the broken_link flag. Because the group of pictures is closed, the first B-pictures (if any) can still be decoded correctly.

Extension data

This start code is byte-aligned and is 32 bits long. Its value is

hex: 00 00 01 B5
binary: 0000 0000 0000 0000 0000 0001 1011 0101

It may be preceded by any number of zeros. If it is present then it will be followed by an undetermined number of data bytes terminated by the next start code. These data bytes are reserved for future extensions to this part of ISO/IEC 11172, and should not be generated by encoders. MPEG video decoders should have the capability to discard any extension data found.

User data

A user data start code may follow the optional extension data. This start code is byte-aligned and is 32 bits long. Its value is

hex: 00 00 01 B2
binary: 0000 0000 0000 0000 0000 0001 1011 0010

It may be preceded by any number of zeros. If it is present then it will be followed by an undetermined number of data bytes terminated by the next start code. These data bytes can be used by the encoder for any purpose. The only restriction on the data is that they cannot emulate a start code, even if not byte-aligned.

This means that a string of 23 consecutive zeros must not occur. One way to prevent emulation is to force the most significant bit of alternate bytes to be a 1.

In closed encoder-decoder systems the decoder may be able to use the data. In the more general case, decoders should be capable of discarding the user data.

D.5.3 Picture

The picture layer contains all the coded information for one picture. The header identifies the temporal reference of the picture, the picture coding type, the delay in the video buffer verifier (V BV) and, if appropriate, the range of motion vectors used.

D.5.3.1 Picture header and start code

A *picture* begins with a picture header. The header starts with a picture start code. This code is byte-aligned and is 32 bits long. Its value is:

hex: 00 00 01 00
binary: 0000 0000 0000 0000 0000 0001 0000 0000

It may be preceded by any number of zeros.

D.5.3.2 Temporal reference

The Temporal Reference is a ten-bit number which can be used to define the order in which the pictures must be displayed. It may be useful since pictures are not transmitted in display order, but rather in the order which the decoder needs to decode them. The first picture, in display order, in each group must have Temporal Reference equal to zero. This is incremented by one for each picture in the group.

Some example groups of pictures with their Temporal Reference numbers are given below:

Example (a) in display order	I	B	P	B	P								
	0	1	2	3	4								
Example (a) in decoding order	I	P	B	P	B								
	0	2	1	4	3								
Example (b) in display order	B	B	I	B	B	P	B	B	P	B	B	P	
	0	1	2	3	4	5	6	7	8	9	10	11	
Example (b) in coded order	I	B	B	P	B	B	P	B	B	P	B	B	
	2	0	1	5	3	4	8	6	7	11	9	10	
Example (c) in display order	B	I	B	B	B	B	P	B	I	B	B	I	I
	0	1	2	3	4	5	6	7	8	9	10	11	12
Example (c) in coded order	I	B	P	B	B	B	B	I	B	I	B	B	I
	1	0	6	2	3	4	5	8	7	11	9	10	12

Figure D.19 -- Examples of groups of pictures and temporal references

If there are more than 1024 pictures in a group, then the Temporal Reference is reset to zero and then increments anew. This is illustrated below:

B	B	I	B	B	P	...	P	B	B	P	...	P	B	B	P	display order
0	1	2	3	4	5	...	1 022	1 023	0	1	...	472	473	474	475	

Figure D.20 -- Example group of pictures containing 1 500 pictures

D.5.3.3 Picture coding type

A three bit number follows the temporal reference. This is an index into the following table defining the type of picture.

Table D.5 -- Picture types

CODE	PICTURE TYPE
000	Forbidden
001	I-picture
010	P-picture
011	B-picture
100	D Picture
101	Reserved
110	Reserved
111	Reserved

The various types of pictures are described in D.2.3. Codes 101 through 111 are reserved for future extensions to this part of ISO/IEC 11172. Decoders should be capable of discarding all pictures of this type, and scan for the next picture start code, group start code or sequence start code. Code 000 will never be used to avoid start code emulation.

D.5.3.4 VBV delay

For constant bit rate operation, vbv_delay can be used at the start of decoding and after a random access to ensure that the correct number of bits have been read by the decoder before the first picture is displayed.

The buffer fullness is not specified in bits but rather in units of time. The vbv_delay is a 16-bit number defining the time needed in units of 1/90 000 s to fill the input buffer of the model decoder from an empty state to the correct state at the bit rate specified in the sequence header.

For example, suppose the vbv_delay had a decimal value of 30000, then the time delay would be:

$$D = 30\,000 / 90\,000 = 1/3 \text{ s}$$

If the channel bit rate were 1,2 Mbits/s then the contents of the buffer before the picture is decoded would be:

$$B = 1\,200\,000 / 3 = 400\,000 \text{ bits}$$

If the decoder determined that its actual buffer fullness differed significantly from this value, then it would have to adopt some strategy for regaining synchronization.

The meaning of vbv_delay is undefined for variable bit rate operation.

D.5.3.5 Full pel forward vector

This is a one bit flag giving the precision of the forward motion vectors. If it is 1 then the precision of the vectors is in integer pels, if it is zero then the precision is half a pel. Thus if the flag is set to one the vectors have twice the range than they do if the flag set to zero.

This flag is present only in the headers of P-pictures and B-pictures. It is absent in I-pictures and D pictures.

D.5.3.6 Forward f-code

This is a three-bit number and, like the full pel forward vector flag, is present only in the headers of P-pictures and B-pictures. It provides information used for decoding the coded forward vectors and controls the maximum size of the forward vectors that can be coded. It can take only values of 1 through 7; a value of zero is forbidden.

Two parameters used in decoding the forward motion vectors are derived from `forward_f_code`, `forward_r_size` and `forward_f`.

The `forward_r_size` is one less than the `forward_f_code` and so can take values 0 through 6.

The `forward_f` parameter is given by table D.6:

Table D.6 -- f_codes

forward/backward_f_code	forward/backward_f
1	1
2	2
3	4
4	8
5	16
6	32
7	64

D.5.3.7 Full pel backward vector

This is a one bit flag giving the precision of the backward motion vectors. If it is 1 then the precision of the vectors is in integer pels, if it is zero then the precision is half a pel. Thus if the flag is set to one the vectors have twice the range than they do if the flag set to zero.

This flag is only present in the headers of B-pictures. It is absent in I-pictures, P-pictures and D pictures.

D.5.3.8 Backward f-code

This is a three-bit number and, like the full pel backward vector flag, is present only in the headers of B-pictures. It provides information used for decoding the coded backward vectors. It can take only values of 1 through 7; a value of zero is forbidden.

The `backward_f` parameter is derived from the `backward_f_code` and is given by table D.6

D.5.3.9 Extra picture information

Extra picture information is the next field in the picture header. Any number of information bytes may be present. An information byte is preceded by a flag bit which is set to 1. Information bytes are therefore generally not byte-aligned. The last information byte is followed by a zero bit. The smallest size of this field is therefore one bit, a 0, that has no information bytes. The largest size is unlimited. The following example has 16 bits of extra information denoted by E:

1 E E E E E E E E E E E E E E E E 0

Where E is an extra information bit.

The extra information bytes are reserved for future extensions to this part of ISO/IEC 11172. The meaning of these bytes is currently undefined, so encoders must not generate such bytes and decoders must be capable of discarding them.

D.5.3.10 Extension data

This start code is byte-aligned and is 32 bits long. Its value is:

hex: 00 00 01 B5
binary: 0000 0000 0000 0000 0000 0001 1011 0101

It may be preceded by any number of zeros. If it is present then it will be followed by an undetermined number of data bytes terminated by the next start code. These data bytes are reserved for future extensions to this part of ISO/IEC 11172, and should not be generated by encoders. MPEG video decoders must be capable of discarding them.

D.5.3.11 User data

This start code is byte-aligned and is 32 bits long. Its value is

hex: 00 00 01 B2
binary: 0000 0000 0000 0000 0000 0001 1011 0010

It may be preceded by any number of zeros. If it is present then it will be followed by an undetermined number of data bytes terminated by the next start code. These data bytes can be used by the encoder for any purpose. The only restriction on the data is that they cannot emulate a start code, even if not byte-aligned. One way to prevent emulation is to force the most significant bit of alternate bytes to be a 1.

In closed encoder-decoder systems the decoder may be able to use the data. In the more general case, decoders should be capable of discarding the user data.

D.5.4 Slice

Pictures are divided into slices. Each slice consists of an integral number of macroblocks in raster scan order. Slices can be of different sizes within a picture, and the division in one picture need not be the same as the division in any other picture. Slices can begin and end at any macroblock in a picture subject to the following restrictions. The first slice must begin at the top left of the picture, and the end of the last slice must be the bottom right macroblock of the picture. There can be no gaps between slices, nor can slices overlap. The minimum number of slices in a picture is one, the maximum number is equal to the number of macroblocks.

Each slice starts with a slice start code, the exact value of which defines the vertical position of the slice. This is followed by a code that sets the quantization step-size. At the start of each slice the predictors for the dc coefficient values and the predictors for the vector decoding are all reset. The horizontal position of the start of the slice is given by the macroblock address of the first macroblock in the slice. The result of all this is that, within a picture, a slice can be decoded without information from the previous slices. Therefore, if a data error occurs, decoding can begin again at the subsequent slice.

If the data are to be used in an error free environment, then one slice per picture may be appropriate. If the environment is noisy, then one slice per row of macroblocks may be more desirable, as shown in figure D.21.

1 begin	end 1
2 begin	end 2
3 begin	end 3
4 begin	end 4
5 begin	end 5
6 begin	end 6
7 begin	end 7
8 begin	end 8
9 begin	end 9
10 begin	end 10
11 begin	end 11
12 begin	end 12
13 begin	end 13

Figure D.21 -- Possible arrangement of slices in a 256x192 picture

In this figure and in the next, each strip is one macroblock high, i.e. 16 pels high.

Since each slice header requires 40 bits, there is some penalty for including more than the minimum number of slices. For example, a sequence with a vertical resolution of 240 lines coded at 30 pictures/s requires approximately $40 \times 30 = 1\,200$ bits/s for the slice headers using one slice per picture, and $40 \times 15 \times 30 = 18\,000$ bits/s with one slice per row, an additional overhead of 16 800 bits/s. The calculation is approximate and underestimates the impact, since the inclusion of a slice imposes additional requirements that the macroblock immediately before the slice header be coded, as well as the first macroblock in the slice.

The coding structure permits great flexibility in dividing a picture up into slices. One possible arrangement is shown in figure D.22.

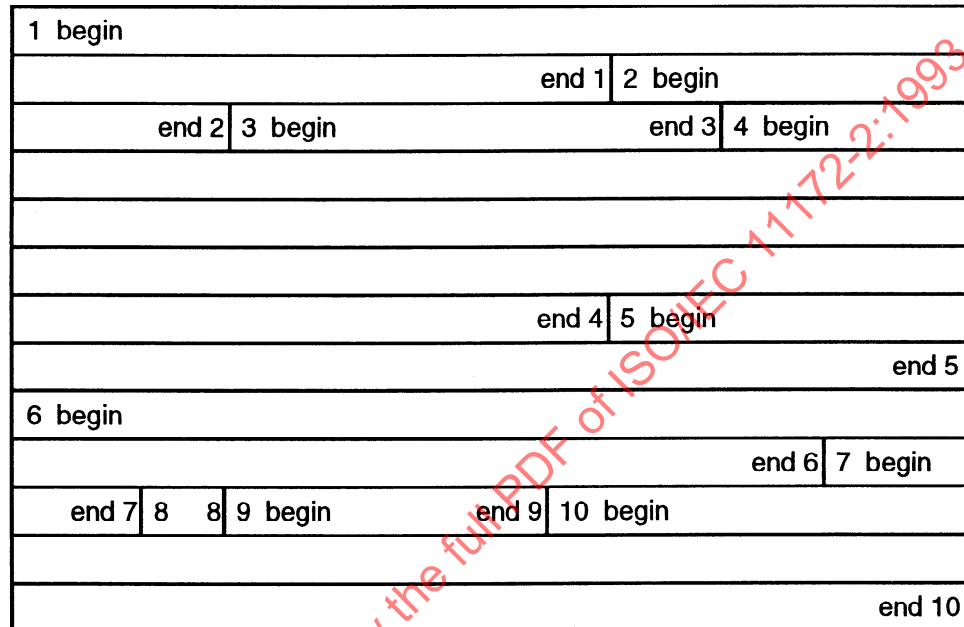


Figure D.22 -- Possible arrangement of slices in a 256x192 picture

This division into slices is given for illustrative purposes only. It is not intended as a suggestion on how to divide a picture into slices.

D.5.4.1 Slice header and start code

Slices start with a slice header. Each slice header starts with a slice start code. This code is byte-aligned and is 32 bits long. The last eight bits can take on a range of values which define the vertical position of the slice in the picture. The permitted slice start codes are:

hex:	from	00 00 01 01
	to	00 00 01 AF
binary:	from	0000 0000 0000 0000 0000 0001 0000 0001
	to	0000 0000 0000 0000 0000 0001 1010 1111

Each slice start code may be preceded by any number of zeros.

The last 8 bits of the slice start code give the slice vertical position, i.e. the vertical position of the first macroblock in the slice in units of macroblocks starting with position 1 at the top of the picture. A useful variable is macroblock row. This is similar to slice vertical position except that row 0 is at the top of the picture. Thus

$$\text{slice vertical position} = \text{macroblock row} + 1$$

For example, a slice start code of 00000101 hex means that the first macroblock in the slice is at vertical position 1 or macroblock row 0, i.e. at the top of the picture. A slice start code of 00000120 hex means

that the first macroblock is at vertical position 32 or macroblock row 31, i.e. at the 496th row of pels. It is possible for two or more slices to have the same vertical position.

The maximum vertical position is 175 units. A slice with this position would require a vertical size of $175 \times 16 = 2\,800$ pels.

The horizontal position of the first macroblock in the slice can be calculated from its macroblock address increment. Thus its position in the picture can be determined without referring to any previous slice or macroblock. Thus a decoder may decode any slice in a picture without having decoded any other slice in the same picture. This feature allows decoders to recover from bit errors by searching for the next slice start code and then resuming decoding.

D.5.4.2 Quantizer scale

The quantizer scale is a five-bit integer which is used by the decoder to calculate the DCT coefficients from the transmitted quantized coefficients. A value of 0 is forbidden, so the quantizer scale can have any value between 1 and 31 inclusive.

Note in addition that the quantizer scale may be set at any macroblock.

D.5.4.3 Extra slice information

Extra slice information forms the last field in the slice header. Any number of information bytes may be present. An information byte is preceded by a flag bit which is set to 1. Information bytes are therefore generally not byte-aligned. The last information byte is followed by a zero bit. The smallest size of this field is therefore one bit, a 0, that has no information bytes. The largest size is unlimited. The following example has 24 bits of extra information denoted by E:

1 E E E E E E E E E E 1 E E E E E E E E E E 1 E E E E E E E E E E 0

The extra information bytes are reserved for future extensions to this part of ISO/IEC 11172. The meaning of these bytes is currently undefined, so encoders must not generate such bytes and decoders must discard them.

The slice header is followed by code defining the macroblocks in the slice.

D.5.5 Macroblock

Slices are divided into macroblocks of 16×16 pels. Macroblocks are coded with a header that contains information on the macroblock address, macroblock type, and the optional quantizer scale. The header is followed by data defining each of the six blocks in the macroblock. It is convenient to discuss the macroblock header fields in the order in which they are coded.

D.5.5.1 Macroblock stuffing

The first field in the macroblock header is "macroblock stuffing". This is an optional field, and may be inserted or omitted at the discretion of the encoder. If present it consists of any number of 11-bit strings with the pattern "0000 0001 111". This stuffing code is used by the encoder to prevent underflow, and is discarded by the decoder. If the encoder determines that underflow is about to occur, then it can insert as many stuffing codes into the first field of the macroblock header it likes.

Note that an encoder has other strategies to prevent buffer underflow. It can insert stuffing bits immediately before a start code. It can reduce the quantizer scale to increase the number of coded coefficients. It can even start a new slice.

D.5.5.2 Macroblock address increment and macroblock escape

Macroblocks have an address which is the number of the macroblock in raster scan order. The top left macroblock in a picture has address 0, the next one to the right has address 1 and so on. If there are M macroblocks in a picture, then the bottom right macroblock has an address M-1.

The address of a macroblock is indicated by transmitting the difference between the addresses of the current macroblock and the previously coded macroblock. This difference is called the macroblock address increment. In I-pictures, all macroblocks are coded and so the macroblock address increment is nearly always one. There is one exception. At the beginning of each slice the macroblock address is set to that of the right hand macroblock of the previous row. At the beginning of the picture it is set to -1. If a slice does not start at the left edge of the picture, then the macroblock address increment for the first macroblock in the slice will be larger than one. For example, the picture of figure D.22 has 16 macroblocks per row. At the start of slice 2 the macroblock address is set to 15 which is the address of the macroblock at the right hand edge of the top row of macroblocks. If the first slice contained 26 macroblocks, 10 of them would be in the second row, so the address of the first macroblock in slice 2 would be 26 and the macroblock address increment would be 11.

Macroblock address increments are coded using the VLC codes in the table in B.1.

It can be seen that there is no code to indicate a macroblock address increment of zero. This is why the macroblock address is set to -1 rather than zero at the top of a picture. The first macroblock will have an increment of one making its address equal to zero.

The macroblock address increments allow the position of the macroblock within the picture to be determined. For example, assume that a slice header has the start code equal to 00 00 01 0A hex, that the picture width is 256 pels, and that a macroblock address increment code 0000111 is in the macroblock header of the first macroblock in the slice. A picture width of 256 pels implies that there are 16 macroblocks per row in this picture. The slice start code tells us that the slice vertical position is 10, and so the macroblock row is 9. The slice header sets the previous macroblock address to the last macroblock on row 8, which has address 143. The macroblock address increment VLC leads to a macroblock address increment of 8, and so the macroblock address of the first macroblock in the slice is $143 + 8 = 151$.

The macroblock row may be calculated from the address:

$$\begin{aligned}\text{macroblock row} &= \text{macroblock address} / \text{macroblock width} \\ &= 151 / 16 \\ &= 9\end{aligned}$$

The division symbol signifies integer truncation, not rounding.

The macroblock column may also be calculated from the address:

$$\begin{aligned}\text{macroblock column} &= \text{macroblock address} \% \text{macroblock width} \\ &= 151 \% 16 \\ &= 7\end{aligned}$$

Columns are numbered from the left of the picture starting at 0.

There are two special codewords: escape and stuffing.

The escape code means "add 33 to the following macroblock address increment". This allows increments greater than 33 to be coded. For example, an increment of 40 would be coded as escape plus an increment of 7:

0000 0001 0000 0010

An increment of 70 would be coded as two escape codes followed by the code for an increment of 4:

0000 0001 0000 0000 0010 0000 11

The stuffing code is included since the decoder must be able to distinguish it from increment codes. It is used by the encoder to prevent underflow, and is discarded by the decoder.

D.5.5.3 Macroblock types

Each of the picture types I, P, and B, have their own macroblock types. See, respectively, D.6.3, D.6.4, and D.6.5 for the codes and their descriptions.

D.5.5.4 Motion horizontal/vertical forward/backward codes

The interpretation of these codes is explained in D.6.2.3.

D.5.5.5 Motion horizontal/vertical forward/backward R

The interpretation of these codes is explained in D.6.2.3.

D.5.5.6 Coded block pattern

This code describes which blocks within the macroblock are coded and transmitted. The interpretation of this code is explained in D.6.4.2.

D.5.5.7 End of macroblock

This code is used only in D-pictures and is described in D.6.6.

D.5.6 Block

A block is an array of 8 by 8 component pel values, treated as a unit and input to the Discrete Cosine Transform (DCT). Blocks of 8 by 8 pels are transformed into arrays of 8 by 8 DCT coefficients using the two dimensional discrete cosine transform.

D.6 Coding MPEG video

D.6.1 Rate control and adaptive quantization

The encoder must control the bit rate so that the model decoder input buffer neither overflows nor underflows. Since the model decoder removes all the bits associated with a picture from the input buffer instantaneously, it is necessary to control only the total number of bits per picture. The encoder should allocate the total numbers of bits among the various types of pictures so that the perceived quality is suitably balanced. The distribution will vary with the scene content and with the particular distribution of the three picture types (I, P and B-pictures).

Within a picture the encoder should allocate the total number of bits available among the macroblocks to maximize the visual quality of the picture.

One method by which an encoder controls the bit rate is to vary the quantizer scale. This is set in each slice header, and may be set at the beginning of any macroblock, giving the encoder excellent control over the bit rate within a picture.

D.6.1.1 Rate control within a sequence

For a typical coding scheme represented by the following group of pictures in display order:

B B I B B P B B P B B P B B P

it has been found that good results can be obtained by matching the visual quality of the I and P-pictures, and by reducing the code size of the B-pictures to save bits giving a generally lower quality for the B-pictures.

The best allocation of bits among the picture types depends on the scene content. Work of the MPEG video committee suggests that allotting P-pictures about 2-5 times as many bits as B-pictures, and allotting I-pictures up to 3 times as many bits as P-pictures gives good results for typical natural scenes. If there is little motion or change in the video, then a greater proportion of the bits should be allotted to the I-pictures.

If there is a lot of motion or change, then the proportion allotted to I-pictures should be reduced and most of the savings given to the P-pictures.

A reasonable encoder algorithm is to start with the foregoing estimates, then reallocate bits dynamically depending on the nature of the video.

D.6.1.2 Rate control within a picture

If the buffer is heading toward overflow, the quantizer scale should be increased. If this action is not sufficient to prevent an impending overflow then, as a last resort, the encoder could discard high frequency DCT coefficients and transmit only low frequency ones. Although this would probably produce visible artifacts in the decoded video, it would in no way compromise the validity of the coded bitstream.

If the buffer is heading toward underflow, the quantizer scale should be reduced. If this is not sufficient, the encoder can insert macroblock stuffing into the bitstream, or add leading zeros to start codes.

Under normal circumstances, the encoder calculates and monitors the state of the model decoder buffer and changes the quantizer scale to avert both overflow and underflow problems.

One simple algorithm that helps accomplish this is to monitor the buffer fullness. Assume that the bits have been allocated among the various picture types, and that an average quantizer scale for each picture type has been established. The actual buffer fullness at any macroblock in a picture can be calculated and compared with the nominal fullness, i.e. the value that would be obtained if the bits were uniformly distributed among all the macroblocks in the picture. If the buffer fullness is larger than the nominal value, then the quantizer scale should be set higher than the average, whereas if the buffer fullness is smaller than the nominal, the quantizer scale should be set lower than the average.

If the quantizer scale is kept constant over a picture, then, for a given number of coding bits, the total mean square error of the coded picture will tend to be close to the minimum. However, the visual appearance of most pictures can be improved by varying the quantizer scale over the picture, making it smaller in smooth areas of the picture and larger in busy areas. This technique reduces the visibility of blockiness in smooth areas at the expense of increased quantization noise in the busy areas where, however, it is masked by the image detail.

Thus a good algorithm for controlling the bitrate within a picture adjusts the quantizer scale depending on both the calculated buffer fullness and on the local image content. Examples of techniques for rate control and quantization may be found in [7][8].

D.6.1.3 Buffer fullness

To give the best visual quality, the encoder should almost fill the input buffer before instructing the decoder to start decoding.

D.6.2 Motion estimation and compensation

D.6.2.1 Motion compensation

P-pictures use motion compensation to exploit temporal redundancy in the video. Decoders construct a predicted block of pels from pels in a previously transmitted picture. Motion within the pictures (e.g. a pan) usually implies that the pels in the previous picture will be in a different position from the pels in the current block, and the displacement is given by motion vectors encoded in the bitstream. The predicted block is usually a good estimate of the current block, and it is usually more efficient to transmit the motion vector plus the difference between the predicted block and the current block, than to transmit a description of the current block by itself.

Consider the following typical group of pictures.

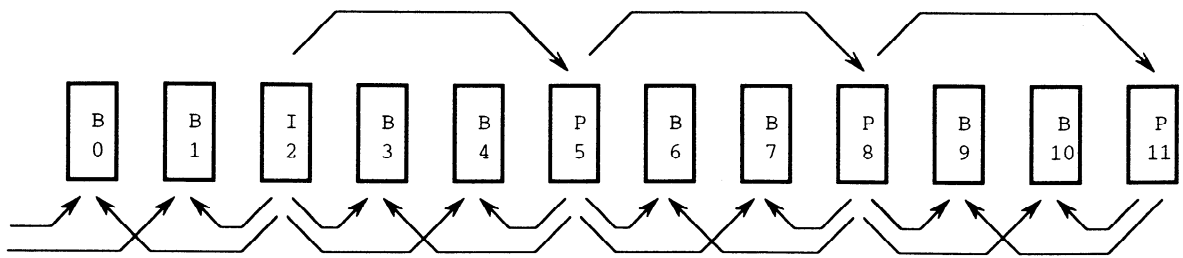


Figure D.23 -- Group of pictures in display order

The I-picture, picture 2, is decoded without requiring any motion vectors. The first P-picture, number 5, is decoded using motion vectors from picture 2. This motion compensation is called forward motion compensation since it is forward in time. Motion vectors define the motion of a macroblock, i.e. the motion of a 16x16 block of luminance pels and the associated chrominance components. Typically, most macroblocks in a P-picture use motion compensation. Non-zero motion vectors are transmitted differentially with reference to the last transmitted motion vector.

The transmitted vectors usually have a precision of half a pel. The maximum range of the vector is set by the `forward_f` parameter in the picture header. Sometimes, if the motion is unusually large, the range may be doubled and the accuracy reduced to integer pels, by the `full_pel_forward_vector` bit in the picture header.

A positive value of the horizontal or vertical component of the motion vector signifies that the prediction is formed from pels in the referenced picture which are spatially to the right or below the pels being predicted.

Not all macroblocks in a P-picture necessarily use motion compensation. Some macroblocks, as defined by the transmitted macroblock type (see table B.2b), may be intra-coded, and these are reconstructed without motion compensation. Full details defining the method of decoding the vectors and constructing the motion-compensated macroblock are given in 2.4.4.2.

P-picture 8 in figure D.23 uses forward motion compensation from picture 5. P-pictures always use forward motion compensation from the last transmitted I or P-picture.

B-pictures may use motion compensation from the previous I or P-picture, from the next (in display order) I or P-picture, or both; i.e., from the last two transmitted I or P-pictures.

Prediction is called forward if reference is made to a picture in the past and called backward if reference is made to a picture in the future. For example, B-picture 3 in figure D.23 uses forward motion compensation from I-picture 2, and backward motion compensation from P-picture 5. B-pictures may use both forward and backward motion compensation and average the result. This operation is called interpolative motion compensation.

All three types of motion compensation are useful, and typically are used in coding B-pictures. Interpolative motion compensation has the advantage of averaging any noise present. Forward or backward motion compensation may be more useful near the edges of pictures, or where a foreground object is passing in front of a fixed or slow moving background.

Note that this technique of coding with P and B-pictures increases the coding efficiency. B-pictures can have greater errors of reconstruction than I or P-pictures to conserve coding bits, but since they are not used as the basis of motion compensation for future pictures, these errors may be tolerated.

D.6.2.2 Motion estimation

Motion compensation in a decoder is straightforward, but motion estimation which includes determining the best motion vectors and which must be performed by the encoder, presents a formidable computational challenge.

Various methods are available to the encoder. The more computationally intensive methods tend to give better results, so there is tradeoff to be made in the encoder: computational power, and hence cost, versus coded video quality.

Using a search strategy the encoder attempts to match the pels in a macroblock with those in a previous or future picture. The vector corresponding to the best match is reported after the search is completed.

D.6.2.2.1 Block matching criteria

In seeking a match, the encoder must decide whether to use the decoded past and future pictures as the reference, or use the original past and future pictures. For motion estimation, use of the decoded pictures by the encoder gives the smallest error in the error picture, whereas use of the original pictures gives the most accurate motion vectors. The choice depends on whether the artifacts of increased noise, or greater spurious motion are judged to be the more objectionable. There is usually little or no difference in quality between the two methods. Note that the decoder does not perform motion estimation. It performs motion compensated prediction and interpolation using vectors calculated in the encoder and stored in the bitstream. In motion compensated prediction and interpolation, both the encoder and decoder must use the decoded pictures as the references.

Several matching criteria are available. The mean square error of the difference between the motion-compensated block and the current block is an obvious choice. Another possible criterion is the mean absolute difference between the motion-compensated block and the current block.

For half pel shifts, the pel values could be interpolated by several methods. Since the decoder uses a simple linear interpolation, there is little reason to use a more complex method in the encoder. The linear interpolation method given in this part of ISO/IEC 11172 is equivalent to the following. Consider four pels having values A, B, D and E as shown in figure D.24:

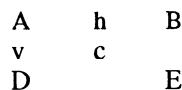


Figure D.24 -- Interpolation of half pel shifts

The value of the horizontally interpolated pel is

$$h = (A + B) // 2$$

where the double division symbol means division with rounding to the nearest integer. Half integer values are to be rounded to the next higher value. Thus if A = 4 and B = 9 then h = 6.5 which is rounded up to 7.

The value of the vertically interpolated pel is

$$v = (A + D) // 2$$

The value of the central interpolated pel is

$$c = (A + B + D + E) // 4$$

D.6.2.2.2 Search range

Once a block matching criterion has been selected, some kind of search strategy must be adopted. This must recognize the limitations of the VLC tables used to code the vectors. The maximum range of the vector depends upon forward_f_code or backward_f_code. The motion vector ranges are given in table D.7.

Table D.7 -- Range of motion vectors

forward_f_code or backward_f_code	Motion vector range	
	full_pel=0	full_pel=1
1	-8 to 7,5	-16 to 15
2	-16 to 15,5	-32 to 31
3	-32 to 31,5	-64 to 63
4	-64 to 63,5	-128 to 127
5	-128 to 127,5	-256 to 255
6	-256 to 255,5	-512 to 511
7	-512 to 511,5	-1 024 to 1 023

The range depends on the value of full_pel_forward_vector or full_pel_backward_vector in the picture header. Thus if all the motion vectors were found to be 15 pels or less, the encoder would usually select half pel accuracy and a forward_f_code or backward_f_code value of 2.

The search must be constrained to take place within the boundaries of the decoded reference picture. Motion vectors which refer to pels outside the picture are not allowed. Any bitstream which refers to such pels does not conform to this part of ISO/IEC 11172.

D.6.2.2.3 2-D search strategy

There are many possible methods of searching another picture for the best match to a current block, and a few simple ones will be described.

The simplest search is a full search. Within the chosen search range, all possible displacements are evaluated using the block matching criterion.

The full search is computationally expensive, and practical encoders may not be able to afford the time required for a full search.

A simple modification of the full search is to search using only integer pel displacements. Once the best integer match has been found, the eight neighbouring half-integer pel displacements are evaluated, and the best one selected as illustrated below:

*	*	*	*	*	y
*	*	*	*	*	y+1
		+	+	+	y+1,5
*	*	+	*	+	y+2
		+	+	+	y+2,5
*	*	*	*	*	y+3
*	*	*	*	*	y+4
x	x+1	x+2	x+3	x+4	

Figure D.25 -- Integer pel and half pel displacements

Assume that the position x+2,y+2 gives the best integer displacement matching using the selected block matching criterion, then the encoder would evaluate the eight positions with half pel displacements marked by + signs in figure D.25. If one of them were a better match then it would become the motion vector, otherwise the motion vector would remain that of the integer displacement x+2,y+2.

If during the integer pel search, two or more positions have the same block matching value, the encoder can adopt a consistent tie-breaking rule.

The modified full search algorithm is approximately an order of magnitude simpler than the full search. Using only integer displacements for the first stage of the search reduces the number of evaluations by a factor of four. In addition, the evaluations are simpler since the pel differences can be calculated directly and do not have to be interpolated.

For some applications even the modified full search may be too time consuming, and a faster search method may be required. One such method is the logarithmic search.

D.6.2.2.4 Logarithmic search

In this search method, grids of 9 displacements are examined, and the search continued based on a smaller grid centered on the position of the best match. If the grids are reduced in size by a factor of 3 at each step then the search is maximally efficient in the sense that any integer shift has a unique selection path to it. This method will find the best match only for a rather limited set of image types. A more robust method is to reduce the size of the grids by a smaller factor at each step, e.g. by a factor of 2. The scaling factors can also be adjusted to match the search ranges of table D.7.

The method will be illustrated with an example. Consider the set of integer shifts in figure D.26:

*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	*	*	1	*	*	*	1	*	*	*	1	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	2	*	2	*	2	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	2	*	1	*	2	*	1	*	*	*	1	*	*	*
*	*	*	*	3	3	3	*	*	*	*	*	*	*	*
*	2	*	2	3	2	3	*	*	*	*	*	*	*	*
*	*	*	*	3	3	3	*	*	*	*	*	*	*	*
*	*	*	1	*	*	*	1	*	*	*	1	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*	*	*	*	*	*

Figure D.26 -- Logarithmic search method for integer pel shifts

The first grid has a spacing of 4 pels. The first step examines pels at shifts of 0, 4, or -4 pels in each direction, marked 1 in figure D.26. The best position is used as the center point of the second grid. Assume it is the pel marked 1 directly to the left of the center pel. The second grid has a spacing of 2 pels. The second step examines pels at shifts of 0, 2, or -2 pels in each direction from the center of the new grid, marked 2 in the figure. The best position is used as the center point of the third grid, assume it is the lower right pel of the second grid. The third grid has a spacing of 1 pel. The third step examines pels at shifts of 0, 1, or -1 pels in each direction from the center of the grid. The best position is used as the center point of the fourth grid. The fourth grid has a spacing of 1/2 pel. The fourth step examines pels at shifts of 0, 1/2, or -1/2 pels in each direction from the center of the grid using the same method as in the modified full search. The best position determines the motion vector.

Some possible grid spacings for various search ranges are given in table D.8.

Table D.8 -- Grid spacings for logarithmic searches

forward f_code	RANGE	STEPS	GRID SPACINGS
1	$\pm 7,5$	4	4 2 1 1/2
2	$\pm 15,5$	5	8 4 2 1 1/2
3	$\pm 31,5$	6	16 8 4 2 1 1/2

For P-pictures only forward searches are performed, but B-pictures require both forward and backward searches. Not all the vectors calculated during the search are necessarily used. In B-pictures either forward or backward motion compensation might be used instead of interpolated motion compensation, and in both P and B-pictures the encoder might decide that a block is better coded as intra, in which case no vectors are transmitted.

D.6.2.2.5 Telescopic search

Even with the faster methods of the modified full search, or the logarithmic search, the search might be quite expensive. For example, if the encoder decides to use a maximum search range of 7 pels per picture

interval, and if there are 4 B-pictures preceding a P-picture, then the full search range for the P-picture would be 35 pels. This large search range may exceed the capabilities of the encoder.

One way of reducing the search range is to use a telescopic search technique. This is best explained by illustrating with an example. Consider the group of pictures in figure D.27.

I	B	B	B	P	B	B	B	P	B	B	B	P
0	1	2	3	4	5	6	7	8	9	10	11	12

Figure D.27 -- Example group of pictures in display order

The encoder might proceed using its selected block matching criterion and D search strategy. For each P-picture and the preceding B-pictures, it first calculates all the forward vectors, then calculates all the backward vectors. The first set of pictures consists of pictures 0 through 4.

To calculate the complete set of forward vectors, the encoder first calculates all the forward vectors from picture 0 to picture 1 using a 2-D search strategy centered on zero displacement. It next calculates all the forward vectors from picture 0 to picture 2 using a 2-D search strategy centered on the displacements calculated for the corresponding block of picture 1. It next calculates all the forward vectors from picture 0 to picture 3 using a 2-D search strategy centered on the displacements calculated for the corresponding block of picture 2. Finally, it calculates all the forward vectors from picture 0 to picture 4 using a 2-D search strategy centered on the displacements calculated for the corresponding block of picture 3.

To calculate the complete set of backward vectors, the encoder first calculates all the backward vectors from picture 4 to picture 3 using a 2-D search strategy centered on zero displacement. It next calculates all the backward vectors from picture 4 to picture 2 using a 2-D search strategy centered on the displacements calculated for the corresponding block of picture 3. Finally, it calculates all the backward vectors from picture 4 to picture 1 using a 2-D search strategy centered on the displacements calculated for the corresponding block of picture 2.

Further methods of motion estimation are given by Netravali and Haskell [1].

D.6.2.3 Coding of motion vectors

The motion vector of a macroblock tends to be well correlated with the vector of the previous macroblock. For example, in a pan all vectors would be roughly the same. Motion vectors are coded using a DPCM technique to make use of this correlation.

In P-pictures the motion vector used for DPCM, the prediction vector, is set to zero at the start of each slice and at each intra-coded macroblock. Note that macroblocks which are coded as predictive but which have no motion vector, also set the prediction vector to zero.

In B-pictures there are two motion vectors, forward and backward. Each vector is coded relative to the predicted vector of the same type. Both motion vectors are set to zero at the start of each slice and at each intra-coded macroblock. Note that predictive macroblocks which have only a forward vector do not affect the value of the predicted backward vector. Similarly, predictive macroblocks which have only a backward vector do not affect the value of the predicted forward vector.

The range of the vectors is set by two parameters. The `full_pel_forward_vector` and `full_pel_backward_vector` flags in the picture header determine whether the vectors are defined in half-pel or integer-pel units.

A second parameter, `forward_f_code` or `backward_f_code`, is related to the number of bits appended to the VLC codes in table D.9.

Table D.9 -- Differential motion code.

VLC code	Value
0000 0011 001	-16
0000 0011 011	-15
0000 0011 101	-14
0000 0011 111	-13
0000 0100 001	-12
0000 0100 011	-11
0000 0100 11	-10
0000 0101 01	-9
0000 0101 11	-8
0000 0111	-7
0000 1001	-6
0000 1011	-5
0000 11	-4
0001 1	-3
0011	-2
011	-1
1	0
010	1
0010	2
0001 0	3
0000 110	4
0000 1010	5
0000 1000	6
0000 0110	7
0000 0101 10	8
0000 0101 00	9
0000 0100 10	10
0000 0100 010	11
0000 0100 000	12
0000 0011 110	13
0000 0011 100	14
0000 0011 010	15
0000 0011 000	16

Advantage is taken of the fact that the range of displacement vector values is constrained. Each VLC represents a pair of difference values. Only one of the pair will yield a motion vector falling within the permitted range.

The range of the vector is limited to the values shown in table D.7. The values obtained by decoding the differential values must be kept within this range by adding or subtracting a modulus which depends on the f value as shown in table D.10.

Table D.10 -- Modulus for motion vectors

forward_f_code or backward_f_code	MODULUS
1	32
2	64
3	128
4	256
5	512
6	1 024
7	2 048

The use of the modulus, which refers only to the numbers in tables D.8 through D.10, will be illustrated by an example. Assume that a slice has the following vectors, expressed in the units set by the full pel flag.

3 10 30 30 -14 -16 27 24

The range is such that an f value of 2 can be used. The initial prediction is zero, so the differential values are

3 7 20 0 -44 -2 43 -3

The differential values are reduced to the range -32 to +31 by adding or subtracting the modulus 64 corresponding to the forward_f_code of 2.

3 7 20 0 20 -2 -21 -3

To create the codeword, $(mvd + (\text{sign}(mvd) * (\text{forward_f} - 1)))$ is divided by forward_f. The signed quotient of this division is used to find a variable length codeword from table D.9. Then the absolute value of the remainder is used to generate a fixed length code that is concatenated with the variable length code. The codes generated by this example are shown below:

Value	VLC Code
3	0010 0
7	0000 1100
20	0000 0100 101
0	1
20	0000 0100 101
-2	0111
-21	0000 0100 0110
-3	0011 0

D.6.3 Coding I-pictures

In coding I-pictures, the encoder has two main decisions to make that are not mandated by this part of ISO/IEC 11172. These are: how to divide the picture up into slices, and how to set the quantizer scale.

D.6.3.1 Slices in I-pictures

Division of the picture into slices is described in D.5.4.

D.6.3.2 Macroblocks in I-pictures

D.6.3.2.1 Macroblock types in I-pictures

There are two types of macroblock in I-pictures. Both use intra coding. One uses the current quantizer scale, whereas the other defines a new value for the quantizer scale. They are identified in the coded bitstream by the VLC codes given in table D.11.

Table D.11 -- Macroblock type VLC for I-pictures (table B.2a.)

TYPE	QUANT	VLC
intra-d		1
intra-q	1	01

The types are referred to names in this annex. Intra-d is the default type where the quantizer scale is not changed. Intra-q sets the quantizer scale.

In order to allow for possible future extension to MPEG video, the VLC for intra-q is 01 rather than 0. Additional types could be added to this table without interfering with the existing entries. The VLC table is thus open for future additions, and not closed. A policy of making the coding tables open in this way was adopted by in developing this part of ISO/IEC 11172. The advantage of future extension was judged to be worth the slight coding inefficiency.

D.6.3.2.2 Quantizer scale

If the macroblock type is intra-q, then the macroblock header contains a five-bit integer which defines the quantizer scale. This is used by the decoder to calculate the DCT coefficients from the transmitted quantized coefficients. A value of 0 is forbidden, so the quantizer scale can have any value between 1 and 31 inclusive.

Note that also the quantizer scale is set in a slice header.

If the block type is intra-d, then no quantizer scale is transmitted and the decoder uses the previously set value. For a discussion on strategies encoders might use to set the quantizer scale, see D.6.1.

Note that the cost of transmitting a new quantizer scale is six bits: one for the extra length of the macroblock type code, and five to define the value. Although this is normally a small fraction of the bits allocated to coding each macroblock, the encoder should exercise some restraint and avoid making a large number of very small changes.

D.6.3.3 DCT transform

The DCT is illustrated in figure D.28.

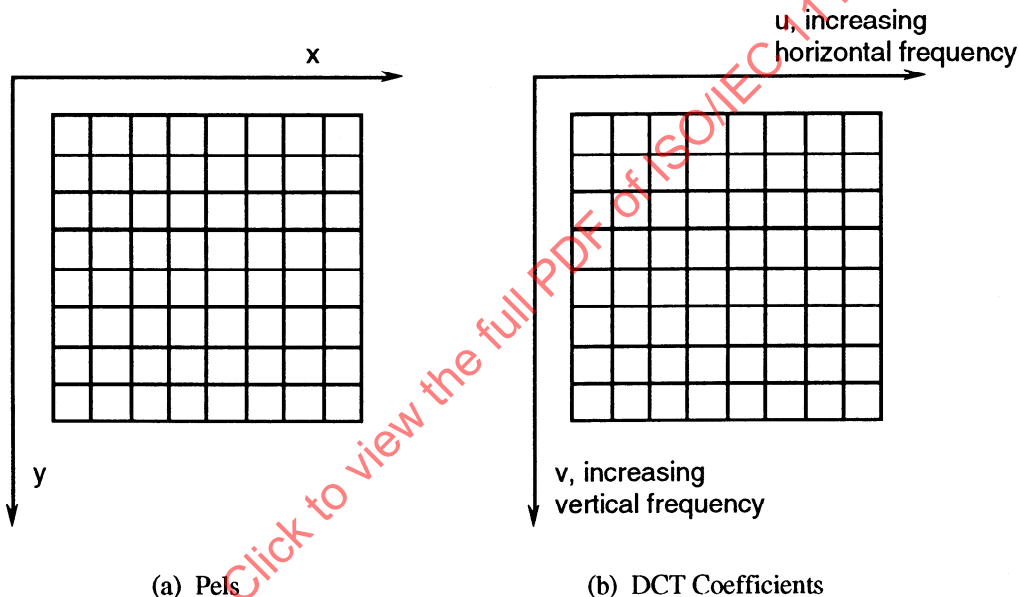


Figure D.28 -- Transformation of pels to coefficients

The pels are shown in raster scan order, whereas the coefficients are arranged in frequency order. The top left coefficient is the dc term and is proportional to the average value of the component pel values. The other coefficients are called ac coefficients. The ac coefficients to the right of the dc coefficient represent increasing horizontal frequencies, whereas ac coefficients below the dc coefficient represent increasing vertical frequencies. The remaining ac coefficients contain both horizontal and vertical frequency components. Note that an image containing only vertical lines contains only horizontal frequencies.

The coefficient array contains all the information of the pel array and the pel array can be exactly reconstructed from the coefficient array, except for information lost by the use of finite arithmetic precision.

The two-dimensional DCT is defined as

$$F(u,v) = \frac{1}{4} \sum_{x=0}^7 \sum_{y=0}^7 f(x,y) \cos(\pi(2x+1)u/16) \cos(\pi(2y+1)v/16)$$

with: $u, v, x, y = 0, 1, 2, \dots, 7$
 where x, y = spatial coordinates in the pel domain
 u, v = coordinates in the transform domain
 $C(u) = 1/\sqrt{2}$ for $u = 0$
 $C(v) = 1/\sqrt{2}$ for $v = 0$
 $= 1$ otherwise

This transform is separable, i.e. a one-dimensional DCT transform may be applied first in the horizontal direction and then in the vertical direction. The formula for the one dimensional transform is:

$$F(u) = \frac{1}{2} C(u) \sum_{x=0}^7 f(x) \cos(\pi(2x+1)u/16)$$

$$C(u) = 1/\sqrt{2} \text{ for } u = 0 \\ = 1 \text{ otherwise}$$

Fast DCT transforms exist, analogous to fast Fourier transforms. See reference [3].

The input pel values have a range from 0 to 255, giving a dynamic range for the dc coefficient from 0 to 2 040. The maximum dynamic range for any ac coefficient is about -1 000 to 1 000. Note that for P and B-pictures the component pels represent difference values and range from -255 to 255. This gives a maximum dynamic range for any coefficient of about -2 000 to 2 000. The encoder may thus represent the coefficients using 12 bits whose values range from -2 048 to 2 047.

D.6.3.4 Quantization

Each array of 8 by 8 coefficients produced by the DCT transform operation is quantized to produce an 8 by 8 array of quantized coefficients. Normally the number of non-zero quantized coefficients is quite small, and this is one of the main reasons why the compression scheme works as well as it does.

The coefficients are quantized with a uniform quantizer. The characteristic of this quantizer, only for I-blocks, is shown below:

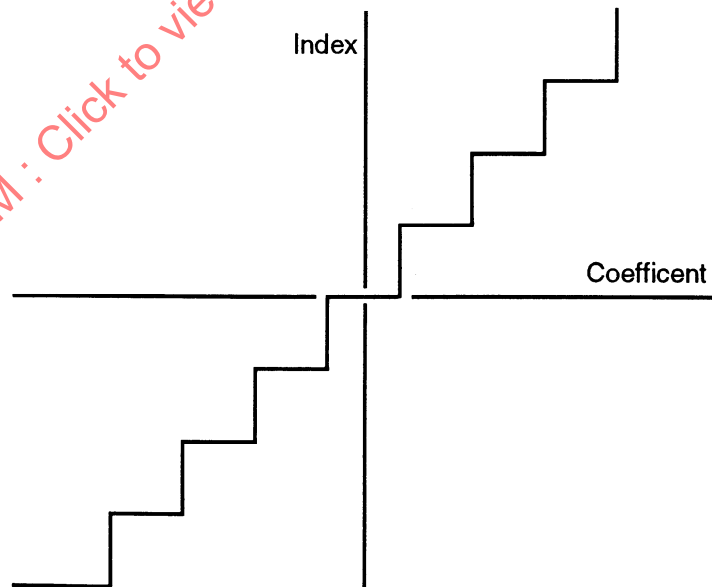


Figure D.29. -- Uniform quantizer characteristics

The value of the coefficient is divided by the quantizer step size and rounded to the nearest whole number to produce the quantized coefficient. Half integer values may be rounded up or down without directly affecting image quality. However, rounding towards zero tends to give the smallest code size and so is preferred. For example, with a step size of 16 all coefficients with values between 25 and 40 inclusive would give a quantized coefficient of 2.

The quantizer step size is derived from the quantization matrix and the quantizer scale. It can thus be different for different coefficients, and may change between macroblocks. The only exception is the dc coefficient which is treated differently.

The eye is quite sensitive to large area luminance errors, and so the accuracy of coding the dc value is fixed. The quantizer step size for the dc coefficients of the luminance and chrominance components is fixed at eight. The dc quantized coefficient is obtained by dividing the dc coefficient by eight and rounding to the nearest whole number. This effectively quantizes the average dc value to one part in 256 for the reconstructed pels.

For example, a dc coefficient of 21 is quantized to a value of 3, independent of the value of the quantizer scale.

The ac coefficients are quantized using the intra quantization matrix. The quantized coefficient $i[u,v]$ is produced by quantizing the coefficient $c[u,v]$ for I-blocks. One equation is given by the formula:

$$i[u,v] = 8 * c[u,v] // (q * m[u,v])$$

where $m[u,v]$ is the corresponding element of the intra quantization matrix, and q is the quantizer scale. The quantized coefficient is limited to the range -255 to +255.

The intra quantization matrix might be the default matrix, or it might have been downloaded in the sequence header.

D.6.3.5 Coding of quantized coefficients

The top left coefficient in figure D.28b is called the dc coefficient, the remainder are called ac coefficients. The dc coefficient is correlated with the dc coefficient of the preceding block, and advantage is taken of this in coding. The ac coefficients are not well correlated, and are coded independently.

After the dc coefficient of a block has been quantized it is coded losslessly by a DPCM technique. Coding of the luminance blocks within a macroblock follows the raster scan order of figure D.5, 0 to 3. Thus the dc value of block 3 becomes the dc predictor for block 0 of the following macroblock. The dc value of each chrominance block is coded using the dc value of the corresponding block of the previous macroblock as a predictor. At the beginning of each slice, all three dc predictors for Y, Cb and Cr, are set to 1 024 (128*8).

The differential dc values thus generated are categorized according to their absolute value as shown in table D.12.

Table D.12. -- Differential dc size and VLC

DIFFERENTIAL DC (absolute value)	SIZE	VLC CODE (luminance)	VLC CODE (chrominance)
0	0	100	00
1	1	00	01
2 to 3	2	01	10
4 to 7	3	101	110
8 to 15	4	110	1110
16 to 31	5	1110	1111 0
32 to 63	6	1111 0	1111 10
64 to 127	7	1111 10	1111 110
128 to 255	8	1111 110	1111 1110

The size is transmitted using a VLC. This VLC is different for luminance and chrominance since the statistics are different.

The size defines the number of additional bits required to define the level uniquely. Thus a size of 6 is followed by 6 additional bits. These bits define the level in order, from low to high. Thus the first of these extra bits gives the sign: 0 for negative and 1 for positive. A size of zero requires no additional bits.

The additional codes are given in table D.13.

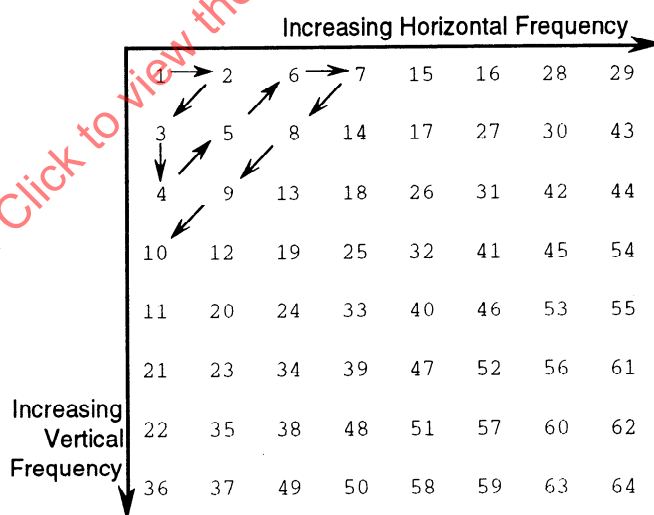
Table D.13. -- Differential dc additional code

DIFFERENTIAL DC	SIZE	ADDITIONAL CODE
-255 to -128	8	00000000 to 01111111
-127 to -64	7	0000000 to 0111111
-63 to -32	6	000000 to 011111
-31 to -16	5	00000 to 01111
-15 to -8	4	0000 to 0111
-7 to -4	3	000 to 011
3 to -2	2	00 to 01
-1	1	0
0	0	
1	1	1
2 to 3	2	10 to 11
4 to 7	3	100 to 111
8 to 15	4	1000 to 1111
16 to 31	5	10000 to 11111
32 to 63	6	100000 to 111111
64 to 127	7	1000000 to 1111111
128 to 255	8	10000000 to 11111111

For example, a luminance dc change of 10 would be coded as 1101010. table D.12 shows that the first three bits 110 indicate that the size is 4. This means that four additional bits are required to define the exact value. The next bit is a 1, and table D.13 shows that the differential dc value must be somewhere between 8 and 15 inclusive. The last three bits, 010, show that the exact value is 10.

The decoder reconstructs dc quantized coefficients by following the inverse procedure.

The ac quantized coefficients are coded using a run length and level technique. The quantized coefficients are first scanned in the zigzag order shown in figure D.30.

**Figure D.30. -- Quantized coefficient block in zigzag scan order**

The scanning order starts at 1, passes through 2, 3 etc in order, eventually reaching 64 in the bottom right corner. The length of a run is the number of zero quantized coefficients skipped over. For example, the quantized coefficients in figure D.31 produce the list of run lengths and levels in table D.14.